

Knights Corner Instruction Set Reference Manual

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Chapter 1 Introduction

This document describes new vector instructions for the co-processor code-named Knights Corner.

The major features of the new vector instructions described herein are:

- **A high performance 64 bit execution environment** Knights Corner provides a 64 bit execution environment (see Figure [2.1\)](#page-36-0) similar to that found in the Intel 64° Intel[®] Architecture Software Developer's Manual. Additionally, Knights Corner instructions provides basic support for float64 and int64 logical operations.
- **32 new vector registers** Knights Corner's 64 bit environment offers 32 512-bit wide vector SIMD registers tailored to boost the performance of high performance computing applications. The 512-bit vector SIMD instruction extensions provide comprehensive, native support to handle 32 bit and 64 bit floating-point and integer data, including a rich set of conversions for native data types.
- **Ternary instructions** Most instructions are ternary, with two sources and a different destination. Multiply&add instructions are ternary with three sources, one of which is also the destination.
- **Vector mask support** Knights Corner instructions introduces 8 vector mask registers that allow for conditional execution over the 16 (or 8) elements in a vector instruction, and merging of the results into the destination. Masks allow vectorizing loops that contain conditional statements. Additionally, Knights Corner instructions provides support for updating the value of the vector masks with special vector instructions such as *vcmpmps*.
- **Coherent memory model** The Knights Corner instructions operates in a memory address space that follows the standard defined by the *Intel[®]* 64 achitecture. This feature eases the process of developing vector code.
- **Gather/Scatter support** The Knights Corner instructions features specific gather/scatter instructions that allow manipulation of irregular data patterns of memory (by fetching sparse locations of memory into a dense vector register or vice-versa) thus enabling vectorization of algorithms with complex data structures.

Chapter 2

Instructions Terminology and State

The vector streaming SIMD instruction extensions are designed to enhance the performance of *Intel® 64* processors for scientific and engineering applications.

This chapter introduces Knights Corner instructions terminology and relevant processor state.

2.1 Overview of the Knights Corner instructions Extensions

2.1.1 What are vectors?

The vector is the basic working unit of the Knights Corner instructions. Most instructions use at least one vector. A vector is defined as a sequence of packed data elements. For Knights Corner instructions the size of a vector is 64 bytes. As the support data types are float 32, int 32, float 64 and int64, then a vector consists on either 16 doubleword-size elements or alternatively, 8 quadword-size elements. Only doubleword and quadword elements are supported in Knights Corner instructions.

The number of Knights Corner instructions registers is 32.

Additionally, Knights Corner instructions features *vector masks*. Vector masks allow any set of elements in the destination to be protected from updates during the execution of any operation. A subset of this functionality is the ability to control the vector length of the operation being performed (that is, the span of elements being modified, from the first to the last one); however, it is not necessary that the elements that are modified be consecutive.

2.1.2 Vector mask registers

Most Knights Corner instructions vector instructions use a special extra source, known as the write-mask, sourced from a set of 8 registers called *vector mask registers*. These registers contain one bit for each element that can be held by a regular Knights Corner instructions vector register.

Elements are always either float32, int32, float64 or int64 and the vector size is set to 64 bytes. Therefore, a vector register holds either 8 or 16 elements; accordingly, the length of a vector mask register is 16 bits. For 64

bit datatype instructions, only the 8 least significant bits of the vector mask register are used.

A vector mask register affects an instruction for which it is the write-mask operand at element granularity (either 32 or 64 bits). That means that every element-sized operation and element-sized destination update by a vector instruction is predicated on the corresponding bit of the vector mask register used as the write-mask operand. That has two implications:

- **The instruction's operation is not performed for an element**if the corresponding write-mask bit is not set. This implies that no exception or violation can be caused by an operation on a masked-off element.
- **A destination element is not updated** if the corresponding write-mask bit is not set. Thus, the mask in effect provides a merging behavior for Knights Corner instructions vector register destinations, thereby potentially converting destinations into implicit sources, whenever a write-mask containing any 0-bits is used.

This merging behavior, and the associated performance hazards, can also occur when writing a vector to memory via a vector store. Vectors are written on a per element basis, based on the vector mask register used as a write-mask. Therefore, no exception or violation can be caused by a write to a masked-off element of a destination vector operand.

The sticky bits implemented in the MXCSR to indicate that floating-point exceptions occurred, are set based soley upon operations on non-masked vector elements.

The value of a given mask register can be set up as a direct result of a vector comparison instruction, transferred from a GP register, or calculated as a direct result of a logical operation between two masks.

Vector mask registers can be used for purposes other than write-masking. For example, they can be used to to set the EFLAGS based on the 0/0xFFFF/other status of the OR of two vector mask registers. A number of the Knights Corner instructions are provided to support such uses of the vector mask register.

2.1.2.1 Vector mask *k*0

The only exception to the vector mask rules described above is mask *k*0. Mask *k*0 cannot be selected as the writemask for a vector operation; the encoding that would be expected to select mask *k*0 instead selects an implicit mask of 0xFFFF, thereby effectively disabling masking. Vector mask *k*0 can still be used as any non-write-mask operand for any instruction that takes vector mask operands; it just can't ever be selected as a write-mask.

2.1.2.2 Example of use

Here's an example of a masked vector operation.

The initial state of vector registers zmm0, zmm1, and zmm2 is:

Given this state, we will execute the following instruction:

vpaddd zmm2 {k3}, zmm0, zmm1

The vpaddd instruction adds vector elements of 32 bit integers. Since elements are not operated upon when the corresponding bit of the mask is not set, the temporary result would be:

where "**********" indicates that no operation is performed.

This temporary result is then written into the destination vector register, zmm2, using vector mask register k3 as the write-mask, producing the following final result:

Note that for a 64 bit instruction (say *vaddpd*), only the 8 LSB of mask k3 (0x03) would be used to identify the write-mask operation on each one of the 8 elements of the source/destination vectors.

2.1.3 Understanding Knights Corner instructions

Knights Corner instructions can be classified depending on the nature of their operands. The majority of the Knights Corner instructions operate on vector registers, with a vector mask register serving as a write-mask. However, in most cases these instructions may have one of the vector source operands stored in either memory or a vector register, and may additionally have one or more non-vector (scalar) operands, such as a *Intel® 64* general purpose register or an immediate value. Additionally, some instructions use vector mask registers as destinations and/or explicit sources. Finally, Knights Corner instructions adds some new scalar instructions.

From the point of view of instruction formats, there are four main types of Knights Corner instructions:

- Vector Instructions
- Vector Memory Instructions
- Vector Mask Instructions
- New Scalar Instructions

2.1.3.1 Knights Corner instructions Vector Instructions

Vector instructions operate on vectors that are sourced from either registers or memory and that can be modified prior to the operation via predefined swizzle and convert functions. The destination is usually a vector register, though some vector instructions may have a vector mask register as either a second destination or the primary destination.

All these instructions work in an element-wise manner: the ϐirst element of the ϐirst source vector is operated on together with the first element of the second source vector, and the result is stored in the first element of the destination vector, and so on for the remaining 15 (or 7) elements.

As described above, the vector mask() register that serves as the write-mask for a vector instruction determines which element locations are actually operated upon; the mask can disable the operation and update for any combination of element locations.

Most vector instructions have three different vector operands (typically, two sources and one destination) except those instructions that have a single source and thus use only two operands. Additionally, most vector instructions feature an extra operand in the form of the vector mask() register that serves as the write-mask. Thus, we can categorize Knights Corner instructions vector instructions based on the number of vector sources they use:

Vector-Converted Vector/Memory. Vector-converted vector/memory instructions, such as vaddps (which adds two vectors), are ternary operations that take two different sources, a vector register and a converted vector/memory operand, and a separate destination vector register, as follows:

```
zmm0 <= OP(zmm1, S(zmm2, m))
```
where zmm1 is a vector operand that is used as the first source for the instruction, $S(zmm2,m)$ is a converted vector/memory operand that is used as the second source for the instruction, and the result of performing operation OP on the two source operands is written to vector destination register *zmm*0.

A converted vector/memory operand is a source vector operand that it is obtained through the process of applying a *swizzle/conversion function* to either a Knights Corner instructions vector or a memory operand. The details of the swizzle/conversion function are found in section [2.2](#page-26-0); note that its behavior varies depending on whether the operand is a register or a memory location, and, for memory operands, on whether the instruction performs a floating-point or integer operation. Each source memory operand must have an address that is aligned to the number of bytes of memory actually accessed by the operand (that is, before the swizzle/convert is performed); otherwise, a #GP fault will result.

Converted Vector/Memory. Converted vector/memory instructions, such as vcvtpu2ps (which converts a vector of unsigned integers to a vector of floats), are binary operations that take a single vector source, as follows:

 $zmm0 \leq D P(S(zmm1, m))$

Vector-Vector-Converted Vector/Memory. Vector-vector-converted vector/memory instructions, of which vfmadd*ps (multiply-add of three vectors) is a good example, are similar to the *vector-converted vector/memory* family of instructions; here, however, the destination vector register is used as a third source as well:

zmm $0 \leq DP(zmm0, zmm1, S(zmm2, m))$

2.1.3.2 Knights Corner instructions Vector Memory Instructions:

Vector Memory Instructions perform vector loads from and vector stores to memory, with extended conversion support.

As with regular vector instructions, vector memory instructions transfer data from/to memory in an elementwise fashion, with the elements that are actually transferred dictated by the contents of the vector mask that is selected as the write-mask.

There are two basic groups of Knights Corner instructions vector memory instructions, vector loads/broadcasts and vector stores.

Vector Loads/Broadcasts. A vector load/broadcast reads a memory source, performs a predefined *load conversion* function, and replicates the result (in the case of broadcasts) to form a 64-byte 16-element vector (or 8-element for 64 bit datatypes). This vector is then conditionally written element-wise to the vector destination register, with the writes enabled or disabled according to the corresponding bits of the vector mask register selected as the write-mask.

The size of the memory operand is a function of the type of conversion and the number of replications to be performed on the memory operand. We call this special memory operand an *up-converted memory operand*. Each source memory operand must have an address that is aligned to the number of bytes of memory actually accessed by the operand (that is, before the swizzle/convert is performed); otherwise, a #GP fault will result.

A Vector Load operates as follows:

zmm $0 \le U(m)$

where $U(m)$ is an up-converted memory operand whose contents are replicated and written to destination register *zmm*0. The mnemonic dictates the degree of replication and the conversion table.

A special sub-case of these instructions are **Vector Gathers**. **Vector Gathers** are a special form of vector loads where, instead of a consecutive chunks of memory, we load a sparse set of memory operands (as many as the vector elements of the destination). Every one of those memory operands must obey the alignment rules; otherwise, a #GP fault will result if the related write-mask bit is not disabled (set to 0).

A Vector Gather operates as follows:

zmm $0 \leq U(mv)$

where $U(mv)$ is a set of up-converted memory operands described by a base address, a vector of indices and an immediate scale to apply for each index. Every one of those operands is conditionally written to destination vector *zmm*0 (based on the value of the write-mask).

Vector Stores. A vector store reads a vector register source, performs a predefined *store conversion* function, and writes the result to the destination memory location on a per-element basis, with the writes enabled or disabled according to the corresponding bits of the vector mask register selected as the write-mask.

The size of the memory destination is a function of the type of conversion associated with the mnemonic. We call this special memory operand a *down-converted memory operand*. Each memory destination operand must have an address that is aligned to the number of bytes of memory accessed by the operand (pre-conversion, if conversion is performed); otherwise, a #GP fault will result.

A Vector Store operates as follows:

 $m \le D(zmm0)$

where $zmm0$ is the vector register source whose full contents are down-converted (denoted by $D()$), and written to memory.

A special sub-case of these instructions are **Vector Scatters**. **Vector Scatters** are a special form of vector stores where, instead of writing the source vector into a consecutive chuck of memory, we store each vector element into a different memory location. Every one of those memory destinations must obey the alignment rules; otherwise, a #GP fault will result if the related write-mask bit is not disabled (set to 0).

A Vector Scatter operates as follows:

 $mv \le D(zmm0)$

where $zmm0$ is the vector register source whose full or partial contents are down-converted (denoted by $D($)), and written to the set of memory locations mv , specified by a base address, a vector of indices and an immediate scale which is applied to every index. Every one of those down-converted elements are conditionally stored in the memory locations based on the value of the write-mask.

2.1.3.3 Knights Corner instructions vector mask Instructions

Vector mask instructions allow programmers to set, copy, or operate on the contents of a given vector mask.

There are three types of vector mask instructions:

- **Mask read/write instructions:** These instruction move data between a general-purpose integer register and a vector mask register, or between two vector mask registers.
- **Flag instructions:** This category, consisting of instructions that modify EFLAGS based on vector mask registers, actually contains only one instruction, kortest.
- **Mask logical instructions:** These instructions perform standard bitwise logical operations between vector mask registers.

2.1.3.4 Knights Corner instructions New Scalar Instructions

In addition to vector, vector memory, and vector mask instructions, Knights Corner instructions adds a few scalar instructions as well. These instructions are useful for increasing the performance of some critical algorithms;

for example, any code that suffers reduced performance due to cache-miss latency can benefit from the new prefetch instructions.

2.2 Knights Corner instructions Swizzles and Converts

Data transformation, in the form of certain data conversions or element rearrangements (for loads, both at once) of one operand, can be performed for free as part of most Knights Corner instructions vector instructions.

Three sorts of data transformations are available:

- **Data Conversions:** Sources from memory can be converted to either 32 bit signed or unsigned integer or 32 bit floating-point before being used. Supported data types in memory are float16, sint8, uint8, sint16, and uint16 for load-op instructions
- **Broadcast:** If the source memory operand contains fewer than the total number of elements, it can be broadcast (repeated) to form the full number of elements of the effective source operand (16 for 32 bit instructions, 8 for 64 bit instructions). Broadcast can be combined with load-type conversions only; loadop instructions can do one or the other: either broadcast, or swizzle and/or up-conversion. There are two broadcast granularities:
	- **– 1-element granularity** where the 1 element of the source memory operand are broadcast 16 times to form a full 16-element effective source operand (for 32 bit instructions), or 8 times to form a full 8-element effective source operand (for 64 bit instructions).
	- **– 4-element granularity** where the 4 elements of the source memory operand is broadcast 4 times to form a full 16-element effective source operand (for 32 bit instructions), or 2 times to form a full 8-element effective source operand (for 64 bit instructions).

Broadcast is very useful for instructions that mix vector and scalar sources, where one of the sources is common across the different operations.

• **Swizzles:** Sources from registers can undergo swizzle transformations (that is, they can be permuted), although only 8 swizzles are available, all of which are limited to permuting within 4-element sets (either of 32 bits or 64 bits each).

Knights Corner instructions also introduces the concept of **Rounding Mode Override** or **Static (per instruction) Rounding Mode**, which efficiently supports the feature of determining the rounding mode for arithmetic operations on a per-instruction basis. Thus one can choose the rounding mode without having to perform costly MXCSR save-modify-restore operations.

Knights Corner extends the swizzle functionality for register-register operands in order to provide rounding mode override capabilities for Knights Corner floating-point instructions instead of obeying the MXCSR.RC bits. All four rounding modes are available via swizzle attribute: Round-up, Round-down, Round-toward-zero and Round-to-nearest. The option is not available for instructions with memory operands. On top of these options, Knights Corner introduces the SAE (suppress-all-exceptions) attribute feature. An instruction with SAE set will not raise any kind of floating-point exception flags, independent of the inputs.

In addition to those transformations, all Knights Corner instructions memory operands may have a special attribute, called the *EH* hint (eviction hint), that indicates to the processor that the data is non-temporal - that is, it is unlikely to be reused soon enough to benefit from caching in the 1st-level cache and should be given priority for eviction. This is, however, a hint, and the processor may implement it in any way it chooses, including ignoring the hint entirely.

Table [2.1](#page-27-1) shows the assembly language syntax used to indicate the presence or absence of the EH hint.

Data transformations can only be performed on one source operand at most; for instructions that take two or three source operands, the other operands are always used unmodified, exactly as they're stored in their source registers. In no case do any of the Knights Corner instructions allow using data conversion and swizzling at the same time. Broadcasts, on the other hand, can be combined with data conversions when performing vector loads.

Not all instructions can use all of the different data transformations. Load-op instructions (such as vector arithmetic instructions), vector loads, and vector stores have different data transformation capabilities. We can categorize these transformation capabilities into three families:

- **Load-Op SwizzUpConv:** For a register source, swizzle; for a memory operand, either: (a) broadcast, or (b) convert to 32 bit floats or 32 bit signed or unsigned integers. This is used by vector arithmetic instructions and other load-op instructions. There are two versions, one for 32 bit floating-point instructions and another for 32 bit integer instructions; in addition, the available data transformations differ for register and memory operands.
- Load UpConv: Convert from a memory operand to 32 bit floats or 32 bit signed or unsigned integers; used by vector loads and broadcast instructions. For 32 bit floats, there are three different conversion tables based on three different input types. See Section [2.2.2,](#page-29-0) Load UpConvert. There is no load conversion support for 64 bit datatypes.
- **DownConv:** Convert from 32 bit floats or 32 bit signed or unsigned integers to a memory operand; used by vector stores. For 32 bit floats, there are three different conversion tables based on three different output types. See Section [2.2.3,](#page-31-0) Down-Conversion.

There is no store conversion support for 64 bit datatypes.

2.2.1 Load-Op Swizzle/Convert

Vector load-op instructions can swizzle, broadcast, or convert one of the sources; we will refer to this as the *swizzle/convert* source, and we will use SwizzUpConv to describe the swizzle/convert function itself. The available SwizzUpConv transformations vary depending on whether the operand is memory or a register, and also in the case of conversions from memory depending on whether the vector instruction is 32 bit integer, 32 bit floating-point, 64 bit integer or 64 bit floating-point. 3 bits are used to select among the different options, so eight options are available in each case.

When the swizzle/convert source is a register, SwizzUpConv allows the choice of one of eight swizzle primitives (one of the eight being the identity swizzle). These swizzle functions work on either 4-byte or 8-byte elements within 16-byte/32-byte boundaries. For 32 bit instructions, that means certain permutations of each set of four elements (16 bytes) are supported, replicated across the four sets of four elements. When the swizzle/convert source is a register, the functionality is the same for both integer and floating-point 32 bit instructions. Table [2.2](#page-28-0) shows the available register-source swizzle primitives.

Table 2.2: **32 bit Register SwizzUpConv** swizzle primitives. Notation: dcba denotes the 32 bit elements that form one 128-bit block in the source (with 'a' least significant and 'd' most significant), so aaaa means that the least significant element of the 128-bit block in the source is replicated to all four elements of the same 128bit block in the destination; the depicted pattern is then repeated for all four 128-bit blocks in the source and destination. We use 'ponm lkji hgfe dcba' to denote a full Knights Corner instructions source register, where 'a' is the least significant element and '*p'* is the most significant element. However, since each 128-bit block performs the same permutation for register swizzles, we only show the least significant block here. Note that in this table as well as in subsequent ones from this chapter $S_2S_1S_0$ are bits 6-4 from MVEX prefix encoding (see Figure [3.3](#page-41-0)

Table 2.3: **64 bit Register SwizzUpConv** swizzle primitives. Notation: dcba denotes the 64 bit elements that form one 256-bit block in the source (with 'a' least significant and 'd' most significant), so aaaa means that the least significant element of the 256-bit block in the source is replicated to all four elements of the same 256bit block in the destination; the depicted pattern is then repeated for the two 256-bit blocks in the source and destination. We use 'hgfe dcba' to denote a full Knights Corner instructions source register, where 'a' is the least significant element and 'h' is the most significant element. However, since each 256-bit block performs the same permutation for register swizzles, we only show the least significant block here.

For 64 bit instructions, that means certain permutations of each set of four elements (32 bytes) are supported, replicated across the two sets of four elements. When the swizzle/convert source is a register, the functionality is the same for both integer and floating-point 64 bit instructions. Table [2.3](#page-28-1) shows the available register-source swizzle primitives.

When the source is a memory location, load-op swizzle/convert can perform either no transformation, 2 different broadcasts, or four data conversions. Vector load-op instructions cannot both broadcast and perform data conversion at the same time. The conversions available differ depending on whether the associated vector instruction is integer or floating-point, and whether the natural data type is 32 bit or 64 bit. (Note however that there are no load conversions for 64 bit destination data types.)

Source memory operands may have sizes smaller than 64 bytes, expanding to the full 64 bytes of a vector source by means of either broadcasting (replication) or data conversion.

Each source memory operand must have an address that is aligned to the number of bytes of memory actually accessed by the operand (that is, before conversion or broadcast is performed); otherwise, a #GP fault will result. Thus, for SwizzUpConv, any of 4-byte, 16-byte, 32-byte, or 64-byte alignment may be required.

$S_2S_1S_0$	Function:	Usage
000	no conversion	[rax]
001	broadcast 1 element (x16)	[rax] $\{1\t016\}$
010	broadcast 4 elements (x4)	[rax] {4to16}
011	float16 to float32	[rax] {float16}
100	uint ₈ to float ₃₂	[rax] $\{$ uint $8\}$
101	reserved	N/A
110	uint16 to float32	[rax] $\{$ uint $16\}$
111	sint16 to float32	[rax] $\{sint16\}$

Table 2.4: **32 bit Floating-point Load-op SwizzUpConv***f*³² swizzle/conversion primitives. We use 'ponm lkji hgfe dcba' to denote a full Knights Corner instructions source register, with each letter referring to a 32 bit element, where 'a' is the least significant element and 'p' is the most significant element. So, for example, 'dcba dcba dcba dcba' shows that the source elements are copied to the destination by replicating the lower 128 bits of the source (the four least significant elements) to each 128-bit block of the destination.

Table 2.5: **32 bit Integer Load-op SwizzUpConv***i*³² (Doubleword) swizzle/conversion primitives. We use 'ponm lkji hgfe dcba' to denote a full Knights Corner instructions source register, with each letter referring to a 32 bit element, where 'a' is the least significant element and 'p' is the most significant element. So, for example, 'dcba dcba dcba dcba' shows that the source elements are copied to the destination by replicating the lower 128 bits of the source (the four least significant elements) to each 128-bit block of the destination.

Table [2.4](#page-29-1) shows the available 32 bit floating-point swizzle primitives.

SwizzUpConv conversions to float32s are exact.

Table [2.5](#page-29-2) shows the available 32 bit integer swizzle primitives.

Table [2.6](#page-30-0) shows the available 64 bit floating-point swizzle primitives.

Finally, Table [2.7](#page-30-1) shows the available 64 bit integer swizzle primitives.

2.2.2 Load Up-convert

Vector load/broadcast instructions can perform a wide array of data conversions on the data being read from memory, and can additionally broadcast (replicate) that data across the elements of the destination vector register depending on the instructions. The type of broadcast depends on the opcode/mnemonic being used. We

Table 2.6: **64 bit Floating-point Load-op SwizzUpConv***f*⁶⁴ swizzle/conversion primitives. We use 'hgfe dcba' to denote a full Knights Corner instructions source register, with each letter referring to a 64 bit element, where 'a' is the least significant element and 'h' is the most significant element. So, for example, 'dcba dcba' shows that the source elements are copied to the destination by replicating the lower 256 bits of the source (the four least significant elements) to each 256-bit block of the destination.

Table 2.7: **64 bit Integer Load-op SwizzUpConv***i*⁶⁴ (Quadword) swizzle/conversion primitives. We use 'hgfe dcba' to denote a full Knights Corner instructions source register, with each letter referring to a 64 bit element, where 'a' is the least significant element and 'h' is the most significant element. So, for example, 'dcba dcba' shows that the source elements are copied to the destination by replicating the lower 256 bits of the source (the four least significant elements) to each 256-bit block of the destination.

will refer to this conversion process as up-conversion, and we will use UpConv to describe the load conversion function itself.

Based on that, load instructions could be divided into the following categories:

- *regular loads*: load 16 elements (32 bits) or 8 elements (64 bits), convert them and write into the destination vector
- *broadcast 4-elements*: load 4 elements, convert them (possible only for 32 bit data types), replicate them four times (32 bits) or two times (64 bits) and write into the destination vector
- *broadcast 1-element*: load 1 element, convert it (possible only for 32 bit data types), replicate it 16 times (32 bits) or 8 times (64 bits) and write into the destination vector

Therefore, unlike load-op swizzle/conversion, Load UpConv can perform both data conversion and broadcast simultaneously. We will refer to this process as *up-conversion*, and we will use Load UpConv to describe the load conversion function itself.

When a *broadcast 1-element* is selected, the memory data, after data conversion, has a size of 4 bytes, and is broadcast 16 times across all 16 elements of the destination vector register. In other words, one vector element

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is fetched from memory, converted to a 32 bit float or integer, and replicated to all 16 elements of the destination register. Using the notation where the contents of the source register are denoted {ponm lkji hgfe dcba}, with each letter referring to a 32 bit element ('a' being the least significant element and 'p' being the most significant element), the source elements map to the destination register as follows:

{aaaa aaaa aaaa aaaa}

When *broadcast 4-element* is selected, the memory data, after data conversion, has a size of 16 bytes, and is broadcast 4 times across the four 128-bit sets of the destination vector register. In other words, four vector elements are fetched from memory, converted to four 32 bit floats or integers, and replicated to all four 4-element sets in the destination register. For this broadcast, the source elements map to the destination register as follows:

{dcba dcba dcba dcba}

Table [2.8](#page-31-1) shows the different 32 bit Load up-conversion instructions in function of the broadcast function and the conversion datatype. Similarly, Table [2.10](#page-32-2) shows the different 64 bit Load up-conversion instructions in function of the broadcast function and datatype.

Table 2.8: **32 bit Load UpConv** load/broadcast instructions per datatype. Elements may be 1, 2, or 4 bytes in memory prior to data conversion, after which they are always 4 bytes. We use 'ponm lkji hgfe dcba' to denote a full Knights Corner instructions source register, with each letter referring to a 32 bit element, where 'a' is the least significant element and 'p' is the most significant element. So, for example, 'dcba dcba dcba dcba' shows that the source elements are copied to the destination by replicating the lower 128 bits of the source (the four least significant elements) to each 128-bit block of the destination.

As with SwizzUpConv, UpConv may have source memory operands with sizes smaller than 64-bytes, which are expanded to a full 64-byte vector by means of broadcast and/or data conversion. Each source memory operand must have an address that is aligned to the number of bytes of memory actually accessed by the operand (that is, before conversion or broadcast is performed); otherwise, a #GP fault will result. Thus, any of 1-byte, 2-byte, 4-byte, 8-byte, 16-byte, 32-byte, or 64-byte alignment may be required.

Table [2.9](#page-32-1) shows the available data conversion primitives for 32 bit Load UpConv and for the different datatypes supported.

Table [2.11](#page-33-0) shows the 64 bit counterpart of Load UpConv. As shown, no 64 bit conversions are available but the pure "no-conversion" option.

2.2.3 Down-Conversion

Vector store instructions can perform a wide variety of data conversions to the data on the way to memory. We will refer to this process as *down-conversion*, and we will use DownConv to describe the store conversion function itself.

DownConv may have destination memory operands with sizes smaller than 64 bytes, as a result of data conversion. Each destination memory operand must have an address that is aligned to the number of bytes of memory actually accessed by the operand (that is, after data conversion is performed); otherwise, a #GP fault will result.

UpConv*f*³² (FP32)

Table 2.9: **32 bit Load UpConv** conversion primitives.

Table 2.10: **64 bit Load UpConv** load/broadcast instructions per datatype. Elements are always 8 bytes. We use 'hgfe dcba' to denote a full Knights Corner instructions source register, with each letter referring to a 64 bit element, where 'a' is the least significant element and 'h' is the most significant element. So, for example, 'dcba dcba' shows that the source elements are copied to the destination by replicating the lower 256 bits of the source (the four least significant elements) to each 256-bit block of the destination.

Thus, any of 1-byte, 2-byte, 4-byte, 8-byte, 16-byte, 32-byte, or 64-byte alignment may be required.

Table [2.12](#page-33-1) shows the available data conversion primitives for 32 bit DownConv and for the different supported datatypes.

Table [2.13](#page-34-1) shows the 64 bit counterpart of DownConv. As shown, no 64 bit conversions are available but the pure "no-conversion" option.

UpConv*f*⁶⁴ (FP64)

Table 2.11: **64 bit Load UpConv** conversion primitives.

DownConv*i*³² (INT32)

DownConv*f*³² (FP32)

Table 2.12: 32 bit DownConv conversion primitives. Unless otherwise noted, all conversions from floatingpoint use MXCSR.RC

DownConv*i*⁶⁴ (INT64)

Table 2.13: **64 bit DownConv** conversion primitives.

2.3 Static Rounding Mode

As described before, Knights Corner introduces a new instruction attribute on top of the normal register swizzles called *Static (per instruction) Rounding Mode* or *Rounding Mode override*. This attribute allows statically applying a specific arithmetic rounding mode ignoring the value of RM bits in MXCSR.

Static Rounding Mode can be enabled in the encoding of the instruction by setting the *EH* bit to 1 in a registerregister vector instruction. Table [2.14](#page-35-0) shows the available rounding modes and their encoding. On top of the rounding-mode, Knights Corner also allows to set the SAE ("suppress-all-exceptions") attribute, to disable reporting any floating-point exception flag on MXCSR. This option is available, even if the instruction does not perform any kind of rounding.

Note that some instructions already allow to specify the rounding mode statically via immediate bits. In such case, the immediate bits take precedence over the swizzle-specified rounding mode (in the same way that they take precedence over the MXCSR.RC setting).

Table 2.14: **Static Rounding-Mode Swizzle** available modes plus SAE.

2.4 Knights Corner Execution Environments

Knights Corner's support for 32 bit and 64 bit execution environments are similar to those found in the Intel64[®] Intel® Architecture Software Developer's Manual. The 64 bit execution environment of Knights Corner is shown in Figure [2.1.](#page-36-0) The layout of 512-bit vector registers and vector mask registers are shown in Figure [2.2](#page-37-0). This section describes new features associated with the 512-bit vector registers and the 16 bit vector mask registers.

Knights Corner instructions defines two new sets of registers that hold the new vector state. The Knights Corner instructions extension uses the vector registers, the vector mask registers and/or the *x86 64* general purpose registers.

- **Knights Corner instructions Vector Registers.** The 32 registers each store store 16 doubleword/single precision floating-point entries (or 8 quadword/double precision floating-point entries), and serve as source and destination operands for vector packed floating point and integer operations. Additionally, they may also contain memory pointer offsets used to gather and scatter data from/to memory. These registers are referenced as zmm0 through zmm31.
- **Vector Mask Registers.** These registers specify which vector elements are operated on and written for Knights Corner instructions vector instructions. If the Nth bit of a vector mask register is set, then the Nth element of the destination vector is overridden with the result of the operation; otherwise, the element remains unchanged. A vector mask register can be set using vector compare instructions, instructions to move contents from a GP register, or a special subset of vector mask arithmetic instructions.

Knights Corner vector instructions are able to report exceptions via MXCSR flags but never cause traps as all SIMD floating-point exceptions are always masked (unlike Intel® SSE/Intel® AVX instructions in other processors, that may trap if floating-point exceptions are unmasked, depending on the value of the OM/UM/IM/PM/DM/ZM bits). The reason is that Knights Corner forces the new DUE bit (Disable Unmasked Exceptions) in the MXCSR (bit21) to be set to 1.

On Knights Corner, both single precision and double precision floating-point instructions use MXCSR.DAZ and MXCSR.FZ to decide whether to treat input denormals as zeros or to flush tiny results to zero (the latter are in most cases - but not always - denormal results which are flushed to zero when MXCSR.FZ is set to 1; see the IEEE Standard 754-2008, section 7.5, for a definition of tiny floating-point results).

Table [2.15](#page-38-0) shows the bit layout of the MXCSR control register.

MXCSR bit 20 is reserved, however it is not reported as Reserved by MXCSR MASK. Setting this bit will result in undefined behavior

General-purpose registers. The sixteen general-purpose registers are available in Knights Corner's 64 bit

Figure 2.1: 64 bit Execution Environment

mode execution environment. These registers are identical to those available in the 64 bit execution environment described in the Intel64® Intel® Architecture Software Developer's Manual.

- **EFLAGS register.** R/EFLAGS are updated by instructions according to the Intel64® Intel® Architecture Software Developer's Manual. Additionally, it is also updated by Knights Corner's KORTEST instruction.
- **FCW and FSW registers.** Used by x87 instruction set extensions to set rounding modes, exception masks and flags in the case of the FCW, and to keep track of exceptions in the case of the FSW.
- **x87 stack.** An eight-element stack used to perform floating-point operations on 32/64/80-bit floating-point data using the x87 instruction set.

Table 2.15: **MXCSR** bit layout. Note: MXCSR bit 20 is reserved, however it is not reported as Reserved by MXCSR_MASK. Setting this bit will result in undefined behavior

Chapter 3

Knights Corner Instruction Format

This chapter describes the instruction encoding format and assembly instruction syntax of new instructions supported by Knights Corner.

3.1 Overview

Knights Corner introduces 512-bit vector instructions operating on 512-bit vector registers (zmm0-zmm31), and offers vector mask registers (k0-k7) to support a rich set of conditional operations on data elements within the zmm registers. Vector instructions operating on zmm registers are encoded using a multi-byte prefix encoding scheme, with 62H being the 1st of the multi-byte prefix. This multi-byte prefix is referred to as MVEX in this document.

Instructions operating on the vector mask registers are encoded using another multi-byte prefix, with C4H or C5H being the 1st of the multi-byte prefix. This multi-byte prefix is similar to the VEX prefix that is defined in the "Intel® Architecture Instruction Set Architecture Programming Reference". We will refer to the C4H/C5H based VEX-like prefix as "VEX" in this document. Additionally, Knights Corner also provides a handful of new instructions operating on general-purpose registers but are encoded using VEX. In some cases, new scalar instructions supported by Knights Corner can be encoded with either MVEX or VEX.

3.2 Instruction Formats

Instructions encoded by MVEX have the format shown in Figure [3.1](#page-40-0).

Instructions encoded by VEX have the format shown in Figure [3.2.](#page-40-1)

3.2.1 MVEX/VEX and the LOCK prefix

Any MVEX-encoded or VEX-encoded instruction with a LOCK prefix preceding the multi-byte prefix will generate an invalid opcode exception (#UD).

Figure 3.2: New Instruction Encoding Format with VEX Prefix

3.2.2 MVEX/VEX and the 66H, F2H, and F3H prefixes

Any MVEX-encoded or VEX-encoded instruction with a 66H, F2H, or F3H prefix preceding the multi-byte prefix will generate an invalid opcode exception (#UD).

3.2.3 MVEX/VEX and the REX prefix

Any MVEX-encoded or VEX-encoded instruction with a REX prefix preceding the multi-byte prefix will generate an invalid opcode exception (#UD).

3.3 The MVEX Prefix

The MVEX prefix consists of four bytes that must lead with byte 62H. An MVEX-encoded instruction supports up to three operands in its syntax and is operating on vectors in vector registers or memory using a vector mask register to control the conditional processing of individual data elements in a vector. Swizzling, conversion and other operations on data elements within a vector can be encoded with bit fields in the MVEX prefix, as shown in Figure [3.3](#page-41-0). The functionality of these bit fields is summarized below:

- 64 bit mode register specifier encoding (R, X, B, R', W, V') for memory and vector register operands (encoded in 1's complement form).
	- **–** A vector register as source or destination operand is encoded by combining the R'R bits with the reg field, or the XB bits with the r/m field of the modR/M byte.
	- **–** The base of a memory operand is a general purpose register encoded by combining the B bit with the r/m field. The index of a memory operand is a general purpose register encoded by combining the X bit with the SIB.index field.

– The vector index operand in the gather/scatter instruction family is a vector register, encoded by combining the VX bits with the SIB.index field. MVEX.vvvv is not used in the gather/scatter instruction family.

RXBR': 64-bit mode register specifier associated with reg and r/m encoding in 1's complement form.

mmmm:

0000: Reserved for future use (will #UD)

0001: implied OF leading opcode byte

- 0010: implied OF 38 leading opcode bytes
- 0011: implied OF 3A leading opcode bytes
- 0100-1111: Reserved for future use (will #UD)

W: Opcode extension or 64-bit osize (operand size) promotion.

V'vvvv: A non-destructive register specifier (in 1's complement form) or 11111 if unused.

pp: Compaction of 66/F2/F3 prefix

00: None $01:66$ 10: F3 11: F₂

E: Non-temporal/eviction hint.

SSS: Swizzle/broadcast/up-convert/down-convert/static-rounding controls.

kkk: Vector mask register for masking control.

Figure 3.3: MVEX bitfields

- Non-destructive source register specifier (applicable to the three operand syntax): This is the first source operand in the three-operand instruction syntax. It is represented by the notation, MVEX.vvvv. It can encode any of the lower 16 zmm vector registers, or using the low 3 bits to encode a vector mask register as a source operand. It can be combined with V to encode any of the 32 zmm vector registers
- Vector mask register and masking control: The MVEX.aaa field encodes a vector mask register that is used in controlling the conditional processing operation on the data elements of a 512-bit vector instruction. The MVEX.aaa field does not encode a source or a destination operand. When the encoded value of MVEX.aaa is 000b, this corresponds to "no vector mask register will act as conditional mask for the vector instruction".
- Non-temporal/eviction hint. The MVEX.E field can encode a hint to the processor on a memory referencing instruction that the data is non-temporal and can be prioritized for eviction. When an instruction encoding

does not reference any memory operand, this bit may also be used to control the function of the MVEX.SSS field.

- \bullet Compaction of legacy prefixes (66H, F2H, F3H): This is encoded in the MVEX.pp field.
- Compaction of two-byte and three-byte opcode: This is encoded in the MVEX.mmmm field.
- Register swizzle/memory conversion operations (broadcast/up-convert/down-convert)/static-rounding override: This is encoded in the MVEX.SSS field.
	- **–** Swizzle operation is supported only for register-register syntax of 512-bit vector instruction, and requires MVEX.E = 0, the encoding of MVEX.SSS determines the exact swizzle operation - see Section [2.2](#page-26-0)
	- **–** Static rounding override only applies to register-register syntax of vector ϐloating-point instructions, and requires MVEX.E = 1.

The MVEX prefix is required to be the last prefix and immediately precedes the opcode bytes.

3.3.1 Vector SIB (VSIB) Memory Addressing

In the gather/scatter instruction family, an SIB byte that follows the ModR/M byte can support VSIB memory addressing to an array of linear addresses. VSIB memory addressing is supported only with the MVEX prefix.

In VSIB memory addressing, the SIB byte consists of:

- \bullet The scale field (bit 7:6), which specifies the scale factor.
- The index field (bits 5:3), which is prepended with the 2-bit logical value of the MVEX.VX bits to specify the vector register number of the vector index operand; each element in the vector register specifies an index.
- The base field (bits 2:0) is prepended with the logical value of MVEX.B field to specify the register number of the base register.

3.4 The VEX Prefix

The VEX prefix is encoded in either the two-byte form (the first byte must be C5H) or in the three-byte form (the first byte must be C4H). Beyond the first byte, the VEX prefix consists of a number of bit fields providing specific capability; they are shown in Figure [3.4.](#page-43-0)

The functionality of the bit fields is summarized below:

• 64 bit mode register specifier encoding (R, X, B, W) : The $R/X/B$ bit field is combined with the lower three bits or register operand encoding in the modR/M byte to access the upper half of the 16 registers available in 64 bit mode. The VEX.R, VEX.X, VEX.B fields replace the functionality of REX.R, REX.X, REX.B bit fields. The W bit either replaces the functionality of REX.W or serves as an opcode extension bit. The usage of the VEX.WRXB bits is explained in detail in section 2.2.1.2 of the Intel® 64 and IA-32 Architectures Software developer's manual, Volume 2A. This bit is stored in 1's complement form (bit inverted format).

- Non-destructive source register specifier (applicable to three operand syntax): this is the first source operand in the instruction syntax. It is represented by the notation, VEX.vvvv. It can encode any generalpurpose register, or using only 3 bits it can encode vector mask registers. This field is encoded using 1's complement form (bit inverted form), i.e. RAX/K0 is encoded as 1111B, and R15 is encoded as 0000B.
- Compaction of legacy prefixes (66H, F2H, F3H): This is encoded in the VEX.pp field.
- Compaction of two-byte and three-byte opcode: This is encoded in the VEX.mmmmm field.

The VEX prefix is required to be the last prefix and immediately precedes the opcode bytes. It must follow any other prefixes. If the VEX prefix is present a REX prefix is not supported.

RXB: 64-bit mode register specifier associated with reg and r/m operand encoding. m-mmmm:

00000: Reserved for future use (will #UD) 00001: implied OF leading opcode byte 00010: implied OF 38 leading opcode bytes 00011: implied OF 3A leading opcode bytes 00100-11111: Reserved for future use (will #UD)

W: Opcode extension or 64-bit osize (operand size) promotion.

vvvv: A non-destructive register specifier (in 1's complement form) or 1111 if unused.

pp: Compaction of 66/F2/F3 prefix

00: None $01:66$ 10: F3 11: F2

Figure 3.4: VEX bitfields

3.5 Knights Corner instructions Assembly Syntax

Knights Corner instructions supports up to three operands. The rich encoding fields for swizzle/broadcast/convert/rounding, masking control, and non-temporal hint are expressed as modifier expressions to the respective operands in the assembly syntax. A few common forms for Knights Corner assembly instruction syntax are expressed in the general form:

```
mnemonic vreg{masking modifier}, source1, transform_modifier(vreg/mem)
mnemonic vreg{masking modifier}, source1, transform_modifier(vreg/mem), imm
mnemonic mem{masking modifier}, transform_modifier(vreg)
```
The specific forms to express assembly syntax operands, modifiers, and transformations are listed in Table [3.1.](#page-46-0)

3.6 Notation

The notation used to describe the operation of each instruction is given as a sequence of control and assignment statements in C-like syntax. This document only contains the notation specifically needed for vector instructions. Standard *Intel® 64* notation may be found at *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2* for convenience.

When instructions are represented symbolically, the following notations are used: *label: mnemonic argument1 {write-mask}, argument2, argument3, argument4, ...*

where:

- A *mnemonic* is a reserved name for a class of instruction opcodes which have the same function.
- The operands *argument1*, *argument2*, *argument3*, *argument4*, and so on are optional. There may be from one to three register operands, depending on the opcode. The leftmost operand is always the destination; for certain instructions, such as vfmadd231ps, it may be a source as well. When the second leftmost operand is a vector mask register, it may in certain cases be a destination as well, as for example with the vpsubrsetbd instruction. All other register operands are sources. There may also be additional arguments in the form of immediate operands; for example, the vcvtfxpntdq2ps instructions has a 3-bit immediate field that specifies the exponent adjustment to be performed, if any. The *write-mask* operand specifies the vector mask mask register used to control the selective updating of elements in the destination register or registers.

3.6.1 Operand Notation

In this manual we will consider vector registers from several perspectives. One perspective is is as an array of 64 bytes. Another is as an array of 16 doubleword elements. Another is an array of 8 quadword elements. Yet another is as an array of 512 bits. In the mnemonic operation description pseudo-code, registers will be addressed using bit ranges, such as:

 $i = n*32$ zmm1[i+31:i]

This example refers to the 32 bits of the n-th doubleword element of vector register zmm1.

We will use a similar bit-oriented notation to describe access to vector mask registers. In the case of vector mask registers, we will usually specify a single bit, rather than a range of bits, because vector mask registers are used for predication, carry, borrow, and comparison results, and a single bit per element is enough for any of those purposes.

Using this notation, it is for example possible to test the value of the 12*th* bit in k1 as follows:

if ($k1[11] == 1$) { ... code here ... }

Tables [3.1](#page-46-0) and [3.2](#page-46-1) summarize the notation used for instruction operands and their values.

In Knights Corner instructions, the contents of vector registers are variously interpreted as floating-point values (either 32 or 64 bits), integer values, or simply doubleword values of no particular data type, depending on the instruction semantics.

3.6.2 The Displacement Bytes

Knights Corner introduces a brand new displacement representation that allows for a more compact encoding in unrolled code: compressed displacement of 8-bits, or disp8*N. Such compressed displacement is based on the assumption that the effective displacement is a multiple of the granularity of the memory access, and hence we do not need to encode the redundant low-order bits of the address offset.

Knights Corner instructions using the MVEX prefix (i.e. using encoding 62) have the following displacement options:

- No displacement
- 32 bit displacement: this displacement works exactly the same as the legacy 32 bit displacement and works at byte granularity
- Compressed 8 bit displacement (disp8*N): this displacement format substitutes the legacy 8-bit displacement in Knights Corner instructions using map 62. This displacement assumes the same granularity as the memory operand size (which is dependent on the instructions and the memory conversion function being used). Redundant low-order bits are ignored and hence, 8-bit displacements are reinterpreted so that they are multiplied by the memory operands total size in order to generate the final displacement to be used in calculating the effective address.

Note that the displacements in the MVEX vector instruction prefix are encoded in exactly the same way as regular displacements (so there are no changes in the ModRM/SIB encoding rules), with the only exception that disp8 is overloaded to disp8*N. In other words there are no changes in the encoding rules or encoding lengths, but only in the interpretation of the displacement value by hardware (which needs to scale the displacement by the size of the memory operand to obtain a byte-wise address offset).

3.6.3 Memory size and disp8*N calculation

Table [3.3](#page-47-0) and Table [3.4](#page-48-0) show the size of the vector (or element) being accessed in memory, which is equal to the scaling factor for compressed displacement (disp8*N). Note that some instructions work at element granularity

Table 3.1: Operand Notation

Table 3.2: Vector Operand Value Notation

instead of full vector granularity at memory level, and hence should use the "element level" column in Table [3.3](#page-47-0) and Table [3.4](#page-48-0) (namely VLOADUNPACK, VPACKSTORE, VGATHER, and VSCATTER instructions).

Table 3.3: Size of vector or element accessed in memory for upconversion

Table 3.4: Size of vector or element accessed in memory for downconversion

3.7 EH hint

All vector instructions that access memory provide the option of specifying a cache-line eviction hint, EH.

EH is a performance hint, and may operate in different ways or even be completely ignored in different hardware implementations. Knights Corner is designed to provide support for cache-efficient access to memory locations that have either low temporal locality of access or bursts of a few very closely bunched accesses.

There are two distinct modes of EH hint operation, one for prefetching and one for loads, stores, and load-op

instructions.

The interaction of the EH hint with prefetching is summarized in Table [3.5](#page-49-0).

Table 3.5: Prefetch behavior based on the **EH** (cache-line eviction hint)

The above table describes the effect of the EH bit on gather/scatter prefetches into the targeted cache (e.g. L1 for vgatherpf0dps, L2 for vgatherpf1dps). If vgatherpf0dps misses both L1 and L2, the resulting prefetch into L1 is a non-temporal prefetch into way #N of L1, but the prefetch into L2 is a normal prefetch, not a non-temporal prefetch. If you want the data to be non-temporally fetched into L2, you must use vgatherpf1dps with the EH bit set.

The operation of the EH hint with prefetching is designed to limit the cache impact of streaming data.

Note that regular prefetch instructions (like vprefetch0) do not have an embedded EH hint. Instead, the nontemporal hint is given by the opcode/mnemonic (see VPREFETCHNTA/0/1/2 descriptions for details). The same rules described in Table [3.5](#page-49-0) still apply.

Table [3.6](#page-49-1) summarizes the interaction of the EH hint with load and load-op instructions.

Table 3.6: Load/load-op behavior based on the **EH** bit.

The EH bit, when used with load and load-op instructions, affects only the L1 cache behavior. Any resulting L2 misses are handled normally, regardless of the setting of the EH bit.

Table [3.7](#page-49-2) summarizes the interaction of the EH hint with store instructions. Note that stores that write a full cache-line (no mask, no down-conversion) evict the line from L1 (invalidation) while updating the contents directly into the L2 cache. In any other case, a store with an EH hint works as a load with an EH hint.

Table 3.7: Store behavior based on the **EH** bit.

The EH bit, when used with load and load-op instructions, affects only the L1 cache behavior. Any resulting L2 misses are handled normally, regardless of the setting of the EH bit.

intel. 3.8 Functions and Tables Used

Some mnemonic definitions use auxiliary tables and functions to ease the process of describing the operations of the instruction. The following section describes those tables and functions that do not have an obvious meaning.

3.8.1 MemLoad and MemStore

This document uses two functions, Mem-Load and MemStore, to describe in pseudo-code memory transfers that involve no conversions or broadcasts:

- **MemLoad**: Given an address pointer, this function returns the associated data from memory. Size is de-fined by the explicit destination size in the pseudo-code (see for example LDMXCSR in Appendix [B](#page-655-0))
- MemStore: Given an address pointer, this function stores the associated data to memory. Size is defined by the explicit source data size in the pseudo-code.

3.8.2 SwizzUpConvLoad, UpConvLoad and DownConvStore

In this document, the detailed discussions of memory-accessing instructions that support datatype conversion and/or broadcast (as defined by the UpConv, SwizzUpConv, and DownConv tables in section [2.2\)](#page-26-0) use the functions shown in Table [3.8](#page-50-0) in their Operation sections (the instruction pseudo-code). These functions are used to describe any swizzle, broadcast, and/or conversion that can be performed by the instruction, as well as the actual load in the case of SwizzUpConv and UpConv. Note that *zmm*/*m* means that the source may be either a vector operand or a memory operand, depending on the ModR/M encoding.

Table 3.8: SwizzUpConv, UpConv and DownConv function conventions

The Operation section may use *UpConvSizeOf*, which returns the ϐinal size (in bytes) of an up-converted memory element given a specified up-conversion mode. A specific subset of a memory stream may be used as a parameter for UpConv; Size of the subset is inferred by the size of destination together with the up-conversion mode.

Additionally, the Operation section may also use *DownConvStoreSizeOf*, which returns the final size (in bytes) of a downcoverted vector element given a specified down-conversion mode. A specific subset of a vector register

may be used as a parameter for DownConvStore; for example, DownConvStore(zmm2[31:0]) specifies that the low 32 bits of zmm2 form the parameter for DownConv.

3.8.3 Other Functions/Identifiers

The following identifiers are used in the algorithmic descriptions:

- **Carry** The carry bit from an addition.
- **FpMaxAbs** The greater of the absolute values of two floating-point numbers. See the description of the VGMAXABSPS instruction for further details.
- **FpMax** The greater of two floating-point numbers. See the description of the VGMAXPS instruction for further details.
- **FpMin** The lesser of two floating-point numbers. See the description of the VGMINPS instruction for further details.
- **Abs** The absolute value of a number.
- **IMax** The greater of two signed integer numbers.
- **UMax** The greater of two unsigned integer numbers.
- **IMin** The lesser of two signed integer numbers.
- **UMin** The lesser of two unsigned integer numbers.
- **CvtInt32ToFloat32** Convert a signed 32 bit integer number to a 32 bit floating-point number.
- **CvtInt32ToFloat64** Convert a signed 32 bit integer number to a 64 bit floating-point number.
- **CvtFloat32ToInt32** Convert a 32 bit floating-point number to a 32 bit signed integer number using the specified rounding mode.
- **CvtFloat64ToInt32** Convert a 64 bit floating-point number to a 32 bit signed integer number using the specified rounding mode.
- **CvtFloat32ToUint32** Convert a 32 bit floating-point number to a 32 bit unsigned integer number using the specified rounding mode.
- **CvtFloat64ToUint32** Convert a 64 bit floating-point number to a 32 bit unsigned integer number using the specified rounding mode.
- **CvtFloat32ToFloat64** Convert a 32 bit floating-point number to a 64 bit floating-point number.
- **CvtFloat64ToFloat32** Convert a 64 bit floating-point number to a 32 bit floating-point number using the specified rounding mode.
- **CvtUint32ToFloat32** Convert an unsigned 32 bit integer number to a 32 bit floating-point number.
- **CvtUint32ToFloat64** Convert an unsigned 32 bit integer number to a 64 bit floating-point number.
- **GetExp** Obtains the (un-biased) exponent of a given floating-point number, returned in the form of a 32 bit floating-point number. See the description of the VGETEXPPS instruction for further details.

- **RoundToInt** Rounds a floating-point number to the nearest integer, using the specified rounding mode. The result is a floating-point representation of the rounded integer value.
- **Borrow** The borrow bit from a subtraction.
- **ZeroExtend** Returns a value zero-extended to the operand-size attribute of the instruction.
- **FlushL1CacheLine** Flushes the cache line containing the specified memory address from L1.
- InvalidateCacheLine Invalidate the cache line containing the specified memory address from the whole memory cache hierarchy.
- **FetchL1CacheLine** Prefetches the cache line containing the specified memory address into L1. See the description of the VPREFETCH1 instruction for further details.
- **FetchL2CacheLine** Prefetches the cache line containing the specified memory address into L2. See the description of the VPREFETCH2 instruction for further details.

Chapter 4

Floating-Point Environment, Memory Addressing, and Processor State

This chapter describes the Knights Corner vector floating-point instruction exception behavior and interactions related to system programming.

4.1 Overview

Knights Corner 512-bit vector instructions that operate on floating-point data may signal exceptions related to arithmetic processing. When SIMD floating-point exceptions occur, Knights Corner supports exception reporting using exception flags in the MXCSR register, but traps (unmasked exceptions) are not supported.

Exceptions caused by memory accesses apply to vector floating-point, vector integer, and scalar instructions.

The MXCSR register (see Figure [4.1\)](#page-54-0) in Knights Corner provides:

- Exception flags to indicate SIMD floating-point exceptions signaled by floating-point instructions operating on zmm registers. The flags are: IE, DE, ZE, OE, UE, PE.
- Rounding behavior and control: DAZ, FZ and RC.
- Exception Suppression: DUE (always 1)

4.1.1 Suppress All Exceptions Attribute (SAE)

Knights Corner instructions that process floating-point data support a specific feature to disable floating-point exception signaling, called SAE ("suppress all exceptions"). The SAE mode is enabled via a specific bit in the register swizzle field of the MVEX prefix (by setting the EH bit to 1). When SAE is enabled in the instruction encoding, that instruction does not report any SIMD floating-point exception in the MXCSR register. This feature is only available to the register-register format of the instructions and in combination with static rounding-mode.

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4.1.2 SIMD Floating-Point Exceptions

SIMD floating-point exceptions are those exceptions that can be generated by Knights Corner instructions that operate on floating-point data in zmm operands. Six classes of SIMD floating-point exception flags can be signaled:

- Invalid operation (#I)
- Divide-by-zero (#Z)
- Numeric overflow (#0)
- Numeric underflow (#U)
- Inexact result (Precision) (#P)
- Denormal operand (#D)

4.1.3 SIMD Floating-Point Exception Conditions

The following sections describe the conditions that cause SIMD floating-point exceptions to be signaled, and the masked response of the processor when these conditions are detected.

When more than one exception is encountered, then the following precedence rules are applied 1 .

- 1. Invalid-operation exception caused by sNaN operand
- 2. Any other invalid exception condition different from sNaN input operand

¹Note that Knights Corner has no support for unmasked exceptions, so in this case the exception precedence rules have no effect. All concurrently-encountered exceptions will be reported simultaneously.

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- 3. Denormal operand exception
- 4. A divide-by-zero exception
- 5. Overflow/underflow exception
- 6. Inexact result

All Knights Corner instructions floating-point exceptions are precise and are reported as soon as the instruction completes execution. The status flags from the MXCSR register set by each instruction will be the logical OR of the flags set by each of the up to 16 (or 8) individual operations. The status flags are sticky and can be cleared only via a LDMXCSR instruction.

4.1.3.1 Invalid Operation Exception (#I)

The floating-point invalid-operation exception (#I) occurs in response to an invalid arithmetic operand. The flag (IE) and mask (IM) bits for the invalid operation exception are bits 0 and 7, respectively, in the MXCSR register.

Knights Corner instructions forces all floating-point exceptions, including invalid-operation exceptions, to be masked. Thus, for the #I exception the value returned in the destination register is a QNaN, QNaN Indefinite, Integer Indefinite, or one of the source operands. When a value is returned to the destination operand, it over-writes the destination register specified by the instruction. Table [4.1](#page-56-0) lists the invalid-arithmetic operations that the processor detects for instructions and the masked responses to these operations.

Normally, when one or more of the source operands are QNaNs (and neither is an SNaN or in an unsupported format), an invalid-operation exception is not generated. For VCMPPS and VCMPPD when the predicate is one of lt, le, nlt, or nle, a QNaN source operand does generate an invalid-operation exception.

Note that divide-by-zero exceptions (like all other floating-point exceptions) are always masked in Knights Corner.

4.1.3.2 Divide-By-Zero Exception (#Z)

The processor reports a divide-by-zero exception when a VRCP23PS instruction has a 0 operand.

Note that divide-by-zero exceptions (like all other floating-point exceptions) are always masked in Knights Corner.

4.1.3.3 Denormal Operand Exception (#D)

The processor reports a denormal operand exception when an arithmetic instruction attempts to operate on a denormal operand and the DAZ bit in the MXCSR (the "Denormals Are Zero" bit) is not set to 0 (so that denormal operands are not treated as zeros).

Note that denormal exceptions (like all other floating-point exceptions) are always masked in Knights Corner.

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Table 4.1: Masked Responses of Knights Corner instructions to Invalid Arithmetic Operations

4.1.3.4 Numeric Overflow Exception (#0)

The processor reports a numeric overflow exception whenever the rounded result of an arithmetic instruction exceeds the largest allowable finite value that fits in the destination operand.

Note that overflow exceptions (like all other floating-point exceptions) are always masked in Knights Corner.

4.1.3.5 Numeric Underflow Exception (#U)

The processor signals an underflow exception whenever (a) the rounded result of an arithmetic instruction, calculated assuming unbounded exponent, is less than the smallest possible normalized finite value that will fit in the destination operand (the result is *tiny*), and (b) the final rounded result, calculated with bounded exponent determined by the destination format, is inexact.

Note that underflow exceptions (like all other floating-point exceptions) are always masked in Knights Corner.

The flush-to-zero control bit provides an additional option for handling numeric underflow exceptions in Knights Corner. If set (FZ = 1), tiny results (these are usually, but not always, denormal values) are replaced by zeros of the same sign. If not set (FZ=0) then tiny results will be rounded to 0, a denormalized value, or the smallest normalized floating-point number in the destination format, with the sign of the exact result.

4.1.3.6 Inexact Result (Precision) Exception (#P)

The inexact-result exception (also called the precision exception) occurs if the result of an operation is not exactly representable in the destination format. For example, the fraction 1/3 cannot be precisely represented in binary form. This exception occurs frequently and indicates that some (normally acceptable) accuracy has been lost. The exception is supported for applications that need to perform exact arithmetic only. In flush-to-zero mode, the inexact result exception is signaled for any tiny result. (By definition, tiny results are not zero, and are flushed to zero when MXCSR.FZ = 1 for all instructions that support this mode.)

Note that inexact exceptions (like all other floating-point exceptions) are always masked in Knights Corner.

4.2 Denormal Flushing Control

4.2.1 Denormal control in up-conversions and down-conversions

Instruction up-conversions and down-conversions follow specific denormal flushing rules, i.e. for treating input denormals as zeros and for flushing tiny results to zero:

4.2.1.1 Up-conversions

- Up-conversions from float16 to float32 ignore the MXCSR.DAZ setting and this never treat input denormals as zeros. Denormal exceptions are never signaled (the MXCSR.DE flag is never set by these operations).
- Up-conversions from any small floating-point number (namely, float16) to float32 can never generate a float32 output denormal

4.2.1.2 Down-conversions

• Down-conversions from float32 to float16 follow the MXCSR.DAZ setting to decide whether to treat input denormals as zeros or not. For input denormals, the MXCSR.DE flag is set only if MXCSR.DAZ is not set, otherwise it is left unchanged.

- Down-conversions from float32 to any integer format follow the MXCSR.DAZ setting to decide whether to treat input denormals as zeros or not (this may matter only in directed rounding modes). The MXCSR.DE status flag is never set.
- Down-conversions from float32 to any small floating-point number ignore MXCSR.FZ and always preserve output denormals.

4.3 Extended Addressing Displacements

Address displacements used by memory operands to the Knights Corner instructions vector instructions, as well as MVEX-encoded versions of VPREFETCH and CLEVICT, operate differently than do normal x86 displacements. Knights Corner instructions 8-bit displacements (i.e. when MOD.mod=01) are reinterpreted so that they are multiplied by the memory operand's total size in order to generate the final displacement to be used in calculating the effective address (32 bit displacements, which vector instructions may also use, operate normally, in the same way as for normal x86 instructions). Note that extended 8-bit displacements are still signed integer numbers and need to be sign extended.

A given vector instruction's 8-bit displacement is always multiplied by the total number of bytes of memory the instruction accesses, which can mean multiplication by 64, 32, 16, 8, 4, 2 or 1, depending on any broadcast and/or data conversion in effect. Thus when reading a 64-byte (no conversion, no broadcast) source operand, for example via

vmovaps zmm0, [rsi]

the **encoded** 8-bit displacement is ϐirst multiplied by 64 (shifted left by 6) before being used in the effective address calculation. For

vbroadcastss zmm0, [rsi]{uint16} // {1to16} broadcast of {uint16} data

however, the **encoded** displacement would be multiplied by 2. Note that for MVEX versions of VPREFETCH and CLEVICT, we always use disp8*64; for VEX versions we use the standard x86 disp8 displacement.

The use of disp8*N makes it possible to avoid using 32 bit displacements with vector instructions most of the time, thereby reducing code size and shrinking the required size of the paired-instruction decode window by 3 bytes. Disp8*N overcomes disp8 limitations, as it is simply too small to access enough vector operands to be useful (only 4 64-byte operands). Moreover, although disp8*N can only generate displacements that are multiples of N, that's not a significant limitation, since Knights Corner instructions memory operands must already be aligned to the total number of bytes of memory the instruction accesses in order to avoid raising a #GP fault, and that alignment is exactly what disp8*N results in, given aligned base+index addressing.

4.4 Swizzle/up-conversion exceptions

There is a set of Knights Corner instructions that do not accept all regular forms of memory up-conversion/register swizzling and raise a #UD fault for illegal combinations. The instructions are:

• VALIGND

- VCVTDQ2PD
- VCVTPS2PD
- VCVTUDQ2PD
- VEXP223PS
- VFMADD233PS
- VLOG2PS
- VPERMD
- VPERMF32X4
- VPMADD233D
- VPSHUFD
- VRCP23PS
- VRSQRT23PS

Table [4.2](#page-59-0) summarizes which up-conversion/swizzling primitives are allowed for every one of those instructions:

				Register	Memory
Mnemonic	None	$\{1\text{to}16\}$	${4to16}$	swizzles	Conversions
VALIGND	yes	no	no.	no	no
VCVTDQ2PD	yes	yes	yes	yes	no
VCVTPS2PD	yes	yes	yes	yes	no
VCVTUDQ2PD	yes	yes	yes	yes	no
VEX223PS	yes	no	no	no	no
VFMADD233PS	yes	no	yes	no	no
VLOG2PS	yes	no	no	no	no
VPERMD	yes	no	no	no	no
VPERMF32X4	yes	no	no	no	no
VPMADD233D	yes	no	yes	no	no
VPSHUFD	yes	no	no	no	no
VRCP23PS	yes	no	no.	no	no
VRSQRT23PS	yes	no	no	no	no

Table 4.2: Summary of legal and illegal swizzle/conversion primitives for special instructions.

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4.5 Accessing uncacheable memory

When accessing non cacheable memory, it's important to define the amount of data that is really accessed when using Knights Corner instructions (mainly when Knights Corner instructions instructions are used to access to memory mapped I/O regions). Depending on the memory region accessed, an access may cause that a mapped device behave differently.

Knights Corner instructions, when accessing to uncacheable memory access, can be categorized in four different groups:

- regular memory read operations
- vloadunpackh*/vloadunpackl*
- vgatherd*
- memory store operations

4.5.1 Memory read operations

Any Knights Corner instructions that read from memory, apart from vloadunpackh*/vloadunpackl* and vgatherd, access as many consecutive bytes as dictated by the combination of memory SwizzUpCony modifiers.

4.5.2 vloadunpackh*/vloadunpackl*

vloadunpackh*/vloadunpackl* instructions are exceptions to the general rule. Those two instructions will always access 64 bytes of memory. The memory region accessed is between *effective_address & (0x3F)* and *(effective_address & (0x3F)) + 63* in both cases.

4.5.3 vgatherd*

vgatherd instructions are able to gather to up to 16 32 bit elements. The amount of elements accessed is determined by the number of bits set in the vector mask provided as source. *Vgatherd** instruction will access up to 16 different 64-byte memory regions when gathering the elements. Note that, depending on the implementation, only one 64-byte memory access is performed for a variable number of vector elements located in that region.

Each accessed regions will be between *element_effective_address & (0x3F)* and *(element_effective_address & (0x3F)) + 63*.

4.5.4 Memory stores

All Knights Corner instructions that perform memory store operations, update those memory positions determined by the vector mask operand. Vector mask specifies which elements will be actually stored in memory. DownConv^{*} determine the number of bytes per element that will be modified in memory.

4.6 Floating-point Notes

4.6.1 Rounding Modes

VRNDFXPNTPS and conversion instructions with float32 sources, such as VCVTFXPNTPS2DQ, support four selectable rounding modes: round to nearest (even), round toward negative infinity (round down), round toward positive infinity (round up), and round toward zero, These are the standard IEEE rounding modes; see *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 1*, Section 4.8.4, for details.

Knights Corner introduces general support for all four rounding-modes mandated for binary floating-point arithmetic by the IEEE Standard 754-2008.

4.6.1.1 Swizzle-explicit rounding modes

Knights Corner introduces the option of specifying the rounding-mode per instruction via a specific register swizzle mode (by setting the EH bit to 1). This specific rounding-mode takes precedence over whatever MXCSR.RC specifies.

For those instructions (like VRNDFXPNTPS) where an explicit rounding-mode is specified via immediate, this immediate takes precedence over a swizzle-explicit rounding-mode embedded into the encoding of the instruction.

The priority of the rounding-modes of an instruction hence becomes (from highest to lowest):

- 1. Rounding mode specified in the instruction immediate (if any)
- 2. Rounding mode specified is the instruction swizzle attribute
- 3. Rounding mode specified in RC bits of the MXCSR

4.6.1.2 Definition and propagation of NaNs

The IA-32 architecture defines two classes of NaNs: quiet NaNs (QNaNs) and signaling NaNs (SNaNs). Quiet NaNs have 1 as their first fraction bit, SNaNs have 0 as their first fraction bit. An SNaN is quieted by setting its first first fraction bit to 1. The class of a NaN (quiet or signaling) is preserved when converting between different precisions.

The processor never generates an SNaN as a result of a floating-point operation with no SNaN operands, so SNaNs must be present in the input data or have to be inserted by the software.

QNaNs are allowed to propagate through most arithmetic operations without signaling an exception. Note also that Knights Corner instructions do not trap for arithmetic exceptions, as floating-point exceptions are always masked.

If any operation has one or more NaN operands then the result, in most cases, is a QNaN that is one of the input NaNs, quieted if it is an SNaN. This is chosen as the first NaN encountered when scanning the operands from left to right, as presented in the instruction descriptions from Chapter [6.](#page-69-0)

If any floating-point operation with operands that are not NaNs leads to an indefinite result (e.g. $0/0, 0 \times \infty$, or ∞ − ∞), the result will be ONaN Indefinite: 0xFFC00000 for 32 bit operations and 0xFFF80000000000000 for

64 bit operations.

When operating on NaNs, if the instruction does not define any other behavior, Table [4.3](#page-62-0) describes the NaN behavior for unary and binary instructions. Table [4.4](#page-63-0) shows the NaN behavior for ternary fused multiply and add/sub operations. This table can be derived by considering the operation as a concatenation of two binary operations. The first binary operation, the multiply, produces the product. The second operation uses the product as the first operand for the addition.

Table 4.3: Rules for handling NaNs for unary and binary operations.

4.6.1.3 Signed Zeros

Zero can be represented as a +0 or a *−*0 depending on the sign bit. Both encodings are equal in value. The sign of a zero result depends on the operation being performed and the rounding mode being used.

Knights Corner instructions introduces the fused "multiply and add'' and "multiply and sub'' operations. These consist of a multiplication (whose sign is possibly negated) followed by an addition or subtraction, all calculated with just one rounding error.

The sign of the multiplication result is the exclusive-or of the signs of the multiplier and multiplicand, regardless of the rounding mode (a positive number has a sign bit of 0, and a negative one, a sign bit of 1).

The sign of the addition (or subtraction) result is in general that of the exact result. However, when this result is exactly zero, special rules apply: when the sum of two operands with opposite signs (or the difference of two operands with like signs) is exactly zero, the sign of that sum (or difference) is $+0$ in all rounding modes, except round down; in that case, the sign of an exact zero sum (or difference) is *−*0. This is true even if the operands are zeros, or denormals treated as zeros because MXCSR.DAZ is set to 1. Note that $x + x = x - (-x)$ retains the same sign as x even when x is zero; in particular, $(+0) + (+0) = +0$, and $(-0) + (-0) = -0$, in all rounding modes.

When $(a \times b) \pm c$ is exactly zero, the sign of fused multiply-add/subtract shall be determined by the rules above for a sum of operands. When the exact result of $\pm (a \times b) \pm c$ is non-zero yet the final result is zero because of rounding, the zero result takes the sign of the exact result.

The result for "fused multiply and add" follows by applying the following algorithm:

- (x_d, y_d, z_d) =DAZ applied to (Src1, Src2, Src3) (denormal operands, if any, are treated as zeros of the same sign as the operand; other operands are not changed)
- $Result_d = x_d \times y_d + z_d$ computed exactly then rounded to the destination precision.

(intel)

The interpretation of the sources is slightly different for this instruction. Here the Src1 column and NaN2 are associated with Src3[31:0]. Similarly the Src3 column and NaN3 are associated with Src3[63:32].

• Result = FTZ applied to *Result^d* (tiny results are replaced by zeros of the same sign; other results are not changed).

4.6.2 REX prefix and Knights Corner instructions interactions

The REX prefix is illegal in combination with Knights Corner instructions vector instructions, or with mask and scalar instructions allocated using VEX and MVEX prefixes.

Following the *Intel[®]* 64 behavior, if the REX prefix is followed with any legacy prefix and not located just before the opcode escape, it will be ignored.

4.7 Knights Corner instructions State Save

Knights Corner does not include any explicit instruction to perform context save and restore of Knights Corner state. To perform a context save and restore we may use:

- Vector loads and stores for vector registers
- A combination of kmov plus scalar loads and stores for mask registers
- LDMXCSR/STMXCSR for the MXCSR state register

Note also that vector instructions raise a device-not-available (#NM) exceptions when CR0.TS is set. This allows to perform selective lazy save and restore of state.

4.8 Knights Corner instructions Processor State After Reset

Table [4.5](#page-65-0) shows the state of the flags and other registers following power-up for Knights Corner.

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Chapter 5

Instruction Set Reference

Knights Corner instructions that are described in this document follow the general documentation convention established in this chapter.

5.1 Interpreting Instruction Reference Pages

This section describes the format of information contained in the instruction reference pages in this chapter. It explains notational conventions and abbreviations used in these sections

5.1.1 Instruction Format

The following is an example of the format used for each instruction description in this chapter.

5.1.2 Opcode Notations for MVEX Encoded Instructions

In the Instruction Summary Table, the Opcode column presents the details of each instruction byte encoding using notations described in this section. For MVEX encoded instructions, the notations are expressed in the following form (including the modR/M byte if applicable, and the immediate byte if applicable):

MVEX.[NDS,NDD].[512].[66,F2,F3].0F/0F3A/0F38.[W0,W1] opcode [/r] [/ib]

• *MVEX:* indicates the presence of the MVEX prefix is required. The MVEX prefix consists of 4 bytes with the leading byte 62H.

The encoding of various sub-fields of the MVEX prefix is described using the following notations:

- **–** *NDS,NDD:* speciϐies that MVEX.vvvv ϐield is valid for the encoding of a register operand:
	- * MVEX.NDS: MVEX. vvvv encodes the first source register in an instruction syntax where the content of source registers will be preserved. To encode a vector register in the range zmm16 zmm31, the MVEX.vvvv field is pre-pended with MVEX.V'.
	- * MVEX.NDD: MVEX.vvvv encodes the destination register that cannot be encoded by ModR/M:reg field. To encode a vector register in the range zmm16-zmm31, the MVEX.vvvv field is pre-pended with MVEX.V'.
	- * If none of NDS, NDD is present, MVEX.vvvv must be 1111b (i.e. MVEX.vvvv does not encode an operand).
- **–** $66, F2, F3$: The presence or absence of these value maps to the MVEX.pp field encodings. If absent, this corresponds to MVEX.pp=00B. If present, the corresponding MVEX.pp value affects the "opcode" byte in the same way as if a SIMD prefix (66H, F2H or F3H) does to the ensuing opcode byte. Thus a non-zero encoding of MVEX.pp may be considered as an implied 66H/F2H/F3H prefix.
- **–** *0F,0F3A,0F38:* The presence of these values maps to a valid encoding of the MVEX.mmmm ϐield. Only three encoded values of MVEX.mmmm are defined as valid, corresponding to the escape byte sequence of 0FH, 0F3AH and 0F38H.
- **–** *W0:* MVEX.W=0
- **–** *W1:* MVEX.W=1
- **–** The presence of W0/W1 in the opcode column applies to two situations: (a) it is treated as an extended opcode bit, (b) the instruction semantics support an operand size promotion to 64 bit of a general-purpose register operand or a 32 bit memory operand.
- *opcode:* Instruction opcode.
- */r:* Indicates that the ModR/M byte of the instruction contains a register operand and an r/m operand.
- */vsib:* Indicates the memory addressing uses the vector SIB byte.
- *ib:* A 1-byte immediate operand to the instruction that follows the opcode, ModR/M bytes or scale/indexing bytes.

In general, the encoding of the MVEX.R, MVEX.X, MVEX.B, and MVEX.V' fields are not shown explicitly in the opcode column. The encoding scheme of MVEX.R, MVEX.X, MVEX.B, and MVEX.V' fields must follow the rules defined in Chapter 3.

5.1.3 Opcode Notations for VEX Encoded Instructions

In the Instruction Summary Table, the Opcode column presents the details of each instruction byte encoding using notations described in this section. For VEX encoded instructions, the notations are expressed in the following form (including the modR/M byte if applicable, the immediate byte if applicable):

VEX.[NDS,NDD].[66,F2,F3].0F/0F3A/0F38.[W0,W1] opcode [/r] [/ib]

• *VEX:* indicates the presence of the VEX prefix is required. The VEX prefix can be encoded using the three-byte form (the ϐirst byte is C4H), or using the two-byte form (the ϐirst byte is C5H). The two-byte form of VEX only applies to those instructions that do not require the following fields to be encoded: VEX.mmmmm, VEX.W, VEX.X, VEX.B. Refer to Chapter [3](#page-39-0) for more details on the VEX prefix. The encoding of various sub-fields of the VEX prefix is described using the following notations:

- **–** *NDS,NDD:* speciϐies that VEX.vvvv ϐield is valid for the encoding of a register operand:
	- * VEX.NDS: VEX.vvvv encodes the first source register in an instruction syntax where the content of source registers will be preserved.
	- * VEX.NDD: VEX.vvvv encodes the destination register that cannot be encoded by ModR/M:reg ϐield.
	- * If none of NDS, NDD is present, VEX.vvvv must be 1111b (i.e. VEX.vvvv does not encode an operand). The VEX.vvvv field can be encoded using either the 2-byte or 3-byte form of the VEX prefix.
- **–** *66,F2,F3:* The presence or absence of these value maps to the VEX.pp ϐield encodings. If absent, this corresponds to VEX.pp=00B. If present, the corresponding VEX.pp value affects the "opcode" byte in the same way as if a SIMD prefix (66H, F2H or F3H) does to the ensuing opcode byte. Thus a non-zero encoding of VEX.pp may be considered as an implied 66H/F2H/F3H prefix. The VEX.pp field may be encoded using either the 2-byte or 3-byte form of the VEX prefix.
- *0F,0F3A,0F38:* The presence of these values maps to a valid encoding of the VEX.mmmmm field. Only three encoded values of VEX.mmmmm are defined as valid, corresponding to the escape byte sequence of 0FH, 0F3AH and 0F38H. The effect of a valid VEX.mmmmm encoding on the ensuing opcode byte is same as if the corresponding escape byte sequence on the ensuing opcode byte for non-VEX encoded instructions. Thus a valid encoding of VEX.mmmmm may be consider as an implies escape byte sequence of either 0FH, 0F3AH or 0F38H. The VEX.mmmmm field must be encoded using the 3-byte form of VEX prefix.
- **–** *0F,0F3A,0F38 and 2-byte/3-byte VEX:* The presence of 0F3A and 0F38 in the opcode column implies that opcode can only be encoded by the three-byte form of VEX. The presence of 0F in the opcode column does not preclude the opcode to be encoded by the two-byte of VEX if the semantics of the opcode does not require any subfield of VEX not present in the two-byte form of the VEX prefix.
- **–** *W0:* VEX.W=0
- **–** *W1:* VEX.W=1
- **–** The presence of W0/W1 in the opcode column applies to two situations: (a) it is treated as an extended opcode bit, (b) the instruction semantics support an operand size promotion to 64 bit of a general-purpose register operand or a 32 bit memory operand. The presence of W1 in the opcode column implies the opcode must be encoded using the 3-byte form of the VEX prefix. The presence of W0 in the opcode column does not preclude the opcode to be encoded using the C5H form of the VEX prefix, if the semantics of the opcode does not require other VEX subfields not present in the two-byte form of the VEX prefix.
- *opcode:* Instruction opcode.
- */r:* Indicates that the ModR/M byte of the instruction contains a register operand and an r/m operand.
- *ib:* A 1-byte immediate operand to the instruction that follows the opcode, ModR/M bytes or scale/indexing bytes.
- In general, the encoding of the VEX.R, VEX.X, and VEX.B fields are not shown explicitly in the opcode column. The encoding scheme of VEX.R, VEX.X, and VEX.B fields must follow the rules defined in Chapter 3.

Chapter 6

Instruction Descriptions

This Chapter defines all of the Knights Corner instructions vector instructions. Note: Some instruction descriptions refer to the *SSS* or $S_2S_1S_0$, which are bits 6-4 from the MVEX prefix encoding. See Table [2.14](#page-35-0) for more details

CHAPTER 6. INSTRUCTION DESCRIPTIONS

6.1 Vector Mask Instructions

JKNZD - Jump near if mask is not zero

Description

Checks the value of source mask, and if not all mask bits are set to 0, performs a jump to the target instruction specified by the destination operand. If the condition is not satisfied, the jump is not performed and execution continues with the instruction following the instruction.

The target instruction is specified with a relative offset (a signed offset relative to the current value of the instruction pointer in the EIP register). A relative offset (rel8, rel16, or rel32) is generally specified as a label in assembly code, but at the machine code level, it is encoded as a signed, 8-bit or 32 bit immediate value, which is added to the instruction pointer. Instruction coding is most efficient for offsets of -128 to +127. If the operand-size attribute is 16, the upper two bytes of the EIP register are cleared, resulting in a maximum instruction pointer size of 16 bits.

The instruction does not support far jumps (jumps to other code segments). When the target for the conditional jump is in a different segment, use the opposite condition from the condition being tested for the JKNZD instruction, and then access the target with an unconditional far jump (JMP instruction) to the other segment. For example, the following conditional far jump is illegal:

JKNZD FARLABEL;

To accomplish this far jump, use the following two instructions:

JKZD BEYOND; JMP FARLABEL; BEYOND:

This conditional jump is converted to code fetch of one or two cache lines, regardless of jump address or cacheability.

In 64 bit mode, operand size (OSIZE) is fixed at 64 bits. JMP Short is $RIP = RIP + 8$ -bit offset sign extended to 64 bits. JMP Near is RIP = RIP + 32 bit offset sign extended to 64 bits.

Operation

```
if (k1[15:0]!=0){
  tempEIP = EIP + SignExtend(DEST);
  if(OSIZE == 16){
     tempEIP = tempEIP & 0000FFFFH;}
  if (*tempEIP is not within code segment limit*)
   {
     #GP(0);
  }
  else
   {
     EIP = tempEIP
  }
}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

None.

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

CHAPTER 6. INSTRUCTION DESCRIPTIONS

JKZD - Jump near if mask is zero

Description

Checks the value of source mask, and if all mask bits are set to 0, performs a jump to the target instruction specified by the destination operand. If the condition is not satisfied, the jump is not performed and execution continues with the instruction following the instruction.

The target instruction is specified with a relative offset (a signed offset relative to the current value of the instruction pointer in the EIP register). A relative offset (rel8, rel16, or rel32) is generally specified as a label in assembly code, but at the machine code level, it is encoded as a signed, 8-bit or 32 bit immediate value, which is added to the instruction pointer. Instruction coding is most efficient for offsets of -128 to +127. If the operand-size attribute is 16, the upper two bytes of the EIP register are cleared, resulting in a maximum instruction pointer size of 16 bits.

The instruction does not support far jumps (jumps to other code segments). When the target for the conditional jump is in a different segment, use the opposite condition from the condition being tested for the JKNZD instruction, and then access the target with an unconditional far jump (JMP instruction) to the other segment. For example, the following conditional far jump is illegal:

JKZD FARLABEL;

To accomplish this far jump, use the following two instructions:

JKNZD BEYOND; JMP FARLABEL; BEYOND:

This conditional jump is converted to code fetch of one or two cache lines, regardless of jump address or cacheability.

In 64 bit mode, operand size (OSIZE) is fixed at 64 bits. JMP Short is $RIP = RIP + 8$ -bit offset sign extended to 64 bits. JMP Near is RIP = RIP + 32 bit offset sign extended to 64 bits.

Operation

```
if (k1[15:0]==0)
{
  tempEIP = EIP + SignExtend(DEST);
  if(OSIZE == 16){
     tempEIP = tempEIP & 0000FFFFH;}
  if (*tempEIP is not within code segment limit*)
   {
     #GP(0);
  }
  else
   {
     EIP = tempEIP
  }
}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

None.

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#GP(0) If the memory address is in a non-canonical form.
#NM If CRO.TS[bit 3]=1. If $CR0.TS[bit 3]=1$. If preceded by any REX, F0, F2, F3, or 66 prefixes.

KAND - AND Vector Mask

Description

Performs a bitwise AND between the vector masks k2 and the vector mask k1, and writes the result into vector mask k1.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = k1[n] & k2[n]}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kand (__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KANDN - AND NOT Vector Mask

Description

Performs a bitwise AND between vector mask k2, and the NOT (bitwise logical negation) of vector mask k1, and writes the result into vector mask k1.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = (\sim(k1[n])) & k2[n]}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kandn (__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KANDNR - Reverse AND NOT Vector Mask

Description

Performs a bitwise AND between the NOT (bitwise logical negation) of vector mask k2, and the vector mask k1, and writes the result into vector mask k1.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = -(k2[n]) & k1[n]}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kandnr (__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KCONCATH - Pack and Move High Vector Mask

Description

Packs vector masks k1 and k2 and moves the result to the high 32 bits of destination register r64. The rest of the destination register is zeroed.

Operation

 $IMP[15:0] = k2[15:0]$ $IMP[31:16] = k1[15:0]$ $r64[31:0] = 0$ $r64[63:32] = TMP$

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__int64 _mm512_kconcathi_64(__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#NM If CR0.TS[bit 3]=1. If preceded by any REX, F0, F2, F3, or 66 prefixes. If destination is a memory operand.

KCONCATL - Pack and Move Low Vector Mask

Description

Packs vector masks k1 and k2 and moves the result to the low 32 bits of destination register r64. The rest of the destination register is zeroed.

Operation

 $IMP[15:0] = k2[15:0]$ $IMP[31:16] = k1[15:0]$ $r64[31:0] = TMP$ $r64[63:32] = 0$

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__int64 _mm512_kconcatlo_64(__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#NM If CR0.TS[bit 3]=1. If preceded by any REX, F0, F2, F3, or 66 prefixes. If destination is a memory operand.

KEXTRACT - Extract Vector Mask From Register

Description

Extract the 16-bit field selected by imm8[1:0] from general purpose register r64 and write the result into destination mask register k1.

Operation

```
index = imm8[1:0] * 16k1[15:0] = r64[index+15:index]
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kextract_64(__int64, const in);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#NM If CR0.TS[bit 3]=1. If preceded by any REX, F0, F2, F3, or 66 prefixes. If source is a memory operand.

KMERGE2L1H - Swap and Merge High Element Portion and Low Portion of Vector Masks

Description

Move high element from vector mask register k1 into low element of vector mask register k1, and insert low element of k2 into the high portion of vector mask register k1.

Operation

```
tmp = k1[15:8]k1[15:8] = k2[7:0]
k1[7:0] = tmp
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kmerge2l1h (__mmask16, __mmask16 k2);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

CHAPTER 6. INSTRUCTION DESCRIPTIONS

KMERGE2L1L - Move Low Element Portion into High Portion of Vector **Mask**

Description

Insert low element from vector mask register k2 into high element of vector mask register k1. Low element of k1 remains unchanged.

Operation

 $k1[15:8] = k2[7:0]$ *k1[7:0] remains unchanged*

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kmerge2l1l (__mmask16, __mmask16 k2);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KMOV - Move Vector Mask

Description

Either the vector mask register $k2$ or the general purpose register $r32$ is read, and its contents written into destination general purpose register r32 or vector mask register k1; however, general purpose register to general purpose register copies are not supported. When the destination is a general purpose register, the 16 bit value that is copied is zeroextended to the maximum operand size in the current mode.

Operation

```
if(DEST is a general purpose register) {
 DEF [63:16] = 0DEST[15:0] = k2[15:0]} else if(DEST is vector mask and SRC is a general purpose register) {
 k1[15:0] = SRC[15:0]
} else {
 k1[15:0] = k2[15:0]
}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

KNOT - Not Vector Mask

Description

Performs the bitwise NOT of the vector mask k2, and writes the result into vector mask k1.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = -k2[n]}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_knot(__mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KOR - OR Vector Masks

Description

Performs a bitwise OR between the vector mask k2, and the vector mask k1, and writes the result into vector mask k1.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = k1[n] | k2[n]}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kor(__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KORTEST - OR Vector Mask And Set EFLAGS

Description

Performs a bitwise OR between the vector mask register k2, and the vector mask register k1, and sets CF and ZF based on the operation result.

ZF flag is set if both sources are 0x0. CF is set if, after the OR operation is done, the operation result is all 1's.

Operation

```
CF = 1ZF = 1for (n = 0; n < 16; n++) {
  tmp = (k1[n] | k2[n])ZF &= (tmp == 0x0)
  CF &= (tmp == 0x1)}
```
Flags Affected

- The ZF flag is set if the result of OR-ing both sources is all 0s
- The CF flag is set if the result of OR-ing both sources is all 1s
- The OF, SF, AF, and PF flags are set to 0.

Intel® C/C++ Compiler Intrinsic Equivalent

int _mm512_kortestz (__mmask16, __mmask16);

int _mm512_kortestc (__mmask16, __mmask16);

Exceptions

KXNOR - XNOR Vector Masks

Description

Performs a bitwise XNOR between the vector mask k1 and the vector mask k2, and the result is written into vector mask k1.

The primary purpose of this instruction is to provide a way to set a vector mask register to 0xFFFF in a single clock; this is accomplished by selecting the source and destination to be the same mask register. In this case the result will be 0xFFFF regardless of the original contents of the register.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = -(k1[n] \hat{ } k2[n])}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kxnor (__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

KXOR - XOR Vector Masks

Description

Performs a bitwise XOR between the vector mask k2, and the vector mask k1, and writes the result into vector mask k1.

Operation

```
for (n = 0; n < 16; n++) {
 k1[n] = k1[n] k2[n]}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_kxor (__mmask16, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

CHAPTER 6. INSTRUCTION DESCRIPTIONS

6.2 Vector Instructions

VADDNPD - Add and Negate Float64 Vectors

Description

Performs an element-by-element addition between float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3, then negates the result. The final result is written into float64 vector zmm1.

Note that all the operations must be performed before rounding.

Table 6.1: VADDN outcome when adding zeros depending on rounding-mode. See Signed Zeros in Section [4.6.1.3](#page-62-0) for other cases with a result of zero.

> This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
 // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
  RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
 if(k1[n] != 0) {i = 64*n// float64 operation
   zmm1[i+63:i] = (-zmm2[i+63:i]) + (-tmpSrc3[i+63:i])
 }
```


}

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Intel® C/C++ Compiler Intrinsic Equivalent

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VADDNPS - Add and Negate Float32 Vectors

Description

Performs an element-by-element addition between float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3, then negates the result. The final result is written into float32 vector zmm1.

Note that all the operations must be performed before rounding.

Table 6.2: VADDN outcome when adding zeros depending on rounding-mode. See Signed Zeros in Section [4.6.1.3](#page-62-0) for other cases with a result of zero.

> This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
 // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
 if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = (-zmm2[i+31:i]) + (-tmpSrc3[i+31:i])
 }
```


}

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Intel® C/C++ Compiler Intrinsic Equivalent

__m512 _mm512_mask_addn_ps (__m512, __mmask16, __m512, __m512);

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{*f*32}

MVEX.EH=0

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VADDPD - Add Float64 Vectors

Description

Performs an element-by-element addition between float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3. The result is written into float64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm2[i+63:i] + tmpSrc3[i+63:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Intel® C/C++ Compiler Intrinsic Equivalent

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{f64}

MVEX.EH=0

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VADDPS - Add Float32 Vectors

Description

Performs an element-by-element addition between float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] + tmpSrc3[i+31:i]
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Intel® C/C++ Compiler Intrinsic Equivalent

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VADDSETSPS - Add Float32 Vectors and Set Mask to Sign

Description

Performs an element-by-element addition between float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

In addition, the sign of the result for the n-th element is written into the n-th bit of vector mask k1.

It is the sign bit of the final result that gets copied to the destination, as opposed to the result of comparison with zero.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k1. Elements in zmm1 and k1 with the corresponding bit clear in k1 register retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
 // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
 if(k1[n] != 0) {
    i = 32*n
    // float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] + tmpSrc3[i+31:i]
    k1[n] = zmm1[i+31]}
```


}

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Intel® C/C++ Compiler Intrinsic Equivalent

Memory Up-conversion: S_{f32}

Register Swizzle: S_{*f*32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If a memory address referencing the SS segment is

VALIGND - Align Doubleword Vectors

Description

Concatenates and shifts right doubleword elements from vector zmm2 and memory/vector zmm3. The result is written into vector zmm1.

No swizzle, broadcast, or conversion is performed by this instruction.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
src[511:0] = zmm3/mt
// Concatenate sources
tmp[511:0] = src[511:0]tmp[1023:512] = zmm2[511:0]// Shift right doubleword elements
SHIFT = imm8[3:0]tmp[1023:0] = tmp[1023:0] >> (32*SHIFT)// Apply write-mask
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*nzmm1[i+31:i] = tmp[i+31:i]
  }
}
```


Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VBLENDMPD - Blend Float64 Vectors using the Instruction Mask

Description

Performs an element-by-element blending between float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3, using the instruction mask as selector. The result is written into float64 vector zmm1.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector: every element of the destination is conditionally selected between first source or second source using the value of the related mask bit (0 for first source, 1 for second source).

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = tmpSrc3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(tmpSrc3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
   if(k1[n]=1 or *no write-mask*) {
      zmm1[i+63:i] = tmpSrc3[i+63:i]
   } else {
      zmm1[i+63:i] = zmm2[i+63:i]
   }
}
```
SIMD Floating-Point Exceptions

None.

Denormal Handling

Treat Input Denormals As Zeros : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m5128 _mm512_mask_blend_pd (__mmask8, __m5128, __m5128);

Exceptions

(intel)

Real-Address Mode and Virtual-8086

VBLENDMPS - Blend Float32 Vectors using the Instruction Mask

Description

Performs an element-by-element blending between float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3, using the instruction mask as selector. The result is written into float32 vector zmm1.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector: every element of the destination is conditionally selected between first source or second source using the value of the related mask bit (0 for first source, 1 for second source).

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = tmpSrc3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(tmpSrc3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
   if(k1[n]=1 or *no write-mask*) {
      zmm1[i+31:i] = tmpSrc3[i+31:i]
   } else {
      zmm1[i+31:i] = zmm2[i+31:i]}
}
```
SIMD Floating-Point Exceptions

Invalid.

Denormal Handling

Treat Input Denormals As Zeros : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512 _mm512_mask_blend_ps (__mmask16, __m512, __m512);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

(intel)

VBROADCASTF32X4 - Broadcast 4xFloat32 Vector

Description

The 4, 8 or 16 bytes (depending on the conversion and broadcast in effect) at memory address m_t are broadcast and/or converted to a float 32 vector. The result is written into float32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {4to16}
tmpSrc2[127:0] = UpConvLoad<sub>f32</sub>(m<sub>t</sub>)for (n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    i = 32*nj = i & 0x7Fzmm1[i+31:i] = tmpSrc2[j+31:j])
  }
}
```
Flags Affected

Invalid.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

VBROADCASTF64X4 - Broadcast 4xFloat64 Vector

Description

The 32 bytes at memory address m_t are broadcast to a float64 vector. The result is written into float64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {4to8}
tmpSrc2[255:0] = UpConvLoad<sub>f64</sub>(m<sub>t</sub>)for (n = 0; n < 8; n++) {
  if (k1[n] != 0) {
    i = 64*nj = i & 0xFFzmm1[i+63:i] = tmpSrc2[j+63:j])
  }
}
```
Flags Affected

None.

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

VBROADCASTI32X4 - Broadcast 4xInt32 Vector

Description

The 4, 8 or 16 bytes (depending on the conversion and broadcast in effect) at memory address *m^t* are broadcast and/or converted to a int32 vector. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {4to16}
tmpSrc2[127:0] = UpConvLoad<sub>i32</sub>(m<sub>t</sub>)for (n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    i = 32*nj = i & 0x7Fzmm1[i+31:i] = tmpSrc2[j+31:j])
  }
}
```
Flags Affected

None.

Memory Up-conversion: U*i*³²

Intel® C/C++ Compiler Intrinsic Equivalent

VBROADCASTI64X4 - Broadcast 4xInt64 Vector

Description

The 32 bytes at memory address *m^t* are broadcast to a int64 vector. The result is written into int64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {4to8}
tmpSrc2[255:0] = UpConvLoad<sub>i64</sub>(m<sub>t</sub>)for (n = 0; n < 8; n++) {
  if (k1[n] != 0) {
    i = 64*nj = i & 0xFFzmm1[i+63:i] = tmpSrc2[j+63:j])
  }
}
```
Flags Affected

None.

Memory Up-conversion: U*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

VBROADCASTSD - Broadcast Float64 Vector

Description

The 8 bytes at memory address m_t are broadcast to a float64 vector. The result is written into float64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {1to8}
tmpSrcc2[63:0] = UpConvLoad_{f64}(m_t)for (n = 0; n < 8; n++) {
  if (k1[n] != 0) {
    i = 64*nzmm1[i+63:i] = tmpSrc2[63:0]
 }
}
```
Flags Affected

None.

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

VBROADCASTSS - Broadcast Float32 Vector

Description

The 1, 2, or 4 bytes (depending on the conversion and broadcast in effect) at memory address m_t are broadcast and/or converted to a float 32 vector. The result is written into float32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {1to16}
tmpSrc2[31:0] = UpConvLoad<sub>f32</sub>(m<sub>t</sub>)for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*nzmm1[i+31:i] = tmpSrc2[31:0]
  }
}
```
Flags Affected

Invalid.

Memory Up-conversion: U_{*f*32}

Intel® C/C++ Compiler Intrinsic Equivalent

VCMPPD - Compare Float64 Vectors and Set Vector Mask

Description

Performs an element-by-element comparison between float64 vector zmm1 and the float64 vector result of the swizzle/broadcast/conversion from memory or float64 vector zmm2. The result is written into vector mask k2.

Note: If DAZ=1, denormals are treated as zeros in the comparison (original source registers untouched). untouched). +0 equals *−*0. Comparison with NaN returns false.

Infinity of like signs, are considered equals. Infinity values of either signs are considered ordered values.

Table [6.3](#page-142-0) summarizes VCMPPD behavior, in particular showing how various NaN results can be produced.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a

write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

Immediate Format

Operation

```
switch (IMM8[2:0]) {
   case 0: OP ← EQ; break;
   case 1: OP ← LT; break;
   case 2: OP ← LE; break;
   case 3: OP ← UNORD; break;
   case 4: OP ← NEQ; break;
   case 5: OP ← NLT; break;
   case 6: OP ← NLE; break;
   case 7: OP ← ORD; break;
}
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoadf64(zmm2/mt)
}
for (n = 0; n < 8; n++) {
 k2[n] = 0if(k1[n] != 0) {i = 64*n// float64 operation
   k2[n] = (zmm1[i+63:i] OP tmpSrc2[i+63:i]) ? 1 : 0
  }
}
k2[15:8] = 0
```


Instruction Pseudo-ops

Compilers and assemblers may implement the following pseudo-ops in addition to the standard instruction op:

SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VCMPPS - Compare Float32 Vectors and Set Vector Mask

Description

Performs an element-by-element comparison between float32 vector zmm1 and the float32 vector result of the swizzle/broadcast/conversion from memory or float32 vector zmm2. The result is written into vector mask k2.

Note: If DAZ=1, denormals are treated as zeros in the comparison (original source registers untouched). untouched). +0 equals *−*0. Comparison with NaN returns false.

Infinity of like signs, are considered equals. Infinity values of either signs are considered ordered values.

Table [6.4](#page-147-0) summarizes VCMPPS behavior, in particular showing how various NaN results can be produced.

Table 6.4: VCMPPS behavior

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

Immediate Format

Operation

```
switch (IMM8[2:0]) {
   case 0: OP ← EQ; break;
   case 1: OP ← LT; break;
   case 2: OP ← LE; break;
   case 3: OP ← UNORD; break;
   case 4: OP ← NEQ; break;
   case 5: OP ← NLT; break;
   case 6: OP ← NLE; break;
   case 7: OP ← ORD; break;
}
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoadf32(zmm2/mt)
}
for (n = 0; n < 16; n++) {
  k2[n] = 0if(k1[n] != 0) {i = 32*n
   // float32 operation
   k2[n] = (zmm1[i+31:i] OP tmpSrc2[i+31:i]) ? 1 : 0
 }
}
```


Instruction Pseudo-ops

Compilers and assemblers may implement the following pseudo-ops in addition to the standard instruction op:

SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f32}

Register Swizzle: S_{*f*32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VCVTDQ2PD - Convert Int32 Vector to Float64 Vector

Description

Performs an element-by-element conversion from the int32 vector result of the swizzle/broadcast/conversion from memory or int32 vector zmm2 to a float64 vector. The result is written into float64 vector zmm1. The int32 source is read from either the lower half of the source operand (int32 vector zmm2), full memory source (8 elements, i.e. 256 bits) or the broadcast memory source.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[255:0] = zmm2[255:0]} else {
   tmpSrc2[255:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nj = 32*nzmm1[i+63:i] =
        CvtInt32ToFloat64(tmpSrc2[j+31:j])
  }
}
```
SIMD Floating-Point Exceptions

None.

Denormal Handling

Treat Input Denormals As Zeros : Not Applicable

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

_m512d _mm512_cvtepi32lo_pd (__m512i); _m512d _mm512_mask_cvtepi32lo_pd (__m512d, __mmask8, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

VCVTFXPNTDQ2PS - Convert Fixed Point Int32 Vector to Float32 Vector

Description

Performs an element-by-element conversion from the int32 vector result of the swizzle/broadcast/conversion from memory or int32 vector zmm2 to a float32 vector, then performs an optional adjustment to the exponent.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
expadj = IMMS[6:4]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc2[511:0] = zmm2[511:0]} else {
  RoundingMode = MXCSR.RC
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
```


```
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] =
    CvtInt32ToFloat32(tmpSrc2[i+31:i], RoundingMode) / EXPADJ_TABLE[expadj]
 }
}
```
SIMD Floating-Point Exceptions

Precision.

Denormal Handling

Treat Input Denormals As Zeros : Not Applicable

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512 _mm512_cvtfxpnt_round_adjustepi32_ps(__m512i, int, _MM_EXP_ADJ_ENUM);
- __m512 _mm512_mask_cvtfxpnt_round_adjustepi32_ps(__m512, __mmask16, __m512i, int, _MM_EXP_ADJ_ENUM);

Exceptions

#UD Instruction not available in these modes

Protected and Compatibility Mode

- #UD Instruction not available in these modes
- 64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes.

VCVTFXPNTPD2DQ - Convert Float64 Vector to Fixed Point Int32 Vector

Description

Performs an element-by-element conversion and rounding from the float64 vector result of the swizzle/broadcast/conversion from memory or float64 vector zmm2 to a int32 vector . The int32 result is written into the lower half of the destination register zmm1; the other half of the destination is set to zero.

Out-of-range values are converted to the nearest representable value and that NaNs convert to 0, because this makes the calculation of Exp2 more efficient (avoiding problems with converting very large values to integers, where undetected incorrect values could otherwise result from overflow). Table [6.5](#page-159-0) describes what should be the result when dealing with floating-point special number.

Table 6.5: Converting to integer special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
RoundingMode = IMMS[1:0]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*ni = 32*nzmm1[j+31:j] =CvtFloat64ToInt32(tmpSrc2[i+63:i], RoundingMode)
 }
}
zmm1[511:256] = 0
```
SIMD Floating-Point Exceptions

Invalid, Precision.

Denormal Handling

```
Treat Input Denormals As Zeros :
         (MXCSR.DAZ)? YES : NO
Flush Tiny Results To Zero :
```
Not Applicable

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_cvtfxpnt_roundpd_epi32lo(__m512i, __mmask8, __m512d, int);

Exceptions

#PF(fault-code) For a page fault.

VCVTFXPNTPD2UDQ - Convert Float64 Vector to Fixed Point Uint32 Vector

Description

Performs an element-by-element conversion and rounding from the float64 vector result of the swizzle/broadcast/conversion from memory or float64 vector zmm2 to a uint32 vector . The uint32 result is written into the lower half of the destination register zmm1; the other half of the destination is set to zero.

Out-of-range values are converted to the nearest representable value and that NaNs convert to 0, because this makes the calculation of Exp2 more efficient (avoiding problems with converting very large values to integers, where undetected incorrect values could otherwise result from overflow). Table [6.6](#page-163-0) describes what should be the result when dealing with floating-point special number.

Table 6.6: Converting to integer special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
RoundingMode = IMMS[1:0]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*ni = 32*nzmm1[j+31:j] =CvtFloat64ToUint32(tmpSrc2[i+63:i], RoundingMode)
 }
}
zmm1[511:256] = 0
```
SIMD Floating-Point Exceptions

Invalid, Precision.

Denormal Handling

```
Treat Input Denormals As Zeros :
         (MXCSR.DAZ)? YES : NO
Flush Tiny Results To Zero :
```
Not Applicable

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_cvtfxpnt_roundpd_epi32lo(__m512i, __mmask8, __m512d, int);

Exceptions

#PF(fault-code) For a page fault.

VCVTFXPNTPS2DQ - Convert Float32 Vector to Fixed Point Int32 Vector

Description

Performs an element-by-element conversion and rounding from the float 32 vector result of the swizzle/broadcast/conversion from memory or float32 vector zmm2 to a int32 vector , with an optional exponent adjustment before the conversion.

Out-of-range values are converted to the nearest representable value and that NaNs convert to 0, because this makes the calculation of Exp2 more efficient (avoiding problems with converting very large values to integers, where undetected incorrect values could otherwise result from overflow). Table [6.7](#page-167-0) describes what should be the result when dealing with floating-point special number.

Table 6.7: Converting to integer special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
RoundingMode = IMM8[1:0]
expadj = IMMS[6:4]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] =CvtFloat32ToInt32(tmpSrc2[i+31:i] * EXPADJ_TABLE[expadj], RoundingMode)
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Precision.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_cvtfxpnt_round_adjustps_epi32(__m512i, __mmask16, __m512, int, _MM_EXP_ADJ_ENUM);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VCVTFXPNTPS2UDQ - Convert Float32 Vector to Fixed Point Uint32 Vector

Description

Performs an element-by-element conversion and rounding from the float 32 vector result of the swizzle/broadcast/conversion from memory or float32 vector zmm2 to a uint32 vector , with an optional exponent adjustment before the conversion.

Out-of-range values are converted to the nearest representable value and that NaNs convert to 0, because this makes the calculation of Exp2 more efficient (avoiding problems with converting very large values to integers, where undetected incorrect values could otherwise result from overflow). Table [6.8](#page-171-0) describes what should be the result when dealing with floating-point special number.

Table 6.8: Converting to integer special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
RoundingMode = IMM8[1:0]
expadj = IMMS[6:4]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] =CvtFloat32ToUint32(tmpSrc2[i+31:i] * EXPADJ_TABLE[expadj], RoundingMode)
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Precision.

Denormal Handling

```
Treat Input Denormals As Zeros :
     (MXCSR.DAZ)? YES : NO
```
Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_cvtfxpnt_round_adjustps_epi32(__m512i, __mmask16, __m512, int, _MM_EXP_ADJ_ENUM);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VCVTFXPNTUDQ2PS - Convert Fixed Point Uint32 Vector to Float32 Vector

Description

Performs an element-by-element conversion from the uint32 vector result of the swizzle/broadcast/conversion from memory or uint32 vector zmm2 to a float32 vector, then performs an optional adjustment to the exponent.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
expadj = IMMS[6:4]if(source is a register operand and MVEX.EH bit is 1) {
  if(SSS[2]==1) Supress_Exception_Flags() // SAE
 // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc2[511:0] = zmm2[511:0]} else {
  RoundingMode = MXCSR.RC
```


```
tmpSrc2[511:0] = SwizzUpConvLoadi32(zmm2/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] =
    CvtUint32ToFloat32(tmpSrc2[i+31:i], RoundingMode) / EXPADJ_TABLE[expadj]
  }
}
```
SIMD Floating-Point Exceptions

Precision.

Denormal Handling

Treat Input Denormals As Zeros : Not Applicable

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VCVTPD2PS - Convert Float64 Vector to Float32 Vector

Description

Performs an element-by-element conversion and rounding from the float64 vector result of the swizzle/broadcast/conversion from memory or float64 vector zmm2 to a float32 vector . The result is written into float32 vector zmm1. The float32 result is written into the lower half of the destination register zmm1; the other half of the destination is set to zero.

Table 6.9: Converting float64 to float32 special values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc2[511:0] = zmm2[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc2[511:0] = SwizzUpConvLoad_{f64}(zmm2/m_t)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nj = 32*nzmm1[j+31:j] =
```


```
CvtFloat64ToFloat32(tmpSrc2[i+63:i], RoundingMode)
 }
}
zmm1[511:256] = 0
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

CHAPTER 6. INSTRUCTION DESCRIPTIONS

VCVTPS2PD - Convert Float32 Vector to Float64 Vector

Description

Performs an element-by-element conversion and rounding from the float32 vector result of the swizzle/broadcast/conversion from memory or float32 vector zmm2 to a float64 vector . The result is written into float64 vector zmm1. The float32 source is read from either the lower half of the source operand (float 32 vector zmm2), full memory source (8 elements, i.e. 256-bits) or the broadcast memory source.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress Exception Flags() // SAE
   tmpSrc2[255:0] = zmm2[255:0]
} else {
   tmpSrc2[255:0] = SwizzUpConvLoad<sub>f32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nj = 32*nzmm1[i+63:i] =
        CvtFloat32ToFloat64(tmpSrc2[j+31:j])
  }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f32}

Register Swizzle: S_{*f*32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

CHAPTER 6. INSTRUCTION DESCRIPTIONS

VCVTUDQ2PD - Convert Uint32 Vector to Float64 Vector

Description

Performs an element-by-element conversion from the uint32 vector result of the swizzle/broadcast/conversion from memory or uint32 vector zmm2 to a float64 vector. The result is written into float64 vector zmm1. The uint32 source is read from either the lower half of the source operand (uint32 vector zmm2), full memory source (8 elements, i.e. 256-bits) or the broadcast memory source.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[255:0] = zmm2[255:0]} else {
   tmpSrc2[255:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nj = 32*nzmm1[i+63:i] =
        CvtUint32ToFloat64(tmpSrc2[j+31:j])
  }
}
```
SIMD Floating-Point Exceptions

None.

Denormal Handling

Treat Input Denormals As Zeros : Not Applicable

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

_m512d _mm512_cvtepu32lo_pd (__m512i); _m512d _mm512_mask_cvtepu32lo_pd (__m512d, __mmask8, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

VEXP223PS - Base-2 Exponential Calculation of Float32 Vector

Description

Computes the element-by-element base-2 exponential computation of the int32 vector on memory or int32 vector zmm2 with 0.99ULP (relative error). Input int32 values are considered as fixed point numbers with a fraction offset of 24 bits (i.e. 8 MSBs correspond to sign and integer part; 24 LSBs correspond to fractional part). The result is written into float32 vector zmm1.

exp2 of a FP input value is computed as a two-instruction sequence:

- 1. vcvtfxpntps2dq (with exponent adjustment, so that destination format is 32b, with 8b for integer part and 24b for fractional part)
- 2. vexp223ps

All overflows are captured by the combination of the saturating behavior of vcytfxpntps2dq instruction and the detection of MAX_INT/MIN_INT by the vexp223ps instruction. Tiny input numbers are quietly flushed to the fixed-point value 0 by the vcvtfxpntps2dq instruction, which produces an overall output $exp2(0) = 1.0f$.

The overall behavior of the two-instruction sequence is the following:

- *−∞* returns +0*.*0*f*
- *±*0*.*0*f* returns 1*.*0*f* (exact result)
- +*∞* returns +*∞* (#Overϐlow)
- NaN returns 1*.*0*f* (#Invalid)
- n, where n is an integral value returns 2 *ⁿ* (exact result)

Table 6.10: vexp2_1ulp() special int values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
tmpSrc2[511:0] = zmm2/mt
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
}
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
   i = 32*nzmm1[i+31:i] = exp2_1ulp(tmpSrc2[i+31:i])
 }
}
```
SIMD Floating-Point Exceptions

Overflow.

Denormal Handling

Treat Input Denormals As Zeros : Not Applicable

Flush Tiny Results To Zero : YES

Register Swizzle

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

```
__m512 _mm512_exp223_ps (__m512i);
__m512 _mm512_mask_exp223_ps (__m512, __mmask16, __m512i);
```
Exceptions

VFIXUPNANPD - Fix Up Special Float64 Vector Numbers With NaN Passthrough

Description

Performs an element-by-element fix-up of various real and special number types in the float64 vector zmm2 using the 21-bit table values from the result of the swizzle/broadcast/conversion process on memory or int64 vector zmm3. The result is merged into float64 vector zmm1. Unlike in *vfixuppd*, source NaN values are passedthrough as quietized values. Note that, also unlike in *vϔixup*, this quietization translates into a #IE exception flag being reported for input SNaNs.

This instruction is specifically intended for use in fixing up the results of arithmetic calculations involving one source, although it is generally useful for fixing up the results of multiple-instruction sequences to reflect special-number inputs. For example, consider rcp(0). Input 0 to rcp, and you should get inf. However, evaluating rcp via $2x - ax^2$ (Newton-Raphson), where $x = approx(1/0) = \infty$, incorrectly yields NaN. To deal with this, vfixupps can be used after the N-R reciprocal sequence to set the result to ∞ when the input is 0.

Denormal inputs must be treated as zeros of the same sign if DAZ is enabled.

Note that NO_CHANGE_TOKEN leaves the destination (output) unchanged. This means that if the destination is a denormal, its value is not flushed to 0.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
enum TOKEN_TYPE
{
   NO CHANGE TOKEN = 0,
```
 NEG _{INF} $TOKEN$ = 1,

```
NEG_ZERO_TOKEN = 2,
   POS ZERO TOKEN = 3,
   POS INF TOKEN = 4,
   NAN TOKEN = 5,MAX_DOUBLE_TOKEN = 6,
   MIN_DOUBLE_TOKEN = 7,
}
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpzmm3[511:0] = zmm3[511:0]
} else {
   tempzmm3[511:0] = SwizzUpConvLoad<sub>i64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
 if (k1[n] := 0) {
   i = 64*ntsrc[63:0] = zmm2[i+63:i]if (IsNaN(tsrc[63:0])
   {
      zmm1[i+63:i] = QNaN(zmm2[i+63:i])
   }
   else
   {
      // tmp is an int value
      if (tsrc[63:0] == -inf) tmp = 0else if (tsrc[63:0] < 0) tmp = 1
      else if (tsrc[63:0] == -0) tmp = 2
      else if (tsrc[63:0] == +0) tmp = 3
      else if (tsrc[63:0] == inf) tmp = 5
      else /* tsrc[63:0] > 0 */ tmp = 4table[20:0] = tmpzmm3[i+63:i]token = table[(tmp*3)+2: tmp*3] // table is viewed as one 21-bit
                                       // little-endian value.
                                       // token is an int value
                                       // the 7th entry is unused
      // float64 result
      if (token == NEG_INF_TOKEN) zmm1[i+63:i] = -infelse if (token == NEG_ZERO_TOKEN) zmm1[i+63:i] = -0else if (token == POS_ZERO_TOKEN) zmm1[i+63:i] = +0else if (token == POS_INF_TOKEN) zmm1[i+63:i] = +infelse if (token == NAN_TOKEN) zmm1[i+63:i] = QNaN_indefinite
      else if (token == MAX_DOUBLE_TOKEN) zmm1[i+63:i] = NMAX
      else if (token == MIN DOUBLE TOKEN) zmm1[i+63:i] = -NMAX
    else if (token == NO_CHANGE_TOKEN) { /* zmm1[i+63:i] remains unchanged */ }
   }
 }
```


}

SIMD Floating-Point Exceptions

Invalid.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

If preceded by any REX, F0, F2, F3, or 66 prefixes.

VFIXUPNANPS - Fix Up Special Float32 Vector Numbers With NaN Passthrough

Description

Performs an element-by-element fix-up of various real and special number types in the float32 vector zmm2 using the 21-bit table values from the result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is merged into float32 vector zmm1. Unlike in *vfixupps*, source NaN values are passedthrough as quietized values. Note that, also unlike in *vϔixup*, this quietization translates into a #IE exception flag being reported for input SNaNs.

This instruction is specifically intended for use in fixing up the results of arithmetic calculations involving one source, although it is generally useful for fixing up the results of multiple-instruction sequences to reflect special-number inputs. For example, consider rcp(0). Input 0 to rcp, and you should get inf. However, evaluating rcp via $2x - ax^2$ (Newton-Raphson), where $x = approx(1/0) = \infty$, incorrectly yields NaN. To deal with this, vfixupps can be used after the N-R reciprocal sequence to set the result to ∞ when the input is 0.

Denormal inputs must be treated as zeros of the same sign if DAZ is enabled.

Note that NO_CHANGE_TOKEN leaves the destination (output) unchanged. This means that if the destination is a denormal, its value is not flushed to 0.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
enum TOKEN_TYPE
{
   NO CHANGE TOKEN = 0,
```


 NEG _{INF} $TOKEN$ = 1,

```
NEG_ZERO_TOKEN = 2,
   POS ZERO TOKEN = 3,
   POS INF TOKEN = 4,
   NAN_TOKEN = 5,
   MAX_FLOAT_TOKEN = 6,
   MIN_FLOAT_TOKEN = 7,
}
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpzmm3[511:0] = zmm3[511:0]
} else {
   tempzmm3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
 if (k1[n] := 0) {
   i = 32*ntsrc[31:0] = zmm2[i+31:i]if (IsNaN(tsrc[31:0])
   {
      zmm1[i+31:i] = QNaN(zmm2[i+31:i])
   }
   else
   {
      // tmp is an int value
      if (tsrc[31:0] == -inf) tmp = 0else if (tsrc[31:0] < 0) tmp = 1
      else if (tsrc[31:0] == -0) tmp = 2
      else if (tsrc[31:0] == +0) tmp = 3
      else if (tsrc[31:0] == inf) tmp = 5
      else /* tsrc[31:0] > 0 */ tmp = 4table[20:0] = tmpzmm3[i+31:i]token = table[(tmp*3)+2: tmp*3] // table is viewed as one 21-bit
                                       // little-endian value.
                                       // token is an int value
                                       // the 7th entry is unused
      // float32 result
      if (token == NEG_INF_TOKEN) zmm1[i+31:i] = -infelse if (token == NEG_ZERO_TOKEN) zmm1[i+31:i] = -0else if (token == POS_ZERO_TOKEN) zmm1[i+31:i] = +0else if (token == POS_INF_TOKEN) zmm1[i+31:i] = +infelse if (token == NAN_TOKEN) zmm1[i+31:i] = QNaN_indefinite
      else if (token == MAX_FLOAT_TOKEN) zmm1[i+31:i] = NMAX
      else if (token == MIN FLOAT TOKEN) zmm1[i+31:i] = -NMAX
    else if (token == NO_CHANGE_TOKEN) { /* zmm1[i+31:i] remains unchanged */ }
   }
 }
```


}

SIMD Floating-Point Exceptions

Invalid.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512 _mm512_mask_ϐixupnan_ps (__m512, __mmask16, __m512, __m512i);

Exceptions

VFMADD132PD - Multiply Destination By Second Source and Add To First Source Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm1 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, then adds the result to float64 vector zmm2. The final sum is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {
    i = 64*n// float64 operation
    zmm1[i+63:i] = zmm1[i+63:i] * tmpSrc3[i+63:i] + zmm2[i+63:i]
 }
}
```


SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFMADD132PS - Multiply Destination By Second Source and Add To First Source Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm1 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then adds the result to float32 vector zmm2. The final sum is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] |= 0) {
    i = 32*n// float32 operation
    zmm1[i+31:i] = zmm1[i+31:i] * tmpSrc3[i+31:i] + zmm2[i+31:i]
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

VFMADD213PD - Multiply First Source By Destination and Add Second Source Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm2 and float64 vector zmm1 and then adds the result to the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3. The final sum is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm2[i+63:i] * zmm1[i+63:i] + tmpSrc3[i+63:i]
```


} }

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

CHAPTER 6. INSTRUCTION DESCRIPTIONS

CHAPTER 6. INSTRUCTION DESCRIPTIONS

VFMADD213PS - Multiply First Source By Destination and Add Second Source Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm2 and float32 vector zmm1 and then adds the result to the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float 32 zmm3. The final sum is written into float 32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n
    // float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] * zmm1[i+31:i] + tmpSrc3[i+31:i]
```


} }

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

CHAPTER 6. INSTRUCTION DESCRIPTIONS

VFMADD231PD - Multiply First Source By Second Source and Add To Destination Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, then adds the result to float64 vector zmm1. The final sum is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
 // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
  RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
 if(k1[n] != 0) {i = 64*n// float64 operation
   zmm1[i+63:i] = zmm2[i+63:i] * tmpSrc3[i+63:i] + zmm1[i+63:i]
 }
```


}

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If a memory address referencing the SS segment is

VFMADD231PS - Multiply First Source By Second Source and Add To Destination Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then adds the result to float32 vector zmm1. The final sum is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
 // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
 if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] * tmpSrc3[i+31:i] + zmm1[i+31:i]
 }
```
}

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

VFMADD233PS - Multiply First Source By Specially Swizzled Second Source and Add To Second Source Float32 Vectors

Description

This instruction is built around the concept of 4-element sets, of which there are four: elements 0-3, 4-7, 8-11, and 12-15. If we refer to the float32 vector result of the broadcast (no conversion is supported) process on memory or the ϐloat32 vector zmm3 (no swizzle is supported) as t3, then:

Each element 0-3 of float32 vector zmm2 is multiplied by element 1 of t3, the result is added to element 0 of t3, and the final sum is written into the corresponding element $0-3$ of float32 vector zmm1.

Each element 4-7 of float32 vector zmm2 is multiplied by element 5 of t3, the result is added to element 4 of t3, and the final sum is written into the corresponding element 4-7 of float32 vector zmm1.

Each element 8-11 of float 32 vector zmm2 is multiplied by element 9 of t3, the result is added to element 8 of t3, and the final sum is written into the corresponding element $8-11$ of float32 vector zmm1.

Each element 12-15 of float32 vector zmm2 is multiplied by element 13 of t3, the result is added to element 12 of t3, and the final sum is written into the corresponding element 12-15 of float 32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded.

This instruction makes it possible to perform scale and bias in a single instruction without needing to have either scale or bias already loaded in a register. This saves one vector load for each interpolant, representing around ten percent of shader instructions.

For structure-of-arrays (SOA) operation, this instruction is intended to be used with the *{*4*to*16*}* broadcast on src2, allowing all 16 scale and biases to be identical. For array-of-

structures (AOS) vec4 operations, no broadcast is used, allowing four different scales and biases, one for each vec4.

No conversion or swizzling is supported for this instruction. However, all broadcasts except *{*1*to*16*}* are supported (i.e. 16to16 and 4to16).

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*nbase = (n & 0x03) * 32scale[31:0] = tmpSrc3[base+63:base+32]bias[31:0] = tmpSrc3[base+31:base]// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] * scale[31:0] + bias[31:0]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512 _mm512_fmadd233_ps (__m512, __m512);
- __m512 _mm512_mask_fmadd233_ps (__m512, __mmask16, __m512, __m512);

Exceptions

 64

Real-Address Mode and Virtual-8086

VFMSUB132PD - Multiply Destination By Second Source and Subtract First Source Float64 Vectors

Description

Performs an element-by-element multiplication of float64 vector zmm1 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, then subtracts float64 vector zmm2 from the result. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm1[i+63:i] * tmpSrc3[i+63:i] - zmm2[i+63:i]
  }
}
```


Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFMSUB132PS - Multiply Destination By Second Source and Subtract First Source Float32 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm1 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then subtracts float32 vector zmm2 from the result. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = zmm1[i+31:i] * tmpSrc3[i+31:i] - zmm2[i+31:i]
  }
}
```


Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

VFMSUB213PD - Multiply First Source By Destination and Subtract Second Source Float64 Vectors

Description

Performs an element-by-element multiplication of float64 vector zmm2 and float64 vector zmm1, then subtracts the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3 from the result. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm2[i+63:i] * zmm1[i+63:i] - tmpSrc3[i+63:i]
 }
}
```


Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFMSUB213PS - Multiply First Source By Destination and Subtract Second Source Float32 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm2 and float32 vector zmm1, then subtracts the float 32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3 from the result. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
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   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] * zmm1[i+31:i] - tmpSrc3[i+31:i]
 }
}
```
Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

VFMSUB231PD - Multiply First Source By Second Source and Subtract Destination Float64 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then subtracts float32 vector zmm1 from the result. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
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   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm2[i+63:i] * tmpSrc3[i+63:i] - zmm1[i+63:i]
  }
}
```


Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFMSUB231PS - Multiply First Source By Second Source and Subtract Destination Float32 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then subtracts float32 vector zmm1 from the result. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

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   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] * tmpSrc3[i+31:i] - zmm1[i+31:i]
  }
}
```


Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

VFNMADD132PD - Multiply Destination By Second Source and Subtract From First Source Float64 Vectors

Description

Performs an element-by-element multiplication of float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, then subtracts the result from float64 vector zmm1. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
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  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {
```


```
i = 64*n// float64 operation
   zmm1[i+63:i] = -(zmm1[i+63:i] * tmpSrc3[i+63:i]) + zmm2[i+63:i]
 }
}
```
Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

```
Treat Input Denormals As Zeros :
         (MXCSR.DAZ)? YES : NO
```

```
Flush Tiny Results To Zero :
        (MXCSR.FZ)? YES : NO
```
Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFNMADD132PS - Multiply Destination By Second Source and Subtract From First Source Float32 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then subtracts the result from float32 vector zmm1. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
```


```
i = 32*n// float32 operation
   zmm1[i+31:i] = -(zmm1[i+31:i] * tmpSrc3[i+31:i]) + zmm2[i+31:i]
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VFNMADD213PD - Multiply First Source By Destination and Subtract From Second Source Float64 Vectors

Description

Performs an element-by-element multiplication of float64 vector zmm1 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, then subtracts the result from float64 vector zmm2. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress Exception Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {
    i = 64 * n
```


```
// float64 operation
   zmm1[i+63:i] = -(zmm2[i+63:i] * zmm1[i+63:i]) + tmpSrc3[i+63:i]}
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFNMADD213PS - Multiply First Source By Destination and Subtract From Second Source Float32 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm1 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, then subtracts the result from float32 vector zmm2. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress Exception Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n
```


```
// float32 operation
   zmm1[i+31:i] = -(zmm2[i+31:i] * zmm1[i+31:i]) + tmpSrc3[i+31:i]
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VFNMADD231PD - Multiply First Source By Second Source and Subtract From Destination Float64 Vectors

Description

Performs an element-by-element multiplication of float64 vector zmm2 and float64 vector zmm1, then subtracts the result from the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {
```


```
i = 64*n// float64 operation
   zmm1[i+63:i] = -(zmm2[i+63:i] * tmpSrc3[i+63:i]) + zmm1[i+63:i]
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

```
Treat Input Denormals As Zeros :
         (MXCSR.DAZ)? YES : NO
```

```
Flush Tiny Results To Zero :
        (MXCSR.FZ)? YES : NO
```
Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFNMADD231PS - Multiply First Source By Second Source and Subtract From Destination Float32 Vectors

Description

Performs an element-by-element multiplication of float32 vector zmm2 and float32 vector zmm1, then subtracts the result from the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
```


```
i = 32*n// float32 operation
   zmm1[i+31:i] = -(zmm2[i+31:i] * tmpSrc3[i+31:i]) + zmm1[i+31:i]
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VFNMSUB132PD - Multiply Destination By Second Source, Negate, and Subtract First Source Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm1 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, negates, and subtracts float64 vector zmm2. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

Table 6.11: VFNMSUB outcome when adding zeros depending on rounding-mode

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
  RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = (-(zmm1[i+63:i] * tmpSrc3[i+63:i]) - zmm2[i+63:i])
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

```
Treat Input Denormals As Zeros :
        (MXCSR.DAZ)? YES : NO
```
Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFNMSUB132PS - Multiply Destination By Second Source, Negate, and Subtract First Source Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm1 and the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, negates, and subtracts float32 vector zmm2. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

Table 6.12: VFNMSUB outcome when adding zeros depending on rounding-mode

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n
    // float32 operation
    zmm1[i+31:i] = (-(zmm1[i+31:i] * tmpSrc3[i+31:i]) - zmm2[i+31:i])
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VFNMSUB213PD - Multiply First Source By Destination, Negate, and Subtract Second Source Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm2 and float64 vector zmm1, negates, and subtracts the float 64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3. The final sum is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

Table 6.13: VFNMSUB outcome when adding zeros depending on rounding-mode

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
  RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = (-(zmm2[i+63:i] * zmm1[i+63:i]) - tmpSrc3[i+63:i])
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

```
Treat Input Denormals As Zeros :
        (MXCSR.DAZ)? YES : NO
```
Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If a memory address referencing the SS segment is

VFNMSUB213PS - Multiply First Source By Destination, Negate, and Subtract Second Source Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm2 and float32 vector zmm1, negates, and subtracts the float32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3. The final sum is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

Table 6.14: VFNMSUB outcome when adding zeros depending on rounding-mode

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n
    // float32 operation
    zmm1[i+31:i] = (-(zmm2[i+31:i] * zmm1[i+31:i]) - tmpSrc3[i+31:i])
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VFNMSUB231PD - Multiply First Source By Second Source, Negate, and Subtract Destination Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or vector float64 zmm3, negates, and subtracts float64 vector zmm1. The final result is written into float64 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

Table 6.15: VFMADDN outcome when adding zeros depending on rounding-mode

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
  RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
  RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = (-(zmm2[i+63:i] * tmpSrc3[i+63:i]) - zmm1[i+63:i])
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

```
Treat Input Denormals As Zeros :
        (MXCSR.DAZ)? YES : NO
```
Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is

VFNMSUB231PS - Multiply First Source By Second Source, Negate, and Subtract Destination Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm2 and the float 32 vector result of the swizzle/broadcast/conversion process on memory or vector float32 zmm3, negates, and subtracts float32 vector zmm1. The final result is written into float32 vector zmm1.

Intermediate values are calculated to infinite precision, and are not truncated or rounded. All operations must be performed previous to final rounding.

Table 6.16: VFMADDN outcome when adding zeros depending on rounding-mode

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n
    // float32 operation
    zmm1[i+31:i] = (-(zmm2[i+31:i] * tmpSrc3[i+31:i]) - zmm1[i+31:i])
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VGATHERDPD - Gather Float64 Vector With Signed Dword Indices

Description

A set of 8 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector $VINDEX$ with scale $SCALE$ are converted to a float64 vector. The result is written into float64 vector zmm1.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Note that accessed element by will always access 64 bytes of memory. The memory region accessed by each element will always be between *elemen_linear_address & (∼0x3F)* and *(element_linear_address & (∼0x3F)) + 63* boundaries.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully loaded.

The instruction will #GP fault if the destination vector zmm1 is the same as index vector *V INDEX*.

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 8; n++) {
 if ( ktemp[n] != 0) {
```


```
i = 64*nj = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[j+31:j] * SCALE)
    pointer[63:0] = mv_t[n]zmm1[i+63:i] = UpConvLoad<sub>f64</sub>(pointer)
    k1[n] = 0}
}
k1[15:8] = 0
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VGATHERDPS - Gather Float32 Vector With Signed Dword Indices

Description

A set of 16 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector $VINDEX$ with scale $SCALE$ are converted to a float 32 vector. The result is written into float 32 vector zmm1.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Note that accessed element by will always access 64 bytes of memory. The memory region accessed by each element will always be between *elemen_linear_address & (∼0x3F)* and *(element_linear_address & (∼0x3F)) + 63* boundaries.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully loaded.

The instruction will #GP fault if the destination vector zmm1 is the same as index vector *V INDEX*.

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
 if ( ktemp[n] != 0) {
```


```
i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]zmm1[i+31:i] = UpConvLoad<sub>f32</sub>(pointer)
    k1[n] = 0}
}
```
SIMD Floating-Point Exceptions

Invalid.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512 _mm512_i32gather_ps (__m512i, void const*, int);
- __m512 _mm512_mask_i32gather_ps (__m512, __mmask16, __m512i, void const*, int);
- __m512 _mm512_i32extgather_ps (__m512i, void const*, _MM_UPCONV_PS_ENUM, int, int);
- __m512 _mm512_mask_i32extgather_ps (__m512, __mmask16, __m512i, void const*, _MM_UPCONV_PS_ENUM, int, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

(intel)

VGATHERPF0DPS - Gather Prefetch Float32 Vector With Signed Dword Indices Into L1

Description

A set of 16 ϐloat32 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* are prefetched from memory to L1 level of cache. If any memory access causes any type of memory exception, the memory access will be considered as completed (destination mask updated) and the exception ignored.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the prefetch sequence have been prefetched and hence, the write-mask bits all are zero).

Note that accessed element by will always access 64 bytes of memory. The memory region accessed by each element will always be between *elemen_linear_address & (∼0x3F)* and *(element_linear_address & (∼0x3F)) + 63* boundaries.

This instruction has special disp $8*N$ and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully loaded.

Note that both gather and scatter prefetches set the access bit (A) in the related TLB page entry. Scatter prefetches (which prefetch data with RFO) do not set the dirty bit (D).

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
exclusive = 0
evicthintpre = MVEX.EH
```


```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
  if (ktemp[n] != 0) {
    i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
   pointer[63:0] = mv_t[n]FetchL1cacheLine(pointer, exclusive, evicthintpre)
   k1[n] = 0}
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_prefetch_i32gather_ps (__m512i, void const*, int, int);
- void _mm512_mask_prefetch_i32gather_ps (__m512i, __mmask16, void const*, int, int);
- void _mm512_prefetch_i32extgather_ps (__m512i, void const*, _MM_UPCONV_PS_ENUM, int, int);
- void _mm512_mask_prefetch_i32extgather_ps (__m512i, __mmask16, void const*, _MM_UPCONV_PS_ENUM, int, int);

Exceptions

Real-Address Mode and Virtual-8086

VGATHERPF0HINTDPD - Gather Prefetch Float64 Vector Hint With Signed Dword Indices

Description

The instruction specifies a set of 8 float64 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* as a performance hint that a real gather instruction with the same set of sources will be invoked. A programmer may execute this instruction before a real gather instruction to improve its performance.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults. This instructions does not modify any kind of architectural state (including the write-mask).

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Operation

```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 8; n++) {
  if (k1[n] := 0) {
    i = 64*ni = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[j+31:j] * SCALE)
    pointer[63:0] = mv_t[n]HintPointer(pointer)
  }
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

None

Exceptions

Real-Address Mode and Virtual-8086

If no write mask is provided or selected write-mask is k0.

VGATHERPF0HINTDPS - Gather Prefetch Float32 Vector Hint With Signed Dword Indices

Description

The instruction specifies a set of 16 float32 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* as a performance hint that a real gather instruction with the same set of sources will be invoked. A programmer may execute this instruction before a real gather instruction to improve its performance.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults. This instructions does not modify any kind of architectural state (including the write-mask).

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Operation

```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]HintPointer(pointer)
 }
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

None

Exceptions

Real-Address Mode and Virtual-8086

If no write mask is provided or selected write-mask is k0.

VGATHERPF1DPS - Gather Prefetch Float32 Vector With Signed Dword Indices Into L2

Description

A set of 16 ϐloat32 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* are prefetched from memory to L2 level of cache. If any memory access causes any type of memory exception, the memory access will be considered as completed (destination mask updated) and the exception ignored.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the prefetch sequence have been prefetched and hence, the write-mask bits all are zero).

Note that accessed element by will always access 64 bytes of memory. The memory region accessed by each element will always be between *elemen_linear_address & (∼0x3F)* and *(element_linear_address & (∼0x3F)) + 63* boundaries.

This instruction has special disp $8*N$ and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully loaded.

Note that both gather and scatter prefetches set the access bit (A) in the related TLB page entry. Scatter prefetches (which prefetch data with RFO) do not set the dirty bit (D).

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
exclusive = 0
evicthintpre = MVEX.EH
```


```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
  if (ktemp[n] != 0) {
    i = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]FetchL2cacheLine(pointer, exclusive, evicthintpre)
    k1[n] = 0}
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_prefetch_i32gather_ps (__m512i, void const*, int, int);
- void _mm512_mask_prefetch_i32gather_ps (__m512i, __mmask16, void const*, int, int);
- void _mm512_prefetch_i32extgather_ps (__m512i, void const*, _MM_UPCONV_PS_ENUM, int, int);
- void _mm512_mask_prefetch_i32extgather_ps (__m512i, __mmask16, void const*, _MM_UPCONV_PS_ENUM, int, int);

Exceptions

Real-Address Mode and Virtual-8086

VGETEXPPD - Extract Float64 Vector of Exponents from Float64 Vector

Description

Performs an element-by-element exponent extraction from the Float64 vector result of the swizzle/broadcast/conversion process on memory or Float64 vector zmm2. The result is written into Float64 vector zmm1.

GetExp() returns the (un-biased) exponent n in floating-point format. That is, when $X =$ 1/16, GetExp() returns the value *−*4, represented as C0800000 in IEEE single precision (for the single-precision version of the instruction). If the source is denormal, VGETEXP will normalize it prior to exponent extraction (unless DAZ=1).

GetExp() function follows Table [6.17](#page-311-0) when dealing with floating-point special number.

Table 6.17: GetExp() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad_{f64}(zmm2/m_t)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nzmm1[i+63:i] = GetExp(tmpSrc2[i+63:i])
```


} }

SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

VGETEXPPS - Extract Float32 Vector of Exponents from Float32 Vector

Description

Performs an element-by-element exponent extraction from the Float32 vector result of the swizzle/broadcast/conversion process on memory or Float32 vector zmm2. The result is written into Float32 vector zmm1.

GetExp() returns the (un-biased) exponent n in floating-point format. That is, when $X =$ 1/16, GetExp() returns the value *−*4, represented as C0800000 in IEEE single precision (for the single-precision version of the instruction). If the source is denormal, VGETEXP will normalize it prior to exponent extraction (unless DAZ=1).

GetExp() function follows Table [6.18](#page-314-0) when dealing with floating-point special number.

Input	Result
NaN	quietized input NaN
$+\infty$	$+\infty$
$+0$	\mathbf{x}
-0	∞
	∞

Table 6.18: GetExp() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] = GetExp(tmpSrc2[i+31:i])
```


} }

SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

VGETMANTPD - Extract Float64 Vector of Normalized Mantissas from Float64 Vector

Description

Performs an element-by-element conversion of the Float64 vector result of the swizzle/broadcast/conversion process on memory or Float64 vector zmm2 to Float64 values with the mantissa normalized to the interval specified by *interv* and sign dictated by the sign control parameter *sc*. The result is written into Float64 vector zmm1. Denormal values are explicitly normalized.

The formula for the operation is:

 $GetMant(x) = \pm 2^k |x \text{.} significant|$

where:

 $1 \leq |x \leq x \leq 1$

Exponent k is dependent on the interval range defined by *interv* and whether the exponent of the source is even or odd. The sign of the final result is determined by sc and the source sign.

GetMant() function follows Table [6.19](#page-317-0) when dealing with floating-point special numbers.

Table 6.19: GetMant() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format


```
GetNormalizedMantissa(SRC , SignCtrl, Interv)
{
  // Extracting the SRC sign, exponent and mantissa fields
  SIGN = (SignCtrl[0])? 0 : SRC[63];
  EXP = SRC[63:52];
  FRACT = (DAZ & & (EXP == 0))? 0 : SRC[51:0];
  // Check for NaN operand
  if(IsNaN(SRC)) {
     if(IsSNaN(SRC)) *set I flag*
     return QNaN(SRC)
  }
  // If SignCtrl[1] is set to 1, return NaN and set
  // exception flag if the operand is negative.
  // Note that -0.0 is included
  if(SignCtrl[1] \&c(SRC[63] == 1)) {
     *set I flag*
     return QNaN_Indefinite
  }
  // Check for +/-INF and +/-0
  if( EXP == 0x7FF & FRACTION == 0)
   || (EXP == 0 && FRACTION == 0 ) ) {
     DEST[63:0] = (SIGN << 63) | (EXP[11:0] << 52) | FRACT[51:0];
     return DEST
  }
  // Normalize denormal operands
  // note that denormal operands are treated as zero if
  // DAZ is set to 1
  if((EXP == 0) & \&& (FRACTION != 0)// JBIT is the hidden integral bit
      JBIT = 0; // Zero in case of denormal operands
      EXP = 03FFh; \frac{1}{2} // Set exponent to BIAS
```


```
While(JBIT == 0) {
         JBIT = FRACT[51]; // Obtain fraction MSB
         FRACT = FRACT \ll 1; // Normalize mantissa
         EXP--; \frac{1}{2} and adjust exponent
      }
      *set D flag*
  }
  // Apply normalization intervals
  UNBIASED_EXP = EXP - 03FFh; // get exponent in unbiased form
  IS_ODD_EXP = UNBIASED_EXP[0]; // if the unbiased exponent odd?
  if( (Interv == 10b)
      || ( (Interv == 01b) && IS ODD EXP)
      || ( (Interv == 11b) && (FRACT[51]==1)) ) {
     EXP = 03FEh; \frac{1}{2} // Set exponent to -1 (unbiased)
  }
  else {
     EXP = 03FFh; \frac{1}{2} // Set exponent to 0 (unbiased)
   }
  // form the final destination
  DEST[63:0] = (SIGN << 63) | (EXP[11:0] << 52) | FRACT[51:0];
  return DEST
}
sc = 1MM8[3:2]interv = IMMS[1:0]if(source is a register operand and MVEX.EH bit is 1) {
  if(SSS[2]==1) Supress_Exception_Flags() // SAE
  tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
 if(k1[n] != 0) {i = 64*n// float64 operation
   zmm1[i+63:i] = GetNormalizedMantissa(tmpSrc2[i+63:i], sc, interv)
 }
}
```


SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

VGETMANTPS - Extract Float32 Vector of Normalized Mantissas from Float32 Vector

Description

Performs an element-by-element conversion of the Float32 vector result of the swizzle/broadcast/conversion process on memory or Float32 vector zmm2 to Float32 values with the mantissa normalized to the interval specified by *interv* and sign dictated by the sign control parameter *sc*. The result is written into Float32 vector zmm1. Denormal values are explicitly normalized.

The formula for the operation is:

 $GetMant(x) = \pm 2^k |x \text{.} significant|$

where:

 $1 \leq |x \leq x \leq 1$

Exponent k is dependent on the interval range defined by *interv* and whether the exponent of the source is even or odd. The sign of the final result is determined by sc and the source sign.

GetMant() function follows Table [6.20](#page-322-0) when dealing with floating-point special numbers.

Table 6.20: GetMant() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format


```
GetNormalizedMantissa(SRC , SignCtrl, Interv)
{
  // Extracting the SRC sign, exponent and mantissa fields
  SIGN = (SignCtrl[0])? 0 : SRC[31];
  EXP = SRC[30:23];
  FRACT = (DAZ & (EXP == 0))? 0 : SRC[22:0];
  // Check for NaN operand
  if(IsNaN(SRC)) {
     if(IsSNaN(SRC)) *set I flag*
     return QNaN(SRC)
  }
  // If SignCtrl[1] is set to 1, return NaN and set
  // exception flag if the operand is negative.
  // Note that -0.0 is included
  if(SignCtrl[1] \& (SRC[31] == 1)) {
     *set I flag*
     return QNaN_Indefinite
  }
  // Check for +/-INF and +/-0
  if( EXP == OxFF & FRACTION == 0)
   || ( EXP == 0 \&& FRACTION == 0 ) }
     DEST[31:0] = (SIGN << 31) | (EXP[7:0] << 23) | FRACT[22:0];
     return DEST
  }
   // Apply normalization intervals
  UNBIASED_EXP = EXP - 07Fh; // get exponent in unbiased form
   IS_ODD_EXP = UNBIASED_EXP[0]; // if the unbiased exponent odd?
   if( (Interv == 10b)
      || ( (Interv == 01b) && IS_ODD_EXP)
      || ( (Interv == 11b) && (FRACT[22]==1)) ) {
```


```
EXP = 07Eh; // Set exponent to -1 (unbiased)}
   else {
     EXP = 07Fh; \frac{1}{2} // Set exponent to 0 (unbiased)
   }
  // form the final destination
  DEST[31:0] = (SIGN << 31) | (EXP[7:0] << 23) | FRACT[22:0];
  return DEST
}
sc = 1MM8[3:2]interv = IMMS[1:0]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]
} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
   zmm1[i+31:i] = GetNormalizedMantissa(tmpSrc2[i+31:i], sc, interv)
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : Not Applicable

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VGMAXABSPS - Absolute Maximum of Float32 Vectors

Description

Determines the maximum of the absolute values of each pair of corresponding elements in float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float 32 vector zmm3. The result is written into float 32 vector zmm1.

Abs() returns the absolute value of one float 32 argument. FpMax() returns the bigger of the two float32 arguments, following IEEE in general. NaN has special handling: If one source operand is NaN, then the other source operand is returned (choice made percomponent). If both are NaN, then the unchanged NaN from the first source (here zmm2) is returned. Please note that if first source is a SNaN it won't be quietized, it will be returned without any modification. This differs from the new IEEE 754-08 rules, which states that in case of an input SNaN, its quietized version should be returned instead of the other value.

Another new IEEE 754-08 rule is that $max(-0,+0) = max(+0,-0) = +0$, which honors the sign, in contrast to the comparison rules for signed zero (stated above). D3D10.0 recommends the IEEE 754-08 behavior here, but it will not be enforced; it is permissible for the result of comparing zeros to be dependent on the order of parameters, using a comparison that ignores the signs.

This instruction treats input denormals as zeros according to the DAZ control bit, but it does not flush tiny results to zero.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
FpMaxAbs(A,B)
{
  if ((A == \text{NaN}) \&& (B == \text{NaN}))return Abs(A);
  else if (A == NaN)
    return Abs(B);
  else if (B == \text{NaN})
```


```
return Abs(A);
  else if ((\text{Abs}(A) == +inf) || (\text{Abs}(B) == +inf))return +inf;
  else if (Abs(A) \geq Abs(B))return Abs(A);
  else
    return Abs(B);
}
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = FpMaxAbs(zmm2[i+31:i] , tmpSrc3[i+31:i])
  }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

VGMAXPD - Maximum of Float64 Vectors

Description

Determines the maximum value of each pair of corresponding elements in float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3. The result is written into float64 vector zmm1.

 F_D returns the bigger of the two float 32 arguments, following IEEE in general. NaN has special handling: If one source operand is NaN, then the other source operand is returned (choice made per-component). If both are NaN, then the unchanged NaN from the first source (here zmm2) is returned. Please note that if first source is a SNaN it won't be quietized, it will be returned without any modification. This differs from the new IEEE 754-08 rules, which states that in case of an input SNaN, its quietized version should be returned instead of the other value.

Another new IEEE 754-08 rule is that max $(-0,+0)$ == max $(+0,-0)$ == +0, which honors the sign, in contrast to the comparison rules for signed zero (stated above). D3D10.0 recommends the IEEE 754-08 behavior here, but it will not be enforced; it is permissible for the result of comparing zeros to be dependent on the order of parameters, using a comparison that ignores the signs.

This instruction treats input denormals as zeros according to the DAZ control bit, but it does not flush tiny results to zero.

The following table describes exception flags priority:

Table 6.21: Max exception flags priority

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

intel

```
FpMax(A,B)
{
 if ((A == -0.0) & (B == +0.0)) return B;
 if ((A == +0.0) & & (B == -0.0)) return A;
 if ((A == \text{NaN}) \&& (B == \text{NaN})) return A;
 if (A == \text{NaN}) return B;
 if (B == \text{NaN}) return A;
 if (A == -inf) return B;
 if (B == -inf) return A;
 if (A == +inf) return A;
 if (B == +inf) return B;
 if (A \geq B) return A;
 return B;
}
if(source is a register operand and MVEX.EH bit is 1) {
  if(SSS[2]==1) Supress_Exception_Flags() // SAE
  tmpSrc3[511:0] = zmm3[511:0]} else {
  tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
 if(k1[n] != 0) {i = 64*n// float64 operation
   zmm1[i+63:i] = FpMax(zmm2[i+63:i] , tmpSrc3[i+63:i])
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VGMAXPS - Maximum of Float32 Vectors

Description

Determines the maximum value of each pair of corresponding elements in float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

 F_D returns the bigger of the two float 32 arguments, following IEEE in general. NaN has special handling: If one source operand is NaN, then the other source operand is returned (choice made per-component). If both are NaN, then the unchanged NaN from the first source (here zmm2) is returned. Please note that if first source is a SNaN it won't be quietized, it will be returned without any modification. This differs from the new IEEE 754-08 rules, which states that in case of an input SNaN, its quietized version should be returned instead of the other value.

Another new IEEE 754-08 rule is that max $(-0,+0) = \max(+0,-0) = -10$, which honors the sign, in contrast to the comparison rules for signed zero (stated above). D3D10.0 recommends the IEEE 754-08 behavior here, but it will not be enforced; it is permissible for the result of comparing zeros to be dependent on the order of parameters, using a comparison that ignores the signs.

This instruction treats input denormals as zeros according to the DAZ control bit, but it does not flush tiny results to zero.

The following table describes exception flags priority:

Table 6.22: Max exception flags priority

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

intel

```
FpMax(A,B)
{
 if ((A == -0.0) && (B == +0.0)) return B;
 if ((A == +0.0) & & (B == -0.0)) return A;
 if ((A == \text{NaN}) \&& (B == \text{NaN})) return A;
 if (A == \text{NaN}) return B;
 if (B == \text{NaN}) return A;
 if (A == -inf) return B;
 if (B == -inf) return A;
 if (A == +inf) return A;
 if (B == +inf) return B;
 if (A \geq B) return A;
 return B;
}
if(source is a register operand and MVEX.EH bit is 1) {
  if(SSS[2]==1) Supress_Exception_Flags() // SAE
  tmpSrc3[511:0] = zmm3[511:0]} else {
  tmpSrc3[511:0] = SwizzUpConvLoadf32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
 if(k1[n] != 0) {i = 32*n// float32 operation
   zmm1[i+31:i] = FpMax(zmm2[i+31:i] , tmpSrc3[i+31:i])
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VGMINPD - Minimum of Float64 Vectors

Description

Determines the minimum value of each pair of corresponding elements in float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3. The result is written into float64 vector zmm1.

 $FpMin()$ returns the smaller of the two float 32 arguments, following IEEE in general. NaN has special handling: If one source operand is NaN, then the other source operand is returned (choice made per-component). If both are NaN, then the unchanged NaN from the first source (here zmm2) is returned. Please note that if first source is a SNaN it won't be quietized, it will be returned without any modification. This differs from the new IEEE 754-08 rules, which states that in case of an input SNaN, its quietized version should be returned instead of the other value.

Another new IEEE 754-08 rule is that $min(-0,+0) == min(+0,-0) == -0$, which honors the sign, in contrast to the comparison rules for signed zero (stated above). D3D10.0 recommends the IEEE 754-08 behavior here, but it will not be enforced; it is permissible for the result of comparing zeros to be dependent on the order of parameters, using a comparison that ignores the signs.

This instruction treats input denormals as zeros according to the DAZ control bit, but it does not flush tiny results to zero.

The following table describes exception flags priority:

Table 6.23: Min exception flags priority

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

intel

```
FpMin(A,B)
{
 if ((A == -0.0) && (B == +0.0)) return A;
 if ((A == +0.0) & & (B == -0.0)) return B;
 if ((A == \text{NaN}) \&& (B == \text{NaN})) return A;
 if (A == \text{NaN}) return B;
 if (B == \text{NaN}) return A;
 if (A == -inf) return A;
 if (B == -inf) return B;
 if (A == +inf) return B;
 if (B == +inf) return A;
 if (A < B) return A;
 return B;
}
if(source is a register operand and MVEX.EH bit is 1) {
  if(SSS[2]==1) Supress_Exception_Flags() // SAE
  tmpSrc3[511:0] = zmm3[511:0]} else {
  tmpSrc3[511:0] = SwizzUpConvLoad_{f64}(zmm3/m_t)}
for (n = 0; n < 8; n++) {
 if(k1[n] != 0) {i = 64*n// float64 operation
   zmm1[i+63:i] = FpMin(zmm2[i+63:i] , tmpSrc3[i+63:i])
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VGMINPS - Minimum of Float32 Vectors

Description

Determines the minimum value of each pair of corresponding elements in float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

 $FpMin()$ returns the smaller of the two float 32 arguments, following IEEE in general. NaN has special handling: If one source operand is NaN, then the other source operand is returned (choice made per-component). If both are NaN, then the unchanged NaN from the first source (here zmm2) is returned. Please note that if first source is a SNaN it won't be quietized, it will be returned without any modification. This differs from the new IEEE 754-08 rules, which states that in case of an input SNaN, its quietized version should be returned instead of the other value.

Another new IEEE 754-08 rule is that $min(-0,+0) == min(+0,-0) == -0$, which honors the sign, in contrast to the comparison rules for signed zero (stated above). D3D10.0 recommends the IEEE 754-08 behavior here, but it will not be enforced; it is permissible for the result of comparing zeros to be dependent on the order of parameters, using a comparison that ignores the signs.

This instruction treats input denormals as zeros according to the DAZ control bit, but it does not flush tiny results to zero.

The following table describes exception flags priority:

Table 6.24: Min exception flags priority

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

intel

```
FpMin(A,B)
{
 if ((A == -0.0) & (B == +0.0)) return A;
 if ((A == +0.0) & & (B == -0.0)) return B;
 if ((A == \text{NaN}) \&& (B == \text{NaN})) return A;
 if (A == \text{NaN}) return B;
 if (B == \text{NaN}) return A;
 if (A == -inf) return A;
 if (B == -inf) return B;
 if (A == +inf) return B;
 if (B == +inf) return A;
 if (A < B) return A;
 return B;
}
if(source is a register operand and MVEX.EH bit is 1) {
  if(SSS[2]==1) Supress_Exception_Flags() // SAE
  tmpSrc3[511:0] = zmm3[511:0]} else {
  tmpSrc3[511:0] = SwizzUpConvLoadf32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
 if(k1[n] != 0) {i = 32*n// float32 operation
   zmm1[i+31:i] = FpMin(zmm2[i+31:i] , tmpSrc3[i+31:i])
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512 _mm512_mask_gmin_ps(__m512, __mmask16, __m512, __m512);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VLOADUNPACKHD - Load Unaligned High And Unpack To Doubleword Vector

Description

The high-64-byte portion of the byte/word/doubleword stream starting at the elementaligned address (*mt−*64)is loaded, converted and expanded into the write-mask-enabled elements of doubleword vector zmm1. The number of set bits in the write-mask determines the length of the converted doubleword stream, as each converted doubleword is mapped to exactly one of the doubleword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted doublewords (if any) in the stream that occur at or after the first 64-byte-aligned address following $(m_t - 64)$ (that is, in the high cache line of the memory stream for the current implementation). Elements in zmm1 that don't map to those stream doublewords are left unchanged. The vloadunpackld instruction is used to load the part of the stream before the first 64-byte-aligned address preceding *mt*.

In conjunction with vloadunpackld, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpackld, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpackld v0 {k1}, [rax]
vloadunpackhd v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or

#SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOfi32(SSS[2:0])
foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 16; n++) {
   if(k1[n] != 0) {if (foundNext64BytesBoundary == false) {
         if ( ( ((pointer + (loadOffset+1)*upSize ) % 64) == 0 ) {
            foundNext64BytesBoundary = true
         }
     } else {
         i = 32*nzmm1[i+31:i] = UpConvLoadi32(pointer + upSize*loadOffset)
     }
     loadOffset++
  }
}
```
Flags Affected

None.

Memory Up-conversion: U_{i32}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VLOADUNPACKHPD - Load Unaligned High And Unpack To Float64 Vector

Description

The high-64-byte portion of the quadword stream starting at the element-aligned address (*m^t −* 64) is loaded, converted and expanded into the write-mask-enabled elements of quadword vector zmm1. The number of set bits in the write-mask determines the length of the converted quadword stream, as each converted quadword is mapped to exactly one of the quadword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted quadwords (if any) in the stream that occur at or after the first 64-byte-aligned address following $(m_t - 64)$ (that is, in the high cache line of the memory stream for the current implementation). Elements in zmm1 that don't map to those stream quadwords are left unchanged. The vloadunpacklpd instruction is used to load the part of the stream before the first 64-byte-aligned address preceding *mt*.

In conjunction with vloadunpacklpd, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpacklpd, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpacklpd v0 {k1}, [rax]
vloadunpackhpd v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or #SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte

boundary. Additionally, A/D bits in the page table will not be updated.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOf_{f64}(SSS[2:0])foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 8; n++) {
   if(k1[n] != 0) {if (foundNext64BytesBoundary == false) {
         if ( ( ((pointer + (loadOffset+1)*upSize ) % 64) == 0 ) {
            foundNext64BytesBoundary = true
         }
      } else {
         i = 64*nzmm1[i+63:i] = UpConvLoadf64(pointer + upSize*loadOffset)
      }
      loadOffset++
  }
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VLOADUNPACKHPS - Load Unaligned High And Unpack To Float32 Vector

Description

The high-64-byte portion of the byte/word/doubleword stream starting at the elementaligned address (*mt−*64)is loaded, converted and expanded into the write-mask-enabled elements of doubleword vector zmm1. The number of set bits in the write-mask determines the length of the converted doubleword stream, as each converted doubleword is mapped to exactly one of the doubleword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted doublewords (if any) in the stream that occur at or after the first 64-byte-aligned address following $(m_t - 64)$ (that is, in the high cache line of the memory stream for the current implementation). Elements in zmm1 that don't map to those stream doublewords are left unchanged. The vloadunpacklps instruction is used to load the part of the stream before the first 64-byte-aligned address preceding *mt*.

In conjunction with vloadunpacklps, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpacklps, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpacklps v0 {k1}, [rax]
vloadunpackhps v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or

#SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOff32(SSS[2:0])
foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 16; n++) {
   if(k1[n] != 0) {if (foundNext64BytesBoundary == false) {
         if ( ( ((pointer + (loadOffset+1)*upSize ) % 64) == 0 ) {
            foundNext64BytesBoundary = true
         }
      } else {
         i = 32*n
         zmm1[i+31:i] = UpConvLoad<sub>f32</sub>(pointer + upSize*loadOffset)
      }
      loadOffset++
   }
}
```
SIMD Floating-Point Exceptions

Invalid.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VLOADUNPACKHQ - Load Unaligned High And Unpack To Int64 Vector

Description

The high-64-byte portion of the quadword stream starting at the element-aligned address (*m^t −* 64) is loaded, converted and expanded into the write-mask-enabled elements of quadword vector zmm1. The number of set bits in the write-mask determines the length of the converted quadword stream, as each converted quadword is mapped to exactly one of the quadword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted quadwords (if any) in the stream that occur at or after the first 64-byte-aligned address following $(m_t - 64)$ (that is, in the high cache line of the memory stream for the current implementation). Elements in zmm1 that don't map to those stream quadwords are left unchanged. The vloadunpacklq instruction is used to load the part of the stream before the first 64-byte-aligned address preceding m_t .

In conjunction with vloadunpacklq, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpacklq, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpacklq v0 {k1}, [rax]
vloadunpackhq v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or #SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOfi64(SSS[2:0])
foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 8; n++) {
   if(k1[n] != 0) {if (foundNext64BytesBoundary == false) {
         if ( ( ((pointer + (loadOffset+1)*upSize ) % 64) == 0 ) {
            foundNext64BytesBoundary = true
         }
     } else {
         i = 64*nzmm1[i+63:i] = UpConvLoadi64(pointer + upSize*loadOffset)
     }
     loadOffset++
  }
}
```
Flags Affected

None.

Memory Up-conversion: U*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VLOADUNPACKLD - Load Unaligned Low And Unpack To Doubleword Vector

Description

The low-64-byte portion of the byte/word/doubleword stream starting at the elementaligned address *m^t* is loaded, converted and expanded into the write-mask-enabled elements of doubleword vector zmm1. The number of set bits in the write-mask determines the length of the converted doubleword stream, as each converted doubleword is mapped to exactly one of the doubleword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted doublewords (if any) in the stream that occur before the first 64-byte-aligned address following m_t (that is, in the low cache line of the memory stream in the current implementation). Elements in zmm1 that don't map to those converted stream doublewords are left unchanged. The vloadunpackhd instruction is used to load the part of the stream at or after the first 64-byte-aligned address preceding *mt*.

In conjunction with vloadunpackhd, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpackhd, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpackld v0 {k1}, [rax]
vloadunpackhd v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOfi32(SSS[2:0])
for(n = 0; n < 16; n++) {
  i = 32*n
  if (k1[n] != 0) {
    zmm1[i+31:i] = UpConvLoadi32(mt+upSize*loadOffset)
   loadOffset++
   if ( (m_t + upSize*loadOffset) % 64) == 0) {
     break
   }
 }
}
```
Flags Affected

None.

Memory Up-conversion: U_{i32}

Intel® C/C++ Compiler Intrinsic Equivalent

- _MM_UPCONV_EPI32_ENUM, int);
- __m512i _mm512_loadunpacklo_epi32 (__m512i, void const*);
- __m512i _mm512_mask_loadunpacklo_epi32 (__m512i, __mmask16, void const*);

Exceptions

VLOADUNPACKLPD - Load Unaligned Low And Unpack To Float64 Vector

Description

The low-64-byte portion of the quadword stream starting at the element-aligned address m_t is loaded, converted and expanded into the write-mask-enabled elements of quadword vector zmm1. The number of set bits in the write-mask determines the length of the converted quadword stream, as each converted quadword is mapped to exactly one of the quadword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted quadwords (if any) in the stream that occur before the first 64-byte-aligned address following m_t (that is, in the low cache line of the memory stream in the current implementation). Elements in zmm1 that don't map to those converted stream quadwords are left unchanged. The vloadunpackhq instruction is used to load the part of the stream at or after the first 64-byte-aligned address preceding m_t .

In conjunction with vloadunpackhpd, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpackhpd, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpacklpd v0 {k1}, [rax]
vloadunpackhpd v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding

bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOf_{f64}(SSS[2:0])for(n = 0; n < 8; n^{++}) {
  i = 64*nif (k1[n] := 0) {
    zmm1[i+63:i] = UpConvLoad<sub>f64</sub>(m<sub>t</sub>+upSize*load0ffset)loadOffset++
    if ( ((m_t + upSize*loadOffset) % 64) == 0) {
      break
    }
  }
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512d _mm512_extloadunpacklo_pd (__m512d, void const*, _MM_UPCONV_PD_ENUM, int);
- __m512d _mm512_mask_extloadunpacklo_pd (__m512d, __mmask8, void const*, _MM_UPCONV_PD_ENUM, int);
- __m512d _mm512_loadunpacklo_pd (__m512d, void const*);
- __m512d _mm512_mask_loadunpacklo_pd (__m512d, __mmask8, void const*);

Exceptions

(intel)

VLOADUNPACKLPS - Load Unaligned Low And Unpack To Float32 Vector

Description

The low-64-byte portion of the byte/word/doubleword stream starting at the elementaligned address *m^t* is loaded, converted and expanded into the write-mask-enabled elements of doubleword vector zmm1. The number of set bits in the write-mask determines the length of the converted doubleword stream, as each converted doubleword is mapped to exactly one of the doubleword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted doublewords (if any) in the stream that occur before the ϐirst 64-byte-aligned address following *m^t* (that is, in the low cache line of the memory stream in the current implementation). Elements in zmm1 that don't map to those converted stream doublewords are left unchanged. The vloadunpackhd instruction is used to load the part of the stream at or after the first 64-byte-aligned address preceding m_t .

In conjunction with vloadunpackhps, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpackhps, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpacklps v0 {k1}, [rax]
vloadunpackhps v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.

This instruction is write-masked, so only those elements with the corresponding bit set in

vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOf_{f32}(SSS[2:0])for(n = 0; n < 16; n++) {
  i = 32*nif (k1[n] := 0) {
    zmm1[i+31:i] = UpConvLoad<sub>f32</sub>(m<sub>t</sub>+upSize*load0ffset)loadOffset++
    if ( (m_t + upSize*loadOffset) % 64) == 0) {
      break
    }
  }
}
```
SIMD Floating-Point Exceptions

Invalid.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512 _mm512_extloadunpacklo_ps (__m512, void const*, _MM_UPCONV_PS_ENUM, int);
- __m512 _mm512_mask_extloadunpacklo_ps (__m512, __mmask16, void const*, _MM_UPCONV_PS_ENUM, int);
- __m512 _mm512_loadunpacklo_ps (__m512, void const*);
- __m512 _mm512_mask_loadunpacklo_ps (__m512, __mmask16, void const*);

Exceptions

VLOADUNPACKLQ - Load Unaligned Low And Unpack To Int64 Vector

Description

The low-64-byte portion of the quadword stream starting at the element-aligned address m_t is loaded, converted and expanded into the write-mask-enabled elements of quadword vector zmm1. The number of set bits in the write-mask determines the length of the converted quadword stream, as each converted quadword is mapped to exactly one of the quadword elements in zmm1, skipping over write-masked elements of zmm1.

This instruction only transfers those converted quadwords (if any) in the stream that occur before the first 64-byte-aligned address following m_t (that is, in the low cache line of the memory stream in the current implementation). Elements in zmm1 that don't map to those converted stream quadwords are left unchanged. The vloadunpackhq instruction is used to load the part of the stream at or after the first 64-byte-aligned address preceding m_t .

In conjunction with vloadunpackhq, this instruction is useful for re-expanding data that was packed into a queue. Also in conjunction with vloadunpackhq, it allows unaligned vector loads (that is, vector loads that are only element-wise, not vector-wise, aligned); use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector load would be:

```
// assume memory location is pointed by register rax
vloadunpacklq v0 {k1}, [rax]
vloadunpackhq v0 {k1}, [rax+64]
```
This instruction does not have broadcast support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note that this instruction will always access 64 bytes of memory. The memory region accessed will always be between *linear_address & (∼0x3F)* and *(linear_address & (∼0x3F)) + 63* boundaries.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are modified in zmm1. Elements in zmm1 with the corresponding

bit clear in vector mask1 retain their previous values. However, see above for unusual aspects of the write-mask's operation with this instruction.

Operation

```
loadOffset = 0upSize = UpConvLoadSizeOfi64(SSS[2:0])
for(n = 0; n < 8; n^{++}) {
  i = 64*nif (k1[n] != 0) {
    zmm1[i+63:i] = UpConvLoadi64(mt+upSize*loadOffset)
   loadOffset++
   if ( ((m_t + upSize*loadOffset) % 64) == 0) {
      break
   }
  }
}
```
Flags Affected

None.

Memory Up-conversion: U*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512i _mm512_extloadunpacklo_epi64 (__m512i, void const*, _MM_UPCONV_EPI64_ENUM, int); __m512i _mm512_mask_extloadunpacklo_epi64 (__m512i, __mmask8, void const*, _MM_UPCONV_EPI64_ENUM, int);
- __m512i _mm512_loadunpacklo_epi64 (__m512i, void const*);
- __m512i _mm512_mask_loadunpacklo_epi64 (__m512i, __mmask8, void const*);

Exceptions

(intel)

VLOG2PS - Vector Logarithm Base-2 of Float32 Vector

Description

Computes the element-by-element logarithm base-2 of the float 32 vector on memory or float32 vector zmm2. The result is written into float32 vector zmm1.

- 1. 4ulp of relative error when the source value is within the intervals (0, 0.5) or $(2, \infty]$
- 2. absolute error less than 2^{-21} within the interval $[0.5, 2]$

For an input value of +/ *−* 0 the instruction returns *−∞* and sets the Divide-By-Zero flag (#Z). Negative numbers (including $-\infty$) should return the canonical NaN and set the Invalid flag (#I). Note however that this instruction treats input denormals as zeros of the same sign, so for denormal negative inputs it returns *−∞* and sets the Divide-By-Zero status flag. If any source element is NaN, the quietized NaN source value is returned for that element (and #I is raised for input sNaNs).

Current implementation of this instruction does not support any SwizzUpConv setting other than "no broadcast and no conversion"; any other SwizzUpConv setting will result in an Invalid Opcode exception.

log2_DX() function follows Table [6.25](#page-371-0) when dealing with floating-point special numbers.

Table 6.25: vlog2_DX() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
tmpSrc2[511:0] = zmm2/mt
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
}
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
   i = 32*nzmm1[i+31:i] = vlog2_DX(tmpSrc2[i+31:i])}
}
```
SIMD Floating-Point Exceptions

Invalid, Zero.

Denormal Handling

Treat Input Denormals As Zeros : YES

Flush Tiny Results To Zero : YES

Register Swizzle

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Protected and Compatibility Mode

64 bit Mode

VMOVAPD - Move Aligned Float64 Vector

Description

Moves float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm2 into float64 vector zmm1 or down-converts and stores float64 vector zmm2 into destination memory.

This instruction is write-masked, so only those elements with the corresponding bit(s) set in the vector mask (k1) register are computed and stored into register/memory. Elements in register/memory with the corresponding bit(s) clear in the vector mask register are maintained with the previous value.

Operation

```
DESTINATION IS A VECTOR OPERAND
   if(source is a register operand) {
      if(MVEX.EH==1) {
         tmpSrc2[511:0] = zmm2[511:0]} else {
         tmpSrc2[511:0] = SwizzUpConvLoad<sub>f64</sub> (zmm2)}
   } else {
      tmpSrc2[511:0] = UpConvLoad_{f64}(m_t)}
  for (n = 0; n < 8; n++) {
     if (k1[n] := 0) {
       i = 64*nzmm1[i+63:i] = tmpSrc2[i+63:i])
     }
   }
```
DESTINATION IS A MEMORY OPERAND

 $downSize = DownConvStoreSizeOf_{f64}(SSS[2:0])$

```
for(n = 0; n < 8; n^{++}) {
```
CHAPTER 6. INSTRUCTION DESCRIPTIONS


```
if (k1[n] != 0) {
    i = 64*ntmp = DownConvStoref64(zmm1[i+63:i], SSS[2:0])
   if(downSize == 8) {
     MemStore(m_t+8*n) = tmp[63:0]}
 }
}
```
SIMD Floating-Point Exceptions

DESTINATION IS A VECTOR OPERAND: None. DESTINATION IS A MEMORY OPERAND: None.

Memory Up-conversion: U_{f64}

Register Swizzle: S_{f64}

MVEX.EH=0

Memory Down-conversion: D_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

__m512d _mm512_mask_mov_pd (__m512d, __mmask8, __m512d);

Exceptions

VMOVAPS - Move Aligned Float32 Vector

Description

Moves float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm2 into float32 vector zmm1 or down-converts and stores float32 vector zmm2 into destination memory.

This instruction is write-masked, so only those elements with the corresponding bit(s) set in the vector mask (k1) register are computed and stored into register/memory. Elements in register/memory with the corresponding bit(s) clear in the vector mask register are maintained with the previous value.

Operation

```
DESTINATION IS A VECTOR OPERAND
   if(source is a register operand) {
      if(MVEX.EH==1) {
         tmpSrc2[511:0] = zmm2[511:0]} else {
         tmpSrc2[511:0] = SwizzUpConvLoadf32(zmm2)
     }
   } else {
     tmpSrc2[511:0] = UpConvLoad_{f32}(m_t)}
  for (n = 0; n < 16; n++) {
     if (k1[n] := 0) {
       i = 32*n
      zmm1[i+31:i] = tmpSrc2[i+31:i])
    }
   }
```
DESTINATION IS A MEMORY OPERAND

downSize = DownConvStoreSizeOf*f*32(SSS[2:0])

for($n = 0$; $n < 16$; n^{++}) {


```
if (k1[n] := 0) {
    i = 32*ntmp = DownConvStoref32(zmm1[i+31:i], SSS[2:0])
   if(downSize == 4) {
     MemStore(m_t+4*n) = tmp[31:0]} else if(downSize == 2) {
     MemStore(m_t + 2*n) = tmp[15:0]} else if(downSize == 1) {
     MemStore(m_t+n) = tmp[7:0]}
 }
}
```
SIMD Floating-Point Exceptions

DESTINATION IS A VECTOR OPERAND: Invalid.

DESTINATION IS A MEMORY OPERAND: Overflow, Underflow, Invalid, Precision, Denormal.

Memory Up-conversion: U_{f32}

Register Swizzle: S_{f32}

MVEX.EH=0

Memory Down-conversion: D_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

VMOVDQA32 - Move Aligned Int32 Vector

Description

Moves int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm2 into int32 vector zmm1 or down-converts and stores int32 vector zmm2 into destination memory.

This instruction is write-masked, so only those elements with the corresponding bit(s) set in the vector mask (k1) register are computed and stored into register/memory. Elements in register/memory with the corresponding bit(s) clear in the vector mask register are maintained with the previous value.

Operation

DESTINATION IS A VECTOR OPERAND

```
if(source is a register operand) {
  if(MVEX.EH==1) { }tmpSrc2[511:0] = zmm2[511:0]} else {
      tmpSrc2[511:0] = SwizzUpConvLoadi32(zmm2)
   }
} else {
  tmpSrc2[511:0] = UpConvLoad_{i32}(m_t)}
for (n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    i = 32*nzmm1[i+31:i] = tmpSrc2[i+31:i])
  }
}
```
DESTINATION IS A MEMORY OPERAND

downSize = DownConvStoreSizeOf*i*32(SSS[2:0])

CHAPTER 6. INSTRUCTION DESCRIPTIONS


```
for(n = 0; n < 16; n++) {
  if (k1[n] != 0) {
   i = 32*ntmp = DownConvStorei32(zmm1[i+31:i], SSS[2:0])
    if(downSize == 4) {
     MemStore(m_t+4*n) = tmp[31:0]} else if(downSize == 2) {
     MemStore(m_t+2*n) = tmp[15:0]} else if(downSize == 1) {
     MemStore(m_t+n) = tmp[7:0]}
 }
}
```
Flags Affected

DESTINATION IS A VECTOR OPERAND: None. DESTINATION IS A MEMORY OPERAND: None.

Memory Up-conversion: U_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Memory Down-conversion: D*i*³²

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_mov_epi32 (__m512i, __mmask16, __m512i);

Exceptions

VMOVDQA64 - Move Aligned Int64 Vector

Description

Moves int64 vector result of the swizzle/broadcast/conversion process on memory or int64 vector zmm2 into int64 vector zmm1 or down-converts and stores int64 vector zmm2 into destination memory.

This instruction is write-masked, so only those elements with the corresponding bit(s) set in the vector mask (k1) register are computed and stored into register/memory. Elements in register/memory with the corresponding bit(s) clear in the vector mask register are maintained with the previous value.

Operation

DESTINATION IS A VECTOR OPERAND

```
if(source is a register operand) {
  if(MVEX.EH==1) { }tmpSrc2[511:0] = zmm2[511:0]} else {
      tmpSrc2[511:0] = SwizzUpConvLoadi64(zmm2)
   }
} else {
  tmpSrc2[511:0] = UpConvLoad_{i64}(m_t)}
for (n = 0; n < 8; n++) {
  if (k1[n] != 0) {
    i = 64*nzmm1[i+63:i] = tmpSrc2[i+63:i])
 }
}
```
DESTINATION IS A MEMORY OPERAND

downSize = DownConvStoreSizeOf*i*64(SSS[2:0])


```
for(n = 0; n < 8; n^{++}) {
  if (k1[n] != 0) {
   i = 64*ntmp = DownConvStorei64(zmm1[i+63:i], SSS[2:0])
    if(downSize == 8) {
      MemStore(mt+8*n) = tmp[63:0]
    }
 }
}
```
Flags Affected

DESTINATION IS A VECTOR OPERAND: None. DESTINATION IS A MEMORY OPERAND: None.

Memory Up-conversion: U*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Memory Down-conversion: D*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_mov_epi64 (__m512i, __mmask8, __m512i);

Exceptions

VMOVNRAPD - Store Aligned Float64 Vector With No-Read Hint

Description

Stores float64 vector zmm1 (or a down-converted version of it) into destination memory with a No-Read hint for the case the whole vector is going to be written into memory. This instruction is intended to speed up the case of stores in streaming kernels where we want to avoid wasting memory bandwidth by being forced to read the original content of entire cache lines from memory when we overwrite their whole contents completely.

In Knights Corner, this instruction is able to optimize memory bandwidth in case of a cache miss and avoid reading the original contents of the memory destination operand if the following conditions hold true:

- The instruction does not use a write-mask (MVEX.aaa=000).
- The instruction does not perform any kind of down-conversion (MVEX.SSS=000).

Note that this instruction is encoded by forcing MVEX.EH bit to 0. The Eviction Hint does not have any effect on this instruction.

The No-Read directive is intended as a performance hint and could be ignored by a given processor implementation.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are stored to memory. Elements in the destination memory vector with the corresponding bit clear in k1 register retain their previous value.

Operation

DESTINATION IS A MEMORY OPERAND

```
downSize = DownConvStoreSizeOff64(SSS[2:0])
for(n = 0; n < 8; n^{++}) {
 if (k1[n] := 0) {
    i = 64*ntmp = DownConvStore_{f64}(zmm1[i+63:i], SSS[2:0])if(downSize == 8) {
      MemStore(mt+8*n) = tmp[63:0]
    }
 }
}
```


SIMD Floating-Point Exceptions

None.

Memory Down-conversion: D_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm512_storenr_pd(void*, _m512d);

Exceptions

Real-Address Mode and Virtual-8086

If preceded by any REX, F0, F2, F3, or 66 prefixes.

#NM If CR0.TS[bit 3]=1.

VMOVNRAPS - Store Aligned Float32 Vector With No-Read Hint

Description

Stores float 32 vector zmm1 (or a down-converted version of it) into destination memory with a No-Read hint for the case the whole vector is going to be written into memory. This instruction is intended to speed up the case of stores in streaming kernels where we want to avoid wasting memory bandwidth by being forced to read the original content of entire cache lines from memory when we overwrite their whole contents completely.

In Knights Corner, this instruction is able to optimize memory bandwidth in case of a cache miss and avoid reading the original contents of the memory destination operand if the following conditions hold true:

- The instruction does not use a write-mask (MVEX.aaa=000).
- The instruction does not perform any kind of down-conversion (MVEX.SSS=000).

Note that this instruction is encoded by forcing MVEX.EH bit to 0. The Eviction Hint does not have any effect on this instruction.

The No-Read directive is intended as a performance hint and could be ignored by a given processor implementation.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are stored to memory. Elements in the destination memory vector with the corresponding bit clear in k1 register retain their previous value.

Operation

DESTINATION IS A MEMORY OPERAND

```
downSize = DownConvStoreSizeOff32(SSS[2:0])
for(n = 0; n < 16; n++) {
  if (k1[n] := 0) {
   i = 32*ntmp = DownConvStore_{f32}(zmm1[i+31:i], SSS[2:0])if(downSize == 4) {
     MemStore(m_t+4*n) = tmp[31:0]} else if(downSize == 2) {
     MemStore(m_t+2*n) = tmp[15:0]} else if(downSize == 1) {
```


```
MemStore(m_t+n) = tmp[7:0]}
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Memory Down-conversion: D_{*f*32}

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm512_storenr_ps(void*, _m512);

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

CHAPTER 6. INSTRUCTION DESCRIPTIONS

VMOVNRNGOAPD - Non-globally Ordered Store Aligned Float64 Vector With No-Read Hint

Description

Stores float64 vector zmm1 (or a down-converted version of it) into destination memory with a No-Read hint for the case the whole vector is going to be written into memory, using a weakly-ordered memory consistency model (i.e. stores performed with these instruction are not globally ordered, and subsequent stores from the same thread can be observed before them).

This instruction is intended to speed up the case of stores in streaming kernels where we want to avoid wasting memory bandwidth by being forced to read the original content of entire cache lines from memory when we overwrite their whole contents completely. This instruction takes advantage of the weakly-ordered memory consistency model to increase the throughput at which this type of write operations can be performed. Due to the same reason, a fencing operation implemented with SFENCE, MFENCE or CPUID instructions should be used in conjunction with this instruction if multiple threads are reading/writing the memory operand location (note that Knights Corner does not implement SFENCE nor MFENCE).

In Knights Corner, this instruction is able to optimize memory bandwidth in case of a cache miss and avoid reading the original contents of the memory destination operand if the following conditions hold true:

- The instruction does not use a write-mask (MVEX.aaa=000).
- The instruction does not perform any kind of down-conversion (MVEX.SSS=000).

Note that this instruction is encoded by forcing MVEX.EH bit to 1. The Eviction Hint does not have any effect on this instruction.

The No-Read directive is intended as a performance hint and could be ignored by a given processor implementation.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are stored to memory. Elements in the destination memory vector with the corresponding bit clear in k1 register retain their previous value.

Operation

DESTINATION IS A MEMORY OPERAND

```
downSize = DownConvStoreSizeOf<sub>f64</sub>(SSS[2:0])
for(n = 0; n < 8; n++) {
  if (k1[n] != 0) {
    i = 64*ntmp = DownConvStoref64(zmm1[i+63:i], SSS[2:0])
    if(downSize == 8) {
      MemStore(m_t + 8*n) = tmp[63:0]}
  }
}
```
SIMD Floating-Point Exceptions

None.

Memory Down-conversion: D_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm512_storenrngo_pd(void*, __m512d);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

CHAPTER 6. INSTRUCTION DESCRIPTIONS

(intel)

VMOVNRNGOAPS - Non-globally Ordered Store Aligned Float32 Vector With No-Read Hint

Description

Stores float32 vector zmm1 (or a down-converted version of it) into destination memory with a No-Read hint for the case the whole vector is going to be written into memory, using a weakly-ordered memory consistency model (i.e. stores performed with these instruction are not globally ordered, and subsequent stores from the same thread can be observed before them).

This instruction is intended to speed up the case of stores in streaming kernels where we want to avoid wasting memory bandwidth by being forced to read the original content of entire cache lines from memory when we overwrite their whole contents completely. This instruction takes advantage of the weakly-ordered memory consistency model to increase the throughput at which this type of write operations can be performed. Due to the same reason, a fencing operation implemented with SFENCE, MFENCE or CPUID instructions should be used in conjunction with this instruction if multiple threads are reading/writing the memory operand location (note that Knights Corner does not implement SFENCE nor MFENCE).

In Knights Corner, this instruction is able to optimize memory bandwidth in case of a cache miss and avoid reading the original contents of the memory destination operand if the following conditions hold true:

- The instruction does not use a write-mask (MVEX.aaa=000).
- The instruction does not perform any kind of down-conversion (MVEX.SSS=000).

Note that this instruction is encoded by forcing MVEX.EH bit to 1. The Eviction Hint does not have any effect on this instruction.

The No-Read directive is intended as a performance hint and could be ignored by a given processor implementation.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are stored to memory. Elements in the destination memory vector with the corresponding bit clear in k1 register retain their previous value.

Operation

DESTINATION IS A MEMORY OPERAND

```
downSize = DownConvStoreSizeOff32(SSS[2:0])
for(n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    i = 32*ntmp = DownConvStoref32(zmm1[i+31:i], SSS[2:0])
    if(downSize == 4) {
     MemStore(m_t+4*n) = tmp[31:0]} else if(downSize == 2) {
     MemStore(m_t + 2*n) = tmp[15:0]} else if(downSize == 1) {
     MemStore(m_t+n) = tmp[7:0]}
 }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Memory Down-conversion: D_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm512_storenrngo_ps(void*, _m512);

Exceptions

(intel)

Real-Address Mode and Virtual-8086

VMULPD - Multiply Float64 Vectors

Description

Performs an element-by-element multiplication between float64 vector zmm2 and the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3. The result is written into float64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm2[i+63:i] * tmpSrc3[i+63:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

If preceded by any REX, F0, F2, F3, or 66 prefixes.

VMULPS - Multiply Float32 Vectors

Description

Performs an element-by-element multiplication between float32 vector zmm2 and the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] * tmpSrc3[i+31:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

If preceded by any REX, F0, F2, F3, or 66 prefixes.

VPACKSTOREHD - Pack And Store Unaligned High From Int32 Vector

Description

Packs and down-converts the mask-enabled elements of int32 vector zmm1 into a byte/word/doubleword stream logically mapped starting at element-aligned address (*m^t −* 64), and stores the high-64-byte elements of that stream (those elements of the stream that map at or after the first 64-byte-aligned address following $(m_t - 64)$, the high cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstoreld instruction is used to store the part of the stream before the first 64byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFFFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstoreld, this instruction is useful for packing data into a queue. Also in conjunction with vpackstoreld, it allows unaligned vector stores (that is, vector stores that are only element-wise , not vector-wise, aligned); just use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstoreld [rax] {k1}, v0
vpackstorehd [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp 8^*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or #SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.

intel

```
storeOffset = 0
downSize = DownConvStoreSizeOfi32(SSS[2:0])
foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 16; n++) {
   if(k1[n] != 0) {
     if (foundNext64BytesBoundary == false) {
         if ( ( (pointer + (storeOffset+1)*downSize) % 64) == 0 ) {
            foundNext64BytesBoundary = true
         }
     } else {
         i = 32*ntmp = DownConvStorei32(zmm1[i+31:i], SSS[2:0])
         if(downSize == 4) {
          MemStore(pointer + storeOffset*4) = tmp[31:0]
         } else if(downSize == 2) {
          MemStore(pointer + storeOffset*2) = tmp[15:0]
         } else if(downSize == 1) {
           MemStore(pointer + storeOffset) = tmp[7:0]
         }
     }
     storeOffset++
  }
}
```
Flags Affected

None.

Memory Down-conversion: D_{i32}

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

intel

VPACKSTOREHPD - Pack And Store Unaligned High From Float64 Vector

Description

Packs and down-converts the mask-enabled elements of float64 vector zmm1 into a ϐloat64 stream logically mapped starting at element-aligned address (*mt−*64), and stores the high-64-byte elements of that stream (those elements of the stream that map at or after the first 64-byte-aligned address following $(m_t - 64)$, the high cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorelpd instruction is used to store the part of the stream before the first 64byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorelpd, this instruction is useful for packing data into a queue. Also in conjunction with vpackstorelpd, it allows unaligned vector stores (that is, vector stores that are only element-wise , not vector-wise, aligned); just use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstorelpd [rax] {k1}, v0
vpackstorehpd [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp 8^*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or #SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.


```
storeOffset = 0downSize = DownConvStoreSizeOf_{f64}(SSS[2:0])foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 8; n++) {
   if(k1[n] != 0) {if (foundNext64BytesBoundary == false) {
         if ( ( (pointer + (storeOffset+1)*downSize) % 64) == 0 ) {
           foundNext64BytesBoundary = true
         }
     } else {
         i = 64*ntmp = DownConvStoref64(zmm1[i+63:i], SSS[2:0])
         if(downSize == 8) {
          MemStore(pointer + storeOffset*8) = tmp[63:0]
         }
     }
     storeOffset++
  }
}
```
SIMD Floating-Point Exceptions

None.

Memory Down-conversion: D_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_mask_extpackstorehi_pd (void*, __mmask8, __m512d, _MM_DOWNCONV_PD_ENUM, int);
- void _mm512_packstorehi_pd (void*, _m512d);
- void _mm512_mask_packstorehi_pd (void*, _mmask8, _m512d);

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

VPACKSTOREHPS - Pack And Store Unaligned High From Float32 Vector

Description

Packs and down-converts the mask-enabled elements of float 32 vector zmm1 into a byte/word/doubleword stream logically mapped starting at element-aligned address (*m^t −* 64), and stores the high-64-byte elements of that stream (those elements of the stream that map at or after the first 64-byte-aligned address following $(m_t - 64)$, the high cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorelps instruction is used to store the part of the stream before the first 64byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFFFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorelps, this instruction is useful for packing data into a queue. Also in conjunction with vpackstorelps, it allows unaligned vector stores (that is, vector stores that are only element-wise , not vector-wise, aligned); just use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstorelps [rax] {k1}, v0
vpackstorehps [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp 8^*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or #SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.

intel

```
storeOffset = 0
downSize = DownConvStoreSizeOff32(SSS[2:0])
foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 16; n++) {
   if(k1[n] != 0) {
     if (foundNext64BytesBoundary == false) {
         if ( ( (pointer + (storeOffset+1)*downSize) % 64) == 0 ) {
            foundNext64BytesBoundary = true
         }
     } else {
         i = 32*ntmp = DownConvStoref32(zmm1[i+31:i], SSS[2:0])
         if(downSize == 4) {
          MemStore(pointer + storeOffset*4) = tmp[31:0]
         } else if(downSize == 2) {
          MemStore(pointer + storeOffset*2) = tmp[15:0]
         } else if(downSize == 1) {
           MemStore(pointer + storeOffset) = tmp[7:0]
         }
     }
     storeOffset++
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Memory Down-conversion: D_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_mask_extpackstorehi_ps (void*, _mmask16, _m512, _MM_DOWNCONV_PS_ENUM, int);
- void _mm512_packstorehi_ps (void*, _m512);
- void _mm512_mask_packstorehi_ps (void*, _mmask16, _m512);

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

64 bit Mode

intel

VPACKSTOREHQ - Pack And Store Unaligned High From Int64 Vector

Description

Packs and down-converts the mask-enabled elements of int64 vector zmm1 into a int64 stream logically mapped starting at element-aligned address (*m^t −* 64), and stores the high-64-byte elements of that stream (those elements of the stream that map at or after the first 64-byte-aligned address following $(m_t - 64)$, the high cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorelg instruction is used to store the part of the stream before the first 64byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorelq, this instruction is useful for packing data into a queue. Also in conjunction with vpackstorelq, it allows unaligned vector stores (that is, vector stores that are only element-wise , not vector-wise, aligned); just use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstorelq [rax] {k1}, v0
vpackstorehq [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp 8^*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand. The instruction will not produce any #GP or #SS fault due to address canonicity nor #PF fault if the address is aligned to a 64-byte boundary. Additionally, A/D bits in the page table will not be updated.


```
storeOffset = 0downSize = DownConvStoreSizeOfi64(SSS[2:0])
foundNext64BytesBoundary = false
pointer = m_t - 64
for (n = 0; n < 8; n++) {
   if(k1[n] != 0) {if (foundNext64BytesBoundary == false) {
         if ( ( (pointer + (storeOffset+1)*downSize) % 64) == 0 ) {
           foundNext64BytesBoundary = true
         }
     } else {
         i = 64*ntmp = DownConvStorei64(zmm1[i+63:i], SSS[2:0])
         if(downSize == 8) {
          MemStore(pointer + storeOffset*8) = tmp[63:0]
         }
     }
     storeOffset++
  }
}
```
Flags Affected

None.

Memory Down-conversion: D*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

VPACKSTORELD - Pack and Store Unaligned Low From Int32 Vector

Description

Packs and down-converts the mask-enabled elements of int32 vector zmm1 into a byte/word/doubleword stream logically mapped starting at element-aligned address *mt*, and stores the low-64 byte elements of that stream (those elements of the stream that map before the first 64byte-aligned address following *mt*, the low cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorehd instruction is used to store the part of the stream at or after the first 64-byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFFFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorehd, this instruction is useful for packing data into into a queue. Also in conjunction with vpackstorehd, it allows unaligned vector stores (that is, vector stores that are only element-wise, not vector-wise, aligned); just use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstoreld [rax] {k1}, v0
vpackstorehd [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.

intel

```
storeOffset = 0
downSize = DownConvStoreSizeOfi32(SSS[2:0])
for(n = 0; n < 16; n++) {
 if (k1[n] != 0) {
   i = 32*ntmp = DownConvStorei32(zmm1[i+31:i], SSS[2:0])
   if(downSize == 4) {
     MemStore(mt+4*storeOffset) = tmp[31:0]
   } else if(downSize == 2) {
     MemStore(m_t+2*storeOffset) = tmp[15:0]} else if(downSize == 1) {
     MemStore(m_t+storeOffSet) = tmp[7:0]}
   storeOffset++
    if (((m_t + downSize*storeOffset) % 64) == 0) {
     break
   }
 }
}
```
Flags Affected

None.

Memory Down-conversion: D*i*³²

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

64 bit Mode

intel

VPACKSTORELPD - Pack and Store Unaligned Low From Float64 Vector

Description

Packs and down-converts the mask-enabled elements of float64 vector zmm1 into a float64 stream logically mapped starting at element-aligned address m_t , and stores the low-64-byte elements of that stream (those elements of the stream that map before the first 64-byte-aligned address following m_t , the low cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorehpd instruction is used to store the part of the stream at or after the first 64-byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorehpd, this instruction is useful for packing data into into a queue. Also in conjunction with vpackstorehpd, it allows unaligned vector stores (that is, vector stores that are only element-wise, not vector-wise, aligned); just use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstorelpd [rax] {k1}, v0
vpackstorehpd [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.


```
storeOffset = 0
downSize = DownConvStoreSizeOf_{f64}(SSS[2:0])for(n = 0; n < 8; n^{++}) {
  if (k1[n] != 0) {
    i = 64*ntmp = DownConvStore_{f64}(zmm1[i+63:i], SSS[2:0])if(downSize == 8) {
      MemStore(mt+8*storeOffset) = tmp[63:0]
    }
    storeOffset++
    if (((m_t + downSize*storeOffset) % 64) == 0) {
      break
    }
 }
}
```
SIMD Floating-Point Exceptions

None.

Memory Down-conversion: D_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_extpackstorelo_pd (void*, _m512d, _MM_DOWNCONV_PD_ENUM, int);
- void _mm512_mask_extpackstorelo_pd (void*, __mmask8, __m512d, _MM_DOWNCONV_PD_ENUM, int);
- void _mm512_packstorelo_pd (void*, __m512d);
- void _mm512_mask_packstorelo_pd (void*, __mmask8, __m512d);

Exceptions

(intel)

Real-Address Mode and Virtual-8086

VPACKSTORELPS - Pack and Store Unaligned Low From Float32 Vector

Description

Packs and down-converts the mask-enabled elements of float 32 vector zmm1 into a byte/word/doubleword stream logically mapped starting at element-aligned address *mt*, and stores the low-64-byte elements of that stream (those elements of the stream that map before the first 64-byte-aligned address following m_t , the low cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorehps instruction is used to store the part of the stream at or after the first 64-byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFFFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorehps, this instruction is useful for packing data into into a queue. Also in conjunction with vpackstorehps, it allows unaligned vector stores (that is, vector stores that are only element-wise, not vector-wise, aligned); just use a mask of 0xFFFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstorelps [rax] {k1}, v0
vpackstorehps [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.

intel

```
storeOffset = 0
downSize = DownConvStoreSizeOff32(SSS[2:0])
for(n = 0 ;n < 16; n++) {
  if (k1[n] := 0) {
   i = 32*ntmp = DownConvStoref32(zmm1[i+31:i], SSS[2:0])
    if(downSize == 4) {
      MemStore(m_t+4*storeOffset) = tmp[31:0]} else if(downSize == 2) {
      MemStore(mt+2*storeOffset) = tmp[15:0]
    } else if(downSize == 1) {
      MemStore(m_t+storeOffSet) = tmp[7:0]}
    storeOffset++
    if (((m_t + downSize*storeOffset) % 64) == 0) {
      break
    }
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Memory Down-conversion: D_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_extpackstorelo_ps (void*, __m512, _MM_DOWNCONV_PS_ENUM, int);
- void _mm512_mask_extpackstorelo_ps (void*, __mmask16, __m512, _MM_DOWNCONV_PS_ENUM, int);
- void _mm512_packstorelo_ps (void*, _m512);
- void _mm512_mask_packstorelo_ps (void*, __mmask16, __m512);

Exceptions

Real-Address Mode and Virtual-8086

VPACKSTORELQ - Pack and Store Unaligned Low From Int64 Vector

Description

Packs and down-converts the mask-enabled elements of int64 vector zmm1 into a int64 stream logically mapped starting at element-aligned address *mt*, and stores the low-64 byte elements of that stream (those elements of the stream that map before the first 64byte-aligned address following *mt*, the low cache line in the current implementation). The length of the stream depends on the number of enabled masks, as elements disabled by the mask are not added to the stream.

The vpackstorehq instruction is used to store the part of the stream at or after the first 64-byte-aligned address preceding *mt*.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector, choosing which elements are added to the stream. The one similarity to a write-mask as used in the rest of this document is that the no-write-mask option (encoding 0) is available to select a mask of 0xFF for this instruction. For that reason, the notation and encoding are the same as for a write-mask.

In conjunction with vpackstorehq, this instruction is useful for packing data into into a queue. Also in conjunction with vpackstorehq, it allows unaligned vector stores (that is, vector stores that are only element-wise, not vector-wise, aligned); just use a mask of 0xFF or no write-mask for this purpose. The typical instruction sequence to perform an unaligned vector store would be:

```
// assume memory location is pointed by register rax
vpackstorelq [rax] {k1}, v0
vpackstorehq [rax+64] {k1}, v0
```
This instruction does not have subset support.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note that the address reported by a page fault is the beggining of the 64-byte cache line boundary containing the memory operand.


```
storeOffset = 0
downSize = DownConvStoreSizeOfi64(SSS[2:0])
for(n = 0; n < 8; n^{++}) {
  if (k1[n] != 0) {
    i = 64*ntmp = DownConvStorei64(zmm1[i+63:i], SSS[2:0])
    if(downSize == 8) {
      MemStore(m_t+8*storeOffset) = tmp[63:0]}
    storeOffset++
    if (((m_t + downSize*storeOffset) % 64) == 0) {
      break
    }
 }
}
```
Flags Affected

None.

Memory Down-conversion: D*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm512_mask_packstorelo_epi64 (void*, _mmask8, _m512i);

Exceptions

(intel)

Real-Address Mode and Virtual-8086

VPADCD - Add Int32 Vectors with Carry

Description

Performs an element-by-element three-input addition between int32 vector zmm1, the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3, and the corresponding bit of k2. The result is written into int32 vector zmm1.

In addition, the carry from the sum for the n-th element is written into the n-th bit of vector mask k2.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k2. Elements in zmm1 and k2 with the corresponding bit clear in k1 retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
\mathbf{r}for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    tmpCarry = Carry(zmm1[i+31:i] + k2[n] + tmpSrc3[i+31:i])zmm1[i+31:i] = zmm1[i+31:i] + k2[n] + tmpSrc3[i+31:i]k2[n] = tmpCarry}
}
```
Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

CHAPTER 6. INSTRUCTION DESCRIPTIONS

64 bit Mode #SS(0) If a memory address referencing the SS segment is in a non-canonical form. #GP(0) If the memory address is in a non-canonical form. If a memory operand linear address is not aligned to the data size granularity dictated by SwizzUpConv mode. #PF(fault-code) For a page fault. #NM If CR0.TS[bit 3]=1. #UD If processor model does not implement the specific instruction. If preceded by any REX, F0, F2, F3, or 66 prefixes.

VPADDD - Add Int32 Vectors

Description

Performs an element-by-element addition between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// integer operation
    zmm1[i+31:i] = zmm2[i+31:i] + tmpSrc3[i+31:i]
  }
}
```
Flags Affected

None.

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VPADDSETCD - Add Int32 Vectors and Set Mask to Carry

Description

Performs an element-by-element addition between int32 vector zmm1 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

In addition, the carry from the sum for the n-th element is written into the n-th bit of vector mask k2.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k2. Elements in zmm1 and k2 with the corresponding bit clear in k1 retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    k2[n] = Carry(zmm1[i+31:i] + tmpSrc3[i+31:i])zmm1[i+31:i] = zmm1[i+31:i] + tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode #SS(0) If a memory address referencing the SS segment is in a non-canonical form. #GP(0) If the memory address is in a non-canonical form. If a memory operand linear address is not aligned to the data size granularity dictated by SwizzUpConv mode. #PF(fault-code) For a page fault. #NM If CR0.TS[bit 3]=1. #UD If processor model does not implement the specific instruction. If preceded by any REX, F0, F2, F3, or 66 prefixes.

VPADDSETSD - Add Int32 Vectors and Set Mask to Sign

Description

Performs an element-by-element addition between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

In addition, the sign of the result for the n-th element is written into the n-th bit of vector mask k1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    zmm1[i+31:i] = zmm2[i+31:i] + tmpSrc3[i+31:i]
    k1[n] = zmm1[i+31]}
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_addsets_epi32 (__m512i, __mmask16, __m512i, __m512i, $_mmask16^*$);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VPANDD - Bitwise AND Int32 Vectors

Description

Performs an element-by-element bitwise AND between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] = zmm2[i+31:i] & tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_and_epi32(__m512i, __m512i); __m512i _mm512_mask_and_epi32(__m512i, __mmask16,__m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPANDND - Bitwise AND NOT Int32 Vectors

Description

Performs an element-by-element bitwise AND between NOT int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] = (~(zmm2[i+31:i])) & tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_andnot_epi32 (__m512i, __m512i); __m512i _mm512_mask_andnot_epi32 (__m512i, __mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPANDNQ - Bitwise AND NOT Int64 Vectors

Description

Performs an element-by-element bitwise AND between NOT int64 vector zmm2 and the int64 vector result of the swizzle/broadcast/conversion process on memory or int64 vector zmm3. The result is written into int64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nzmm1[i+63:i] = (~(zmm2[i+63:i])) & tmpSrc3[i+63:i]
  }
}
```
Flags Affected

Memory Up-conversion: S*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPANDQ - Bitwise AND Int64 Vectors

Description

Performs an element-by-element bitwise AND between int64 vector zmm2 and the int64 vector result of the swizzle/broadcast/conversion process on memory or int64 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nzmm1[i+63:i] = zmm2[i+63:i] & tmpSrc3[i+63:i]
  }
}
```
Flags Affected

Memory Up-conversion: S*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPBLENDMD - Blend Int32 Vectors using the Instruction Mask

Description

Performs an element-by-element blending between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3, using the instruction mask as selector. The result is written into int32 vector zmm1.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector: every element of the destination is conditionally selected between first source or second source using the value of the related mask bit (0 for first source, 1 for second source).

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = tmpSrc3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(tmpSrc3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
   if(k1[n] ==1 or *no write-mask*) {
      zmm1[i+31:i] = tmpSrc3[i+31:i]
   } else {
      zmm1[i+31:i] = zmm2[i+31:i]}
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_blend_epi32 (__mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

(intel)

CHAPTER 6. INSTRUCTION DESCRIPTIONS

VPBLENDMQ - Blend Int64 Vectors using the Instruction Mask

Description

Performs an element-by-element blending between int64 vector zmm2 and the int64 vector result of the swizzle/broadcast/conversion process on memory or int64 vector zmm3, using the instruction mask as selector. The result is written into int64 vector zmm1.

The mask is not used as a write-mask for this instruction. Instead, the mask is used as an element selector: every element of the destination is conditionally selected between first source or second source using the value of the related mask bit (0 for first source, 1 for second source).

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = tmpSrc3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i64</sub>(tmpSrc3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
   if(k1[n] ==1 or *no write-mask*) {
      zmm1[i+63:i] = tmpSrc3[i+63:i]
   } else {
      zmm1[i+63:i] = zmm2[i+63:i]}
}
```
Flags Affected

Memory Up-conversion: S*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_mask_blend_epi64 (__mmask8, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPBROADCASTD - Broadcast Int32 Vector

Description

The 1, 2, or 4 bytes (depending on the conversion and broadcast in effect) at memory address *m^t* are broadcast and/or converted to a int32 vector. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {1to16}
tmpSrc2[31:0] = UpConvLoad_{i32}(m_t)for (n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    i = 32*nzmm1[i+31:i] = tmpSrc2[31:0]
 }
}
```
Flags Affected

None.

Memory Up-conversion: U*i*³²

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Protected and Compatibility Mode

VPBROADCASTQ - Broadcast Int64 Vector

Description

The 8 bytes at memory address *m^t* are broadcast to a int64 vector. The result is written into int64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
// {1to8}
tmpSrc2[63:0] = UpConvLoad<sub>i64</sub>(m<sub>t</sub>)for (n = 0; n < 8; n++) {
  if (k1[n] != 0) {
    i = 64*nzmm1[i+63:i] = tmpSrc2[63:0]
  }
}
```
Flags Affected

None.

Memory Up-conversion: U*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

VPCMPD - Compare Int32 Vectors and Set Vector Mask

Description

Performs an element-by-element comparison between int32 vector zmm1 and the int32 vector result of the swizzle/broadcast/conversion from memory or int32 vector zmm2. The result is written into vector mask k2.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

Immediate Format

Operation

```
switch (IMM8[2:0]) {
   case 0: OP ← EQ; break;
   case 1: OP ← LT; break;
   case 2: OP ← LE; break;
   case 4: OP ← NEQ; break;
   case 5: OP ← NLT; break;
```


```
case 6: OP ← NLE; break;
    default: Reserved; break;
}
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]
} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  k2[n] = 0if(k1[n] != 0) {i = 32*n
    // signed integer operation
    k2[n] = (zmm1[i+31:i] OP tmpSrc2[i+31:i]) ? 1 : 0
  }
}
```
Instruction Pseudo-ops

Compilers and assemblers may implement the following pseudo-ops in addition to the standard instruction op:

Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPCMPEQD - Compare Equal Int32 Vectors and Set Vector Mask

Description

Performs an element-by-element compare for equality between int32 vector zmm1 and the int32 vector result of the swizzle/broadcast/conversion from memory or int32 vector zmm2. The result is written into vector mask k2.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  k2[n] = 0if(k1[n] != 0) {
    i = 32*n// signed integer operation
    k2[n] = (zmm1[i+31:i] == tmpSrc2[i+31:i]) ? 1 : 0}
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_cmpeq_epi32_mask (__m512i, __m512i); __mmask16 _mm512_mask_cmpeq_epi32_mask (__mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VPCMPGTD - Compare Greater Than Int32 Vectors and Set Vector Mask

Description

Performs an element-by-element compare for the greater value of int32 vector zmm1 and the int32 vector result of the swizzle/broadcast/conversion from memory or int32 vector zmm2. The result is written into vector mask k2.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  k2[n] = 0if(k1[n] != 0) {
    i = 32*n
    // signed integer operation
    k2[n] = (zmm1[i+31:i] > tmpSrc2[i+31:i]) ? 1 : 0}
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_cmpgt_epi32_mask (__m512i, __m512i); __mmask16 _mm512_mask_cmpgt_epi32_mask (__mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VPCMPLTD - Compare Less Than Int32 Vectors and Set Vector Mask

Description

Performs an element-by-element compare for the lesser value of int32 vector zmm1 and the int32 vector result of the swizzle/broadcast/conversion from memory or int32 vector zmm2. The result is written into vector mask k2.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  k2[n] = 0if(k1[n] != 0) {
    i = 32*n// signed integer operation
    k2[n] = (zmm1[i+31:i] < tmpSrc2[i+31:i]) ? 1 : 0}
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_cmplt_epi32_mask (__m512i, __m512i); __mmask16 _mm512_mask_cmplt_epi32_mask (__mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

 64

VPCMPUD - Compare Uint32 Vectors and Set Vector Mask

Description

Performs an element-by-element comparison between uint32 vector zmm1 and the uint32 vector result of the swizzle/broadcast/conversion from memory or uint32 vector zmm2. The result is written into vector mask k2.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

Immediate Format


```
switch (IMM8[2:0]) {
   case 0: OP ← EQ; break;
   case 1: OP ← LT; break;
   case 2: OP ← LE; break;
   case 4: OP ← NEQ; break;
```


```
case 5: OP ← NLT; break;
    case 6: OP ← NLE; break;
   default: Reserved; break;
}
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoadi32(zmm2/mt)
}
for (n = 0; n < 16; n++) {
 k2[n] = 0if(k1[n] != 0) {i = 32*n// unsigned integer operation
   k2[n] = (zmm1[i+31:i] OP tmpSrc2[i+31:i]) ? 1 : 0
  }
}
```
Instruction Pseudo-ops

Compilers and assemblers may implement the following pseudo-ops in addition to the standard instruction op:

Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode #SS(0) If a memory address referencing the SS segment is in a non-canonical form. #GP(0) If the memory address is in a non-canonical form. If a memory operand linear address is not aligned to the data size granularity dictated by SwizzUpConv mode. #PF(fault-code) For a page fault. #NM If CR0.TS[bit 3]=1. If preceded by any REX, F0, F2, F3, or 66 prefixes.

VPERMD - Permutes Int32 Vectors

Description

Performs an element permutation of elements from int32 vector read from memory or vector zmm3, using int32 vector zmm2 element as source indices. The result of the permutation is written into int32 vector zmm1. Note that index values may be repeated so that the same vector source element can be replicated across many vector destination elements.

No swizzle, broadcast, or conversion is performed by this instruction.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
src[511:0] = zmm3/mt
for (n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    i = 32*nj = zmm2[i+3:i]*32zmm1[i+31:i] = src[j+31:j]}
}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

VPERMF32X4 - Shuffle Vector Dqwords

Description

Shuffles 128-bit blocks of the vector read from memory or vector zmm2/mem using index bits in immediate. The result of the shuffle is written into vector zmm1.

No swizzle, broadcast, or conversion is performed by this instruction.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format


```
src[511:0] = zmm2/m_t// Inter-lane shuffle
for (n = 0; n < 16/4; n++) {
  i = 128*n
  j = 128 * ((perm128 \gg 2*n) & 0x3)tmp[i+127:i] = src[j+127:j]}
// Writemasking
for (n = 0; n < 16; n++) {
  if (k1[n] != 0) {
    zmm1[i+31:i] = tmp[i+31:i]
  }
}
```


Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

VPGATHERDD - Gather Int32 Vector With Signed Dword Indices

Description

A set of 16 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* are converted to a int32 vector. The result is written into int32 vector zmm1.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Note that accessed element by will always access 64 bytes of memory. The memory region accessed by each element will always be between *elemen_linear_address & (∼0x3F)* and *(element_linear_address & (∼0x3F)) + 63* boundaries.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully loaded.

The instruction will #GP fault if the destination vector zmm1 is the same as index vector *V INDEX*.

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
 if ( ktemp[n] != 0) {
```


```
i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
   pointer[63:0] = mv_t[n]zmm1[i+31:i] = UpConvLoadi32(pointer)
   k1[n] = 0}
}
```
Flags Affected

None.

Memory Up-conversion: U_{i32}

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VPGATHERDQ - Gather Int64 Vector With Signed Dword Indices

Description

A set of 8 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* are converted to a int64 vector. The result is written into int64 vector zmm1.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Note that accessed element by will always access 64 bytes of memory. The memory region accessed by each element will always be between *elemen_linear_address & (∼0x3F)* and *(element_linear_address & (∼0x3F)) + 63* boundaries.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully loaded.

The instruction will #GP fault if the destination vector zmm1 is the same as index vector *V INDEX*.

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 8; n++) {
 if ( ktemp[n] != 0) {
```


```
i = 64*nj = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[j+31:j] * SCALE)
    pointer[63:0] = mv_t[n]zmm1[i+63:i] = UpConvLoadi64(pointer)
   k1[n] = 0}
}
k1[15:8] = 0
```
Flags Affected

None.

Memory Up-conversion: U*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512i _mm512_i32logather_epi64 (__m512i, void const*, int);
- __m512i _mm512_mask_i32logather_epi64 (__m512i, __mmask8, __m512i, void const*, int);
- __m512i _mm512_i32loextgather_epi64 (__m512i, void const*, _MM_UPCONV_EPI64_ENUM, int, int);
- __m512i _mm512_mask_i32loextgather_epi64 (__m512i, __mmask8, __m512i, void const*, _MM_UPCONV_EPI64_ENUM, int, int);

Exceptions

Real-Address Mode and Virtual-8086

-
- #UD Instruction not available in these modes

VPMADD231D - Multiply First Source By Second Source and Add To Destination Int32 Vectors

Description

Performs an element-by-element multiplication between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or vector int32 zmm3, then adds the result to int32 vector zmm1. The final sum is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    zmm1[i+31:i] = zmm2[i+31:i] * tmpSrc3[i+31:i] + zmm1[i+31:i]
  }
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512i _mm512_fmadd_epi32 (__m512i, __m512i, __m512i);
- __m512i _mm512_mask_fmadd_epi32 (__m512i, __mmask16, __m512i, __m512i);
- __m512i _mm512_mask3_fmadd_epi32 (__m512i, __m512i, __m512i, __mmask16);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

 64

Protected and Compatibility Mode

VPMADD233D - Multiply First Source By Specially Swizzled Second Source and Add To Second Source Int32 Vectors

Description

This instruction is built around the concept of 4-element sets, of which there are four: elements 0-3, 4-7, 8-11, and 12-15. If we refer to the int32 vector result of the broadcast (no conversion is supported) process on memory or the int32 vector zmm3 (no swizzle is supported) as t3, then:

Each element 0-3 of int32 vector zmm2 is multiplied by element 1 of t3, the result is added to element 0 of t3, and the final sum is written into the corresponding element $0-3$ of int32 vector zmm1.

Each element 4-7 of int32 vector zmm2 is multiplied by element 5 of t3, the result is added to element 4 of t3, and the final sum is written into the corresponding element 4-7 of int32 vector zmm1.

Each element 8-11 of int32 vector zmm2 is multiplied by element 9 of t3, the result is added to element 8 of t 3 , and the final sum is written into the corresponding element 8-11 of int32 vector zmm1.

Each element 12-15 of int32 vector zmm2 is multiplied by element 13 of t3, the result is added to element 12 of t3, and the final sum is written into the corresponding element 12-15 of int32 vector zmm1.

This instruction makes it possible to perform scale and bias in a single instruction without needing to have either scale or bias already loaded in a register. This saves one vector load for each interpolant, representing around ten percent of shader instructions.

For structure-of-arrays (SOA) operation, this instruction is intended to be used with the *{*4*to*16*}* broadcast on src2, allowing all 16 scale and biases to be identical. For array-ofstructures (AOS) vec4 operations, no broadcast is used, allowing four different scales and biases, one for each vec4.

No conversion or swizzling is supported for this instruction. However, all broadcasts except *{*1*to*16*}* are supported (i.e. 16to16 and 4to16).

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
 if (k1[n] := 0) {
   i = 32*nbase = (n & 0x03) * 32scale[31:0] = tmpSrc3[base+63:base+32]bias[31:0] = tmpSrc3[base+31:base]// integer operation
    zmm1[i+31:i] = zmm2[i+31:i] * scale[31:0] + bias[31:0]
 }
}
```
Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

an Invalid Opcode fault is raised

VPMAXSD - Maximum of Int32 Vectors

Description

Determines the maximum value of each pair of corresponding elements in int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad_{i32}(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    zmm1[i+31:i] = IMax(zmm2[i+31:i] , tmpSrc3[i+31:i])
 }
}
```
Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VPMAXUD - Maximum of Uint32 Vectors

Description

Determines the maximum value of each pair of corresponding elements in uint32 vector zmm2 and the uint32 vector result of the swizzle/broadcast/conversion process on memory or uint32 vector zmm3. The result is written into uint32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad_{i32}(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// unsigned integer operation
    zmm1[i+31:i] = UMax(zmm2[i+31:i] , tmpSrc3[i+31:i])
 }
}
```
Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VPMINSD - Minimum of Int32 Vectors

Description

Determines the minimum value of each pair of corresponding elements in int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// signed integer operation
    zmm1[i+31:i] = (zmm2[i+31:i] < tmpSrc3[i+31:i]) ?
              zmm2[i+31:i] : tmpSrc3[i+31:i]
 }
}
```
Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

VPMINUD - Minimum of Uint32 Vectors

Description

Determines the minimum value of each pair of corresponding elements in uint32 vector zmm2 and the uint32 vector result of the swizzle/broadcast/conversion process on memory or uint32 vector zmm3. The result is written into uint32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// unsigned integer operation
    zmm1[i+31:i] = UMin(zmm2[i+31:i] , tmpSrc3[i+31:i])
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPMULHD - Multiply Int32 Vectors And Store High Result

Description

Performs an element-by-element multiplication between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The high 32 bits of the result are written into int32 zmm1 vector.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    tmp[63:0] = zmm2[i+31:i] * tmpSrc3[i+31:i]zmm1[i+31:i] = tmp[63:32]}
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPMULHUD - Multiply Uint32 Vectors And Store High Result

Description

Performs an element-by-element multiplication between uint32 vector zmm2 and the uint32 vector result of the swizzle/broadcast/conversion process on memory or uint32 vector zmm3. The high 32 bits of the result are written into uint32 zmm1 vector.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// unsigned integer operation
    tmp[63:0] = zmm2[i+31:i] * tmpSrc3[i+31:i]zmm1[i+31:i] = tmp[63:32]}
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPMULLD - Multiply Int32 Vectors And Store Low Result

Description

Performs an element-by-element multiplication between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3, and the low 32 bits of the result are written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// signed integer operation
    zmm1[i+31:i] = zmm2[i+31:i] * tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPORD - Bitwise OR Int32 Vectors

Description

Performs an element-by-element bitwise OR between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*nzmm1[i+31:i] = zmm2[i+31:i] | tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_or_epi32 (__m512i, __m512i); __m512i _mm512_mask_or_epi32 (__m512i, __mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPORQ - Bitwise OR Int64 Vectors

Description

Performs an element-by-element bitwise OR between int64 vector zmm2 and the int64 vector result of the swizzle/broadcast/conversion process on memory or int64 vector zmm3. The result is written into int64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nzmm1[i+63:i] = zmm2[i+63:i] | tmpSrc3[i+63:i]
  }
}
```
Flags Affected

Memory Up-conversion: S*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512i _mm512_or_epi64 (__m512i, __m512i); __m512i _mm512_mask_or_epi64 (__m512i, __mmask8, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSBBD - Subtract Int32 Vectors with Borrow

Description

Performs an element-by-element three-input subtraction of the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3, as well as the corresponding bit of k2, from int32 vector zmm1. The result is written into int32 vector zmm1.

In addition, the borrow from the subtraction difference for the n-th element is written into the n-th bit of vector mask k2.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k2. Elements in zmm1 and k2 with the corresponding bit clear in k1 retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// integer operation
    tmpBorrow = Borrow(zmm1[i+31:i] - k2[n] - tmpSrc3[i+31:i])zmm1[i+31:i] = zmm1[i+31:i] - k2[n] - tmpSrc3[i+31:i]k2[n] = tmpBorrow}
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VPSBBRD - Reverse Subtract Int32 Vectors with Borrow

Description

Performs an element-by-element three-input subtraction of int32 vector zmm1, as well as the corresponding bit of k2, from the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

In addition, the borrow from the subtraction for the n-th element is written into the n-th bit of vector mask k2.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k2. Elements in zmm1 and k2 with the corresponding bit clear in k1 retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
\mathbf{r}for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    tmpBorrow = Borrow(tmpSrc3[i+31:i] - k2[n] - zmm1[i+31:i])zmm1[i+31:i] = tmpSrc3[i+31:i] - k2[n] - zmm1[i+31:i]k2[n] = tmpBorrow}
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode #SS(0) If a memory address referencing the SS segment is in a non-canonical form. #GP(0) If the memory address is in a non-canonical form. If a memory operand linear address is not aligned to the data size granularity dictated by SwizzUpConv mode. #PF(fault-code) For a page fault. #NM If CR0.TS[bit 3]=1. #UD If processor model does not implement the specific instruction. If preceded by any REX, F0, F2, F3, or 66 prefixes.

VPSCATTERDD - Scatter Int32 Vector With Signed Dword Indices

Description

Down-converts and stores all 16 elements in int32 vector UNDEF to the memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX*, with scale *SCALE*.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Writes to overlapping destination memory locations are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Only writes to overlapping vector indices are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Writes that are not overlapped may happen in any order. Memory ordering with other instructions follows the Intel-64 memory ordering model. Note that this does not account for non-overlapping indices that map into the same physical address locations.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully stored.

Operation

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
```


```
if (ktemp[n] != 0) {
    i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]tmp = DownConvStorei32(UNDEF[i+31:i], SSS[2:0])
    if(DownConvStoreSizeOfi32(SSS[2:0]) == 4) {
     MemStore(pointer) = tmp[31:0]
    } else if(DownConvStoreSizeOfi32(SSS[2:0]) == 2) {
     MemStore(pointer) = tmp[15:0]
    } else if(DownConvStoreSizeOfi32(SSS[2:0]) == 1) {
     MemStore(pointer) = tmp[7:0]
   }
   k1[n] = 0}
}
```
Flags Affected

None.

Memory Down-conversion: D*i*³²

Intel® C/C++ Compiler Intrinsic Equivalent

```
void _mm512_i32scatter_epi32 (void*, _m512i, _m512i, int);
void _mm512_mask_i32scatter_epi32 (void*, _mmask16, _m512i, _m512i, int);
```
Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

64

Protected and Compatibility Mode

VPSCATTERDQ - Scatter Int64 Vector With Signed Dword Indices

Description

Down-converts and stores all 8 elements in int64 vector UNDEF to the memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX*, with scale *SCALE*.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Writes to overlapping destination memory locations are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Only writes to overlapping vector indices are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Writes that are not overlapped may happen in any order. Memory ordering with other instructions follows the Intel-64 memory ordering model. Note that this does not account for non-overlapping indices that map into the same physical address locations.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully stored.

Operation

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 8; n++) {
```
intel

```
if (ktemp[n] != 0) {
    i = 64*nj = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[j+31:j] * SCALE)
    pointer[63:0] = mv_t[n]tmp = DownConvStorei64(UNDEF[i+63:i], SSS[2:0])
    if(DownConvStoreSizeOf<sub>i64</sub>(SSS[2:0]) == 8) {
      MemStore(pointer) = tmp[63:0]
    }
    k1[n] = 0}
k1[15:8] = 0
```
Flags Affected

}

None.

Memory Down-conversion: D*i*⁶⁴

Intel® C/C++ Compiler Intrinsic Equivalent

```
void _mm512_i32loscatter_epi64 (void*, _m512i, _m512i, int);
```
void _mm512_mask_i32loscatter_epi64 (void*, _mmask8, _m512i, _m512i, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

(intel)

VPSHUFD - Shuffle Vector Doublewords

Description

Shuffles 32 bit blocks of the vector read from memory or vector zmm2/mem using index bits in immediate. The result of the shuffle is written into vector zmm1.

No swizzle, broadcast, or conversion is performed by this instruction.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

Operation

```
src[511:0] = zmm2/mt
// Intra-lane shuffle
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*n// offset within 128-bit chunk
    j = 32*((perm32 >> 2*(n & 0x3)) & 0x3)
    // 128-bit level offset
    j = j + 128*(n \gg 2)zmm1[i+31:i] = src[j+31:j]}
}
```


Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

VPSLLD - Shift Int32 Vector Immediate Left Logical

Description

Performs an element-by-element logical left shift of the result of the swizzle/broadcast/conversion process on memory or vector int32 zmm2, shifting by the number of bits specified in immediate field. The result is stored in int32 vector zmm1.

If the value specified by the shift operand is greater than 31 then the destination operand is set to all 0s.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoadi32(zmm2/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    zmm1[i+31:i] = tmpSrc2[i+31:i] << IMM8[7:0]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSLLVD - Shift Int32 Vector Left Logical

Description

Performs an element-by-element left shift of int32 vector zmm2, shifting by the number of bits specified by the int32 vector result of the swizzle/broadcast/conversion process on memory or vector int32 zmm3. The result is stored in int32 vector zmm1.

If the value specified by the shift operand is greater than 31 then the destination operand is set to all 0s.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    zmm1[i+31:i] = zmm2[i+31:i] << tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSRAD - Shift Int32 Vector Immediate Right Arithmetic

Description

Performs an element-by-element arithmetic right shift of the result of the swizzle/broadcast/conversion process on memory or vector int32 zmm2, shifting by the number of bits specified in immediate field. The result is stored in int32 vector zmm1.

An arithmetic right shift leaves the sign bit unchanged after each shift count, so the final result has the *i*+1 msbs set to the original sign bit, where *i* is the number of bits by which to shift right.

If the value specified by the shift operand is greater than 31 each destination data element is filled with the initial value of the sign bit of the element.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrcc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoadi32(zmm2/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    zmm1[i+31:i] = tmpSrc2[i+31:i] >> IMM8[7:0]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSRAVD - Shift Int32 Vector Right Arithmetic

Description

Performs an element-by-element arithmetic right shift of int32 vector zmm2, shifting by the number of bits specified by the int32 vector result of the swizzle/broadcast/conversion process on memory or vector int32 zmm3. The result is stored in int32 vector zmm1.

An arithmetic right shift leaves the sign bit unchanged after each shift count, so the final result has the *i*+1 msbs set to the original sign bit, where *i* is the number of bits by which to shift right.

If the value specified by the shift operand is greater than 31 each destination data element is filled with the initial value of the sign bit of the element.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// signed integer operation
    zmm1[i+31:i] = zmm2[i+31:i] >> tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSRLD - Shift Int32 Vector Immediate Right Logical

Description

Performs an element-by-element logical right shift of the result of the swizzle/broadcast/conversion process on memory or vector int32 zmm2, shifting by the number of bits specified in immediate field. The result is stored in int32 vector zmm1.

A logical right shift shifts a 0-bit into the msb for each shift count, so the final result has the *i* msbs set to 0, where *i* is the number of bits by which to shift right.

If the value specified by the shift operand is greater than 31 then the destination operand is set to all 0s.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    zmm1[i+31:i] = tmpSrc2[i+31:i] >> IMMS[7:0]}
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSRLVD - Shift Int32 Vector Right Logical

Description

Performs an element-by-element logical right shift of int32 vector zmm2, shifting by the number of bits specified by the int32 vector result of the swizzle/broadcast/conversion process on memory or vector int32 zmm3. The result is stored in int32 vector zmm1.

A logical right shift shifts a 0-bit into the msb for each shift count, so the final result has the *i* msbs set to 0, where *i* is the number of bits by which to shift right.

If the value specified by the shift operand is greater than 31 then the destination operand is set to all 0s.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// signed integer operation
    zmm1[i+31:i] = zmm2[i+31:i] >> tmpSrc3[i+31:i]}
}
```
Flags Affected

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSUBD - Subtract Int32 Vectors

Description

Performs an element-by-element subtraction from int32 vector zmm2 of the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// integer operation
    zmm1[i+31:i] = zmm2[i+31:i] - tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSUBRD - Reverse Subtract Int32 Vectors

Description

Performs an element-by-element subtraction of int32 vector zmm2 from the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// integer operation
    zmm1[i+31:i] = -zmm2[i+31:i] + tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSUBRSETBD - Reverse Subtract Int32 Vectors and Set Borrow

Description

Performs an element-by-element subtraction of int32 vector zmm1 from the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

In addition, the borrow from the subtraction for the n-th element is written into the n-th bit of vector mask k2.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k2. Elements in zmm1 and k2 with the corresponding bit clear in k1 retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    k2[n] = Borrow(tmpSrc3[i+31:i] - zmm1[i+31:i])zmm1[i+31:i] = tmpSrc3[i+31:i] - zmm1[i+31:i]
  }
}
```
Flags Affected

Memory Up-conversion: S_{i32}

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPSUBSETBD - Subtract Int32 Vectors and Set Borrow

Description

Performs an element-by-element subtraction of the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3 from int32 vector zmm1. The result is written into int32 vector zmm1.

In addition, the borrow from the subtraction for the n-th element is written into the n-th bit of vector mask k2.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1 and k2. Elements in zmm1 and k2 with the corresponding bit clear in k1 retain their previous value.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n// integer operation
    k2[n] = Borrow(zmm1[i+31:i] - tmpSrc3[i+31:i])zmm1[i+31:i] = zmm1[i+31:i] - tmpSrc3[i+31:i]}
}
```
Flags Affected

Register Swizzle: S*i*³²

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode #SS(0) If a memory address referencing the SS segment is in a non-canonical form. #GP(0) If the memory address is in a non-canonical form. If a memory operand linear address is not aligned to the data size granularity dictated by SwizzUpConv mode. #PF(fault-code) For a page fault. #NM If CR0.TS[bit 3]=1. #UD If processor model does not implement the specific instruction. If preceded by any REX, F0, F2, F3, or 66 prefixes.

VPTESTMD - Logical AND Int32 Vectors and Set Vector Mask

Description

Performs an element-by-element bitwise AND between int32 vector zmm1 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm2, and uses the result to construct a 16 bit vector mask, with a 0-bit for each element for which the result of the AND was 0, and a 1-bit where the result of the AND was not zero. The final result is written into vector mask k2.

The write-mask does not perform the normal write-masking function for this instruction. While it does enable/disable comparisons, it does not block updating of the destination; instead, if a write-mask bit is 0, the corresponding destination bit is set to 0. Nonetheless, the operation is similar enough so that it makes sense to use the usual write-mask notation. This mode of operation is desirable because the result will be used directly as a write-mask, rather than the normal case where the result is used with a separate writemask that keeps the masked elements inactive.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoadi32(zmm2/mt)
}
for (n = 0; n < 16; n++) {
 k2[n] = 0if(k1[n] != 0) {i = 32*n// signed integer operation
   if ((zmm1[i+31:i] & tmpSrc2[i+31:i]) != 0) {
     k2[n] = 1}
 }
}
```


Flags Affected

None.

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__mmask16 _mm512_test_epi32_mask (__m512i, __m512i); __mmask16 _mm512_mask_test_epi32_mask (__mmask16, __m512i, __m512i);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

VPXORD - Bitwise XOR Int32 Vectors

Description

Performs an element-by-element bitwise XOR between int32 vector zmm2 and the int32 vector result of the swizzle/broadcast/conversion process on memory or int32 vector zmm3. The result is written into int32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]} else {
   tmpSrc3[511:0] = SwizzUpConvLoadi32(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// signed integer operation
    zmm1[i+31:i] = zmm2[i+31:i] ^ tmpSrc3[i+31:i]
  }
}
```
Flags Affected

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VPXORQ - Bitwise XOR Int64 Vectors

Description

Performs an element-by-element bitwise XOR between int64 vector zmm2 and the int64 vector result of the swizzle/broadcast/conversion process on memory or int64 vector zmm3. The result is written into int64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   tmpSrc3[511:0] = zmm3[511:0]
} else {
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>i64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*nzmm1[i+63:i] = zmm2[i+63:i] ^ tmpSrc3[i+63:i]
  }
}
```
Flags Affected

Memory Up-conversion: S*i*⁶⁴

Register Swizzle: S*i*⁶⁴

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

VRCP23PS - Reciprocal of Float32 Vector

Description

Computes the element-by-element reciprocal approximation of the float32 vector on memory or float32 vector zmm2 with 0.912ULP (relative error). The result is written into float32 vector zmm1.

If any source element is NaN, the quietized NaN source value is returned for that element. If any source element is *±∞*, 0*.*0 is returned for that element. Also, if any source element is $±0.0, ±∞$ is returned for that element.

Current implementation of this instruction does not support any SwizzUpConv setting other than "no broadcast and no conversion"; any other SwizzUpConv setting will result in an Invalid Opcode exception.

recip_1ulp() function follows Table [6.26](#page-575-0) when dealing with floating-point special number.

Table 6.26: recip_1ulp() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
tmpSrc2[511:0] = zmm2/mt
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
}
for (n = 0; n < 16; n++) {
```


```
if (k1[n] != 0) {
    i = 32*nzmm1[i+31:i] = recip_1ulp(tmpSrc2[i+31:i])
  }
}
```
SIMD Floating-Point Exceptions

Invalid, Zero.

Denormal Handling

Treat Input Denormals As Zeros : YES

Flush Tiny Results To Zero : YES

Register Swizzle

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

__m512 _mm512_mask_rcp23_ps (__m512, __mmask16, __m512);

Exceptions

Real-Address Mode and Virtual-8086

VRNDFXPNTPD - Round Float64 Vector

Description

Performs an element-by-element rounding of the result of the swizzle/broadcast/conversion from memory or float64 vector zmm2. The rounding result for each element is a float64 containing an integer or fixed-point value, depending on the value of expadj; the direction of rounding depends on the value of RC. The result is written into float64 vector zmm1.

This instruction doesn't actually convert the result to an int64; the results are float64s, just like the input, but are float64s containing the integer or fixed-point values that result from the specified rounding and scaling.

RoundToInt() function follows Table [6.27](#page-578-0) when dealing with floating-point special number.

Table 6.27: RoundToInt() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Immediate Format

Operation

```
RoundingMode = IMMS[1:0]expadj = IMMS[6:4]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad_{f64}(zmm2/m_t)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
   zmm1[i+63:i] =
     RoundToInt(tmpSrc2[i+63:i] * EXPADJ_TABLE[expadj], RoundingMode) /
     EXPADJ_TABLE[expadj]
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Precision.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{f64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512d _mm512_roundfxpnt_adjust_pd (__m512d, int, _MM_EXP_ADJ_ENUM);
- __m512d _mm512_mask_roundfxpnt_adjust_pd (__m512d, __mmask8, __m512d, int , _MM_EXP_ADJ_ENUM);

Exceptions

Real-Address Mode and Virtual-8086

VRNDFXPNTPS - Round Float32 Vector

Description

Performs an element-by-element rounding of the result of the swizzle/broadcast/conversion from memory or float32 vector zmm2. The rounding result for each element is a float32 containing an integer or fixed-point value, depending on the value of expadj; the direction of rounding depends on the value of RC. The result is written into float32 vector zmm1.

This instruction doesn't actually convert the result to an int32; the results are float32s, just like the input, but are float32s containing the integer or fixed-point values that result from the specified rounding and scaling.

RoundToInt() function follows Table [6.28](#page-582-0) when dealing with floating-point special number.

This instruction treats input denormals as zeros according to the DAZ control bit, but does not flush tiny results to zero.

Table 6.28: RoundToInt() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Immediate Format

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Operation

```
RoundingMode = IMM8[1:0]
expadj = IMMS[6:4]if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   tmpSrc2[511:0] = zmm2[511:0]} else {
   tmpSrc2[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm2/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {i = 32*n
    // float32 operation
    zmm1[i+31:i] =
     RoundToInt(tmpSrc2[i+31:i] * EXPADJ_TABLE[expadj], RoundingMode) /
      EXPADJ_TABLE[expadj]
 }
}
```
SIMD Floating-Point Exceptions

Invalid, Precision.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

- __m512 _mm512_roundfxpnt_adjust_ps (__m512, int , _MM_EXP_ADJ_ENUM);
- __m512 _mm512_mask_roundfxpnt_adjust_ps (__m512, __mmask16, __m512, int , _MM_EXP_ADJ_ENUM);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

CHAPTER 6. INSTRUCTION DESCRIPTIONS

64 bit Mode #SS(0) If a memory address referencing the SS segment is in a non-canonical form. #GP(0) If the memory address is in a non-canonical form. If a memory operand linear address is not aligned to the data size granularity dictated by SwizzUpConv mode. #PF(fault-code) For a page fault. #NM If CR0.TS[bit 3]=1. #UD If processor model does not implement the specific instruction. If preceded by any REX, F0, F2, F3, or 66 prefixes.

VRSQRT23PS - Vector Reciprocal Square Root of Float32 Vector

Description

Computes the element-by-element reciprocal square root of the float32 vector on memory or float 32 vector zmm2 with a precision of 0.775ULP (relative error). The result is written into float32 vector zmm1.

If any source element is NaN, the quietized NaN source value is returned for that element. Negative source numbers, as well as *−∞*, return the canonical NaN and set the Invalid Flag (HI) .

Current implementation of this instruction does not support any SwizzUpConv setting other than "no broadcast and no conversion"; any other SwizzUpConv setting will result in an Invalid Opcode exception.

rsqrt_1ulp() function follows Table [6.29](#page-586-0) when dealing with floating-point special number.

For an input value of +/ *−* 0 the instruction returns *−∞* and sets the Divide-By-Zero ϐlag (#Z). Negative numbers should return NaN and set the Invalid flag (#I). Note however that this instruction treats input denormals as zeros of the same sign, so for denormal negative inputs it returns *−∞* and sets the Divide-By-Zero status ϐlag.

Table 6.29: rsqrt_1ulp() special floating-point values behavior

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
tmpSrc2[511:0] = zmm2/mt
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
}
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
   i = 32*nzmm1[i+31:i] = rsqrt_1ulp(tmpSrc2[i+31:i])}
}
```
SIMD Floating-Point Exceptions

Invalid, Zero.

Denormal Handling

Treat Input Denormals As Zeros : YES

Flush Tiny Results To Zero : YES

Register Swizzle

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

VSCALEPS - Scale Float32 Vectors

Description

Performs an element-by-element scale of float32 vector zmm2 by multiplying it by 2^{exp} , where exp is the vector int32 result of the swizzle/broadcast/conversion process on memory or vector int32 zmm3. The result is written into vector float32 zmm1.

This instruction is needed for scaling u and v coordinates according to the mipmap size, which is 2 *mipmap*_*level*, and for the evaluation of Exp2.

Cases where the exponent would go out of range are handled as if multiplication (via vmulps) of zmm2 by 2 *zmm*³ had been performed.

If the result cannot be represented with a float 32, then the properly signed ∞ (for positive scaling operand) or 0 (for negative scaling operand) will be returned.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoadI(zmm3/mt)
}
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*nexp[31:0] = tmpSrc3[i+31:i]// signed int scale operation. float32 multiplication
    zmm1[i+31:i] = zmm2[i+31:i] * 2exp[31:0]
 }
}
```


SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{i32}

Register Swizzle: S_{i32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

```
#UD Instruction not available in these modes
```
Protected and Compatibility Mode

64 bit Mode

#SS(0) If a memory address referencing the SS segment is in a non-canonical form.

VSCATTERDPD - Scatter Float64 Vector With Signed Dword Indices

Description

Down-converts and stores all 8 elements in float64 vector zmm1 to the memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX*, with scale *SCALE*.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Writes to overlapping destination memory locations are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Only writes to overlapping vector indices are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Writes that are not overlapped may happen in any order. Memory ordering with other instructions follows the Intel-64 memory ordering model. Note that this does not account for non-overlapping indices that map into the same physical address locations.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully stored.

Operation

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 8; n++) {
```


```
if (ktemp[n] != 0) {
    i = 64*nj = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[j+31:j] * SCALE)
    pointer[63:0] = mv_t[n]tmp = DownConvStore_{f64}(zmm1[i+63:i], SSS[2:0])if(DownConvStoreSizeOf<sub>f64</sub>(SSS[2:0]) == 8) {
      MemStore(pointer) = tmp[63:0]
    }
    k1[n] = 0}
k1[15:8] = 0
```
SIMD Floating-Point Exceptions

None.

}

Memory Down-conversion: D_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_mask_i32loextscatter_pd (void*, __mmask8, __m512i, __m512d, _MM_DOWNCONV_PD_ENUM, int, int);
- void _mm512_i32loscatter_pd (void*, _m512i, _m512d, int);
- void _mm512_mask_i32loscatter_pd (void*, _mmask8, _m512i, _m512d, int);

Exceptions

Real-Address Mode and Virtual-8086

VSCATTERDPS - Scatter Float32 Vector With Signed Dword Indices

Description

Down-converts and stores all 16 elements in float32 vector zmm1 to the memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX*, with scale *SCALE*.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the gather/scatter sequence have been loaded/stored and hence, the write-mask bits all are zero).

Writes to overlapping destination memory locations are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Only writes to overlapping vector indices are guaranteed to be ordered with respect to each other (from LSB to MSB of the source registers). Writes that are not overlapped may happen in any order. Memory ordering with other instructions follows the Intel-64 memory ordering model. Note that this does not account for non-overlapping indices that map into the same physical address locations.

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully stored.

Operation

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
```


CHAPTER 6. INSTRUCTION DESCRIPTIONS

```
if (ktemp[n] != 0) {
    i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]tmp = DownConvStore_{f32}(zmm1[i+31:i], SSS[2:0])if(DownConvStoreSizeOf<sub>f32</sub>(SSS[2:0]) == 4) {
      MemStore(pointer) = tmp[31:0]
    } else if(DownConvStoreSizeOff32(SSS[2:0]) == 2) {
      MemStore(pointer) = tmp[15:0]
    } else if(DownConvStoreSizeOff32(SSS[2:0]) == 1) {
      MemStore(pointer) = tmp[7:0]
    }
   k1[n] = 0}
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Memory Down-conversion: D_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_mask_i32extscatter_ps (void*, __mmask16, __m512i, __m512, _MM_DOWNCONV_PS_ENUM, int, int);
- void _mm512_i32scatter_ps (void*, $_m512i, m512, int$);
- void _mm512_mask_i32scatter_ps (void*, _mmask16, _m512i, _m512, int);

Exceptions

(intel)

Real-Address Mode and Virtual-8086

VSCATTERPF0DPS - Scatter Prefetch Float32 Vector With Signed Dword Indices Into L1

Description

Prefetches into the L1 level of cache the memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX*, with scale *SCALE*, with request for ownership (exclusive). Up-conversion operand specifies the granularity used by compilers to better encode the instruction if a displacement, using disp8*N feature, is provided when specifying the address. If any memory access causes any type of memory exception, the memory access will be considered as completed (destination mask updated) and the exception ignored. Up-conversion parameter is optional, and it is used to correctly encode disp8*N.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the prefetch sequence have been prefetched and hence, the write-mask bits all are zero).

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after up-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully stored.

Note that both gather and scatter prefetches set the access bit (A) in the related TLB page entry. Scatter prefetches (which prefetch data with RFO) do not set the dirty bit (D).

Operation

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
exclusive = 1
evicthintpre = MVEX.EH
```


```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
  if (ktemp[n] != 0) {
    i = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]FetchL1cacheLine(pointer, exclusive, evicthintpre)
    k1[n] = 0}
}
```
SIMD Floating-Point Exceptions

None.

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_prefetch_i32extscatter_ps (void*, __m512i, _MM_UPCONV_PS_ENUM, int, int);
- void _mm512_mask_prefetch_i32extscatter_ps(void*, __mmask16, __m512i, _MM_UPCONV_PS_ENUM, int, int);
- void _mm512_prefetch_i32scatter_ps(void*, _m512i, int, int);
- void _mm512_mask_prefetch_i32scatter_ps(void*, _mmask16, _m512i, int, int);

Exceptions

Real-Address Mode and Virtual-8086

VSCATTERPF0HINTDPD - Scatter Prefetch Float64 Vector Hint With Signed Dword Indices

Description

The instruction specifies a set of 8 float64 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* as a performance hint that a real scatter instruction with the same set of sources will be invoked. A programmer may execute this instruction before a real scatter instruction to improve its performance.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults. This instructions does not modify any kind of architectural state (including the write-mask).

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Operation

```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 8; n++) {
  if (k1[n] := 0) {
    i = 64*ni = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[j+31:j] * SCALE)
    pointer[63:0] = mv_t[n]HintPointer(pointer)
  }
}
```
SIMD Floating-Point Exceptions

None.

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Memory Up-conversion: U_{f64}

Intel® C/C++ Compiler Intrinsic Equivalent

None

Exceptions

Real-Address Mode and Virtual-8086

If no write mask is provided or selected write-mask is k0.

VSCATTERPF0HINTDPS - Scatter Prefetch Float32 VectorHint With Signed Dword Indices

Description

The instruction specifies a set of 16 float32 memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX* with scale *SCALE* as a performance hint that a real scatter instruction with the same set of sources will be invoked. A programmer may execute this instruction before a real scatter instruction to improve its performance.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults. This instructions does not modify any kind of architectural state (including the write-mask).

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element before up-conversion.

Operation

```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
  if (k1[n] := 0) {
    i = 32*n// mvt[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]HintPointer(pointer)
 }
}
```
SIMD Floating-Point Exceptions

None.

CHAPTER 6. INSTRUCTION DESCRIPTIONS

Memory Up-conversion: U_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

None

Exceptions

Real-Address Mode and Virtual-8086

If no write mask is provided or selected write-mask is k0.

VSCATTERPF1DPS - Scatter Prefetch Float32 Vector With Signed Dword Indices Into L2

Description

Prefetches into the L2 level of cache the memory locations pointed by base address *BASE*_*ADDR* and doubleword index vector *V INDEX*, with scale *SCALE*, with request for ownership (exclusive). Down-conversion operand specifies the granularity used by compilers to better encode the instruction if a displacement, using disp8*N feature, is provided when specifying the address. If any memory access causes any type of memory exception, the memory access will be considered as completed (destination mask updated) and the exception ignored. Down-conversion parameter is optional, and it is used to correctly encode disp8*N.

Note the special mask behavior as only a subset of the active elements of write mask k1 are actually operated on (as denoted by function *SELECT*_*SUBSET*). There are only two guarantees about the function: (a) the destination mask is a subset of the source mask (identity is included), and (b) on a given invocation of the instruction, **at least** one element (the least significant enabled mask bit) will be selected from the source mask.

Programmers should always enforce the execution of a gather/scatter instruction to be re-executed (via a loop) until the full completion of the sequence (i.e. all elements of the prefetch sequence have been prefetched and hence, the write-mask bits all are zero).

This instruction has special disp8*N and alignment rules. N is considered to be the size of a single vector element after down-conversion.

Note also the special mask behavior as the corresponding bits in write mask k1 are reset with each destination element being updated according to the subset of write mask k1. This is useful to allow conditional re-trigger of the instruction until all the elements from a given write mask have been successfully stored.

Note that both gather and scatter prefetches set the access bit (A) in the related TLB page entry. Scatter prefetches (which prefetch data with RFO) do not set the dirty bit (D).

Operation

```
// instruction works over a subset of the write mask
ktemp = SELECT_SUBSET(k1)
exclusive = 1
evicthintpre = MVEX.EH
```


```
// Use mvt as vector memory operand (VSIB)
for (n = 0; n < 16; n++) {
  if (ktemp[n] != 0) {
    i = 32*n\frac{1}{\sqrt{mv_t}}[n] = BASE_ADDR + SignExtend(VINDEX[i+31:i] * SCALE)
    pointer[63:0] = mv_t[n]FetchL2cacheLine(pointer, exclusive, evicthintpre)
    k1[n] = 0}
}
```
SIMD Floating-Point Exceptions

None.

Memory Down-conversion: D_{f32}

Intel® C/C++ Compiler Intrinsic Equivalent

- void _mm512_prefetch_i32extscatter_ps (void*, __m512i, _MM_UPCONV_PS_ENUM, int, int);
- void _mm512_mask_prefetch_i32extscatter_ps(void*, __mmask16, __m512i, _MM_UPCONV_PS_ENUM, int, int);
- void _mm512_prefetch_i32scatter_ps(void*, _m512i, int, int);
- void _mm512_mask_prefetch_i32scatter_ps(void*, _mmask16, _m512i, int, int);

Exceptions

Real-Address Mode and Virtual-8086

VSUBPD - Subtract Float64 Vectors

Description

Performs an element-by-element subtraction from float64 vector zmm2 of the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3. The result is written into float64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {i = 64*n// float64 operation
    zmm1[i+63:i] = zmm2[i+63:i] - tmpSrc3[i+63:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

VSUBPS - Subtract Float32 Vectors

Description

Performs an element-by-element subtraction from float32 vector zmm2 of the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// float32 operation
    zmm1[i+31:i] = zmm2[i+31:i] - tmpSrc3[i+31:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

If preceded by any REX, F0, F2, F3, or 66 prefixes.

VSUBRPD - Reverse Subtract Float64 Vectors

Description

Performs an element-by-element subtraction of float64 vector zmm2 from the float64 vector result of the swizzle/broadcast/conversion process on memory or float64 vector zmm3. The result is written into float64 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f64</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 8; n++) {
  if(k1[n] != 0) {
    i = 64*n// float64 operation
    zmm1[i+63:i] = -zmm2[i+63:i] + tmpSrc3[i+63:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*64}

Register Swizzle: S_{*f*64}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

If preceded by any REX, F0, F2, F3, or 66 prefixes.

VSUBRPS - Reverse Subtract Float32 Vectors

Description

Performs an element-by-element subtraction of float32 vector zmm2 from the float32 vector result of the swizzle/broadcast/conversion process on memory or float32 vector zmm3. The result is written into float32 vector zmm1.

This instruction is write-masked, so only those elements with the corresponding bit set in vector mask register k1 are computed and stored into zmm1. Elements in zmm1 with the corresponding bit clear in k1 retain their previous values.

Operation

```
if(source is a register operand and MVEX.EH bit is 1) {
   if(SSS[2]==1) Supress_Exception_Flags() // SAE
  // SSS are bits 6-4 from the MVEX prefix encoding. For more details, see Table 2.14
   RoundingMode = SSS[1:0]tmpSrc3[511:0] = zmm3[511:0]
} else {
   RoundingMode = MXCSR.RC
   tmpSrc3[511:0] = SwizzUpConvLoad<sub>f32</sub>(zmm3/m<sub>t</sub>)}
for (n = 0; n < 16; n++) {
  if(k1[n] != 0) {
    i = 32*n// float32 operation
    zmm1[i+31:i] = -zmm2[i+31:i] + tmpSrc3[i+31:i]
  }
}
```
SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Denormal Handling

Treat Input Denormals As Zeros : (MXCSR.DAZ)? YES : NO

Flush Tiny Results To Zero : (MXCSR.FZ)? YES : NO

Memory Up-conversion: S_{*f*32}

Register Swizzle: S_{f32}

MVEX.EH=0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Appendix A

Scalar Instruction Descriptions

In this Chapter all the special scalar instructions introduced with the Knights Corner instruction set are described.

Opcode Instruction Description

Description

Invalidates from the first-level cache the cache line containing the specified linear address (updating accordingly the cache hierarchy if the line is dirty). Note that, unlike CLFLUSH, the invalidation is not broadcasted throughout the cache coherence domain.

The MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint intended for performance and may be speculative, thus may be dropped or specify invalid addresses without causing problems. The instruction does not produce any type of memory-related fault.

Operation

FlushL1CacheLine(linear_address)

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void mm clevict (const void*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes Protected and Compatibility Mode #UD Instruction not available in these modes 64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

Description

Invalidates from the second-level cache the cache line containing the specified linear address (updating accordingly the cache hierarchy if the line is dirty). Note that, unlike CLFLUSH, the invalidation is not broadcasted throughout the cache coherence domain.

The MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint intended for performance and may be speculative, thus may be dropped or specify invalid addresses without causing problems. The instruction does not produce any type of memory-related fault.

Operation

FlushL2CacheLine(linear_address)

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void mm clevict (const void*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes Protected and Compatibility Mode #UD Instruction not available in these modes 64 bit Mode If preceded by any REX, F0, F2, F3, or 66 prefixes.

If operand is not a memory location.

DELAY - Stall Thread

Description

Hints that the processor should not fetch/issue instructions for the current thread for the specified number of clock cycles in register source. The maximum number of clock cycles is limited to 2 ³²*−*1 (32 bit counter). The instructions is speculative and could be executed as a NOP by a given processor implementation.

Any of the following events will cause the processor to start fetching instructions for the delayed thread again: the counter counting down to zero, an NMI or SMI, a debug exception, a machine check exception, the BINIT# signal, the INIT# signal, or the RESET# signal. The instruction may exit prematurely due to any interrupt (e.g. an interrupt on another thread on the same core).

This instruction must properly handle the case where the current clock count turns over. This can be accomplished by performing the subtraction shown below and treating the result as an unsigned number.

This instruction should prevent the issuing of additional instructions on the issuing thread as soon as possible, to avoid the otherwise likely case where another instruction on the same thread that was issued 3 or 4 clocks later has to be killed, creating a pipeline bubble.

If, on any given clock, all threads are non-runnable, then any that are non-runnable due to the execution of DELAY may or may not be treated as runnable threads.

Notes about Knights Corner implementation:

• In Knights Corner, the processor won't execute from a "delayed" thread before the delay counter has expired, even if there are non-runnable threads at any given point in time.

Operation

```
START_CLOCK = CURRENT_CLOCK_COUNT
DELAY_SLOTS = SRC
if(DELAY_SLOTS > 0xFFFFFFFF) DELAY_SLOTS = 0xFFFFFFFF
while ( (CURRENT_CLOCK_COUNT - START_CLOCK) < DELAY_SLOTS )
{
     *avoid fetching/issuing from the current thread*
}
```


Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_delay_32 (unsigned int);

void _mm_delay_64 (unsigned _int64);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is a memory location.

LZCNT - Leading Zero Count

Description

Counts the number of leading most significant zero bits in a source operand (second operand) returning the result into a destination (first operand).

LZCNT is an extension of the BSR instruction. The key difference between LZCNT and BSR is that LZCNT provides operand size as output when source operand is zero, while in the case of BSR instruction, if source operand is zero, the content of destination operand are undefined.

ZF flag is set when the most significant set bit is bit OSIZE-1. CF is set when the source has no set bit.

Operation

```
temp = OPERAND_SIZE - 1DEST = 0while( temp \ge 0) AND (SRC[temp] == 0){
   temp = temp - 1DEST = DEST + 1}
if(DEST == OPERAND_SIZE) {
  CF = 1} else {
  CF = 0}
if(DEST == 0) ZF = 1} else {
  ZF = 0}
```
intel

Flags Affected

- ZF flag is set to 1 in case of zero output (most significant bit of the source is set), and to 0 otherwise
- CF flag is set to 1 if input was zero and cleared otherwise.
- The PF, OF, AF and SF flags are set to 0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If second operand is a memory location.

POPCNT - Return the Count of Number of Bits Set to 1

Operation

```
tmp = 0for (i=0; i<OPERAND_SIZE; i++)
{
     if(SRC[i] == 1) tmp = tmp + 1
}
DEST = tmp
```
Flags Affected

- The ZF flag is set according to the result (if SRC==0)
- The OF, SF, AF, CF and PF flags are set to 0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If second operand is a memory location.

SPFLT - Set performance monitor filtering mask

Description

SPFLT enables/disables performance monitoring on the currently executing thread only based on the LSB value of the source.

SPFLT instruction is a model specific instruction and is not part of Intel® Architecture. The bit(s) and register(s) modified are model-specific and may vary by processor implementation.

The PERF_SPFLT_CTRL model-specific register modified by SPFLT instruction may also be read / modified with the RDMSR / WRMSR instructions, when executing at privilege level 0.

The PERF_SPFLT_CTRL MSR is thread specific. SPFLT execution moves LSB of source (EAX) into the USR_PREF bit (bit 63) in the PERF_SPFLT_CTRL MSR. The lower N bits, called CNTR_x_SPFLT_EN (bits N-1:0, 1 per counter), in PERF_SPFLT_CTRL MSR control whether the USR_PREF bit affects enabling of performance monitoring for the corresponding counter.

SPFLT instruction does not modify the CNTR_x_SPFLT_EN bits, where as RDMSR and WRMSR read / modify all bits of the PERF_SPFLT_CTRL MSR.

Enabling Performance countering

On a per thread basis, a performance monitoring counter n is incremented if, and only if:

- 1. PERF GLOBAL CTRL[n] is set to 1
- 2. IA32 PerfEvtSeln[22] is set to 1 (where 'n' is the enabled counter)
- 3. PERF_SPFLT_CTRL[n] is set to 0, or, PERF_SPFLT_CTRL[63] (USR_PREF) is set to 1.
- 4. The desired event is asserted for thread id T

APPENDIX A. SCALAR INSTRUCTION DESCRIPTIONS

Operation

(* i is the thread ID of the current executing thread *) PerfFilterMask[i][0] = SRC[0];

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_spflt_32 (unsigned int); void _mm_spflt_64 (unsigned _int64);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

#UD If processor model does not implement SPFLT. If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is a memory location.

TZCNT - Trailing Zero Count

Description

Searches the source operand (second operand) for the least significant set bit (1 bit). If a least significant 1 bit is found, its bit index is stored in the destination operand; otherwise, the destination operand is set to the operand size.

ZF flag is set when the least significant set bit is bit 0. CF is set when the source has no set bit.

Operation

```
index = 0if( SRC[OPERAND_SIZE-1:0] == 0 )
{
     DEST = OPERAND_SIZE
     CF = 1}
else
{
     while(SRC[index] == 0){
          index = index + 1}
     DEST = index
     CF = 0}
```
Flags Affected

- The ZF is set according to the result
- The CF is set if SRC is zero
- The PF, OF, AF and SF flags are set to 0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

TZCNTI - Initialized Trailing Zero Count

Description

Searches the source operand (second operand) for the least significant set bit (1 bit) greater than bit DEST (where DEST is the destination operand, the first operand). If a least significant 1 bit is found, its bit index is stored in the destination operand; otherwise, the destination operand is set to the operand size. The value of DEST is a signed offset from bit 0 of the source operand. Any negative DEST value will produce a search starting from bit 0, like TZCNT. Any DEST value equal to or greater than (OPERAND_SIZE-1) will cause the destination operand to be set to the operand size.

This instruction allows continuation of searches through bit vectors without having to mask off each least significant 1-bit before restarting, as is required with TZCNT.

The functionality of this instruction is exactly the same as for the TZCNT instruction, except that the search starts at bit DEST+1 rather than bit 0.

CF is set when the specified index goes beyond the operand size or there is no set bit between the index and the MSB bit of the source.

Operation

```
// DEST is a signed operand, no overflow
if (DEF[0SIZE-1:0] < 0) index = 0
else index = DEST + 1
if( ( index > 0PERAND_SIZE-1 ) || (SRC[0PERAND_SIZE-1:index] == 0 ) )
{
    DEST = OPERAND_SIZE
    CF=1}
else
{
    while(SRC[index] == 0){
          index = index + 1}
    DEST = indexCF=0
```


Flags Affected

}

- The ZF is set according to the result
- The CF is set if SRC is zero betwen index and MSB, or index is greater than the operand size.
- The PF, OF, AF and SF flags are set to 0

Intel® C/C++ Compiler Intrinsic Equivalent

Exceptions

Real-Address Mode and Virtual-8086

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If second operand is a memory location.

VPREFETCH0 - Prefetch memory line using T0 hint

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- exclusive: make line Exclusive in the L1 cache (unless it's already Exclusive or Modified in the L1 cache).
- **nthintpre (NTH):** load data into the L1 nontemporal cache rather than the L1 temporal cache. Data will be loaded in the #TIDth way and made MRU. Data should still be cached normally in the L2 and higher caches.

Note that in Knights Corner, the hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 0
nthintpre = 0FetchL1CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCH1 - Prefetch memory line using T1 hint

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- exclusive: make line Exclusive in the L2 cache (unless it's already Exclusive or Modified in the L2 cache).
- **nthintpre (NTH):** load data into the L2 nontemporal cache rather than the L2 temporal cache. Data will be loaded in the #TIDth way and made MRU. Data should still be cached normally in the L2 and higher caches.

Note that in Knights Corner, the hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 0
nthintpre = 0FetchL2CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCH2 - Prefetch memory line using T2 hint

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- exclusive: make line Exclusive in the L2 cache (unless it's already Exclusive or Modified in the L2 cache).
- **nthintpre (NTH):** load data into the L2 nontemporal cache rather than the L2 temporal cache. Data will be loaded in the #TIDth way and made MRU. Data should still be cached normally in the L2 and higher caches.

Note that in Knights Corner, the hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 0
nthintpre = 1
FetchL2CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCHE0 - Prefetch memory line using T0 hint, with intent to write

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- **exclusive:** make line Exclusive in the L1 cache (unless it's already Exclusive or Modified in the L1 cache).
- **nthintpre (NTH):** load data into the L1 nontemporal cache rather than the L1 temporal cache. Data will be loaded in the #TIDth way and made MRU. Data should still be cached normally in the L2 and higher caches.

In Knights Corner, the hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 1
nthintpre = 0FetchL1CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCHE1 - Prefetch memory line using T1 hint, with intent to write

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- **exclusive:** make line Exclusive in the L2 cache (unless it's already Exclusive or Modified in the L2 cache).
- **nthintpre (NTH):** load data into the L2 nontemporal cache rather than the L2 temporal cache. The data will be loaded in the #TIDth way and making the data MRU. Data should still be cached normally in the L2 and higher caches.

The hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 1
nthintpre = 0FetchL2CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCHE2 - Prefetch memory line using T2 hint, with intent to write

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- **exclusive:** make line Exclusive in the L2 cache (unless it's already Exclusive or Modified in the L2 cache).
- **nthintpre (NTH):** load data into the L2 nontemporal cache rather than the L2 temporal cache. Data will be loaded in the #TIDth way and made MRU. Data should still be cached normally in the L2 and higher caches.

Note that in Knights Corner, the hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 1
nthintpre = 1
FetchL2CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCHENTA - Prefetch memory line using NTA hint, with intent to write

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- **exclusive:** make line Exclusive in the L1 cache (unless it's already Exclusive or Modified in the L1 cache).
- **nthintpre (NTH):** load data into the L1 nontemporal cache rather than the L1 temporal cache. The data will be loaded in the #TIDth way and making the data MRU. Data should still be cached normally in the L2 and higher caches.

The hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 1
nthintpre = 1
FetchL1CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

VPREFETCHNTA - Prefetch memory line using NTA hint

Description

This is very similar to the existing IA-32 prefetch instruction, PREFETCH0, as described in *IA-32 Intel*® *Architecture Software Developer's Manual: Volume 2*. If the line selected is already present in the cache hierarchy at a level closer to the processor, no data movement occurs. Prefetches from uncacheable or WC memory are ignored.

In contrast with the existing prefetch instruction, the MVEX form of this instruction uses disp8*64 addressing. Displacements that would normally be 8 bits according to the ModR/M byte are still 8 bits but scaled by 64 so that they have cache-line granularity. VEX forms of this instruction uses regular disp8 addressing.

This instruction is a hint and may be speculative, and may be dropped or specify invalid addresses without causing problems or memory related faults.

This instruction contains a set of hint attributes that modify the prefetching behavior:

- exclusive: make line Exclusive in the L1 cache (unless it's already Exclusive or Modified in the L1 cache).
- **nthintpre (NTH):** load data into the L1 nontemporal cache rather than the L1 temporal cache. Data will be loaded in the #TIDth way and made MRU. Data should still be cached normally in the L2 and higher caches.

In Knights Corner, the hardware drops VPREFETCH if it hits L1 (so it becomes transparent to L2). Consequently, this instructon is not a good solution to avoid hot L1/cold L2 performance problems. Prefetches set the access bit (A) in the related TLB page entry, but prefetches with exclusive access (RFO) do not set the dirty bit (D).

PREFETCH Hint equivalence for Knights Corner hardware

Operation

```
exclusive = 0
nthintpre = 1
FetchL1CacheLine(effective_address, exclusive, nthintpre)
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_prefetch (char const*, int);

Exceptions

Real-Address Mode and Virtual-8086

#UD Instruction not available in these modes

Protected and Compatibility Mode

#UD Instruction not available in these modes

64 bit Mode

If preceded by any REX, F0, F2, F3, or 66 prefixes. If operand is not a memory location.

Appendix B

Knights Corner 64 bit Mode Scalar Instruction Support

In 64 bit mode, Knights Corner supports a subset of the Intel 64 Architecture instructions. The 64 bit mode instructions supported by Knights Corner are listed in this chapter.

B.1 64 bit Mode General-Purpose and X87 Instructions

Knights Corner supports most of the general-purpose register (GPR) and X87 instructions in 64 bit mode. They are listed in Table [B.2.](#page-656-0)

64 bit Mode GPR and X87 Instructions in Knights Corner:

B.2 Knights Corner 64 bit Mode Limitations

In 64 bit mode, Knights Corner supports a subset of the Intel 64 Architecture instructions. The following summarizes Intel 64 Architecture instructions that are not supported in Knights Corner:

- Instructions that operate on MMX registers
- Instructions that operate on XMM registers
- Instructions that operate on YMM registers

GPR and X87 Instructions Not Supported in Knights Corner

B.3 LDMXCSR - Load MXCSR Register

Description

Loads the source operand into the MXCSR control/status register. The source operand is a 32 bit memory location. See MXCSR Control and Status Register in Chapter 10, of the IA-32 Intel Architecture Software Developers Manual, Volume 1, for a description of the MXCSR register and its contents. See chapter 3 of this document for a description of the new Knights Corner's MXCSR feature bits.

The LDMXCSR instruction is typically used in conjunction with the STMXCSR instruction, which stores the contents of the MXCSR register in memory.

The default MXCSR value at reset is 0020 0000H (DUE=1, FZ=0, RC=00, PM=0, UM=0, OM=0, ZM=0, DM=0, IM=0, DAZ=0, PE=0, UE=0, OE=0, ZE=0, DE=0, IE=0).

Any attempt to set to 1 reserved bits in control register MXCSR will produce a #GP fault:

Additionally, any attempt to set MXCSR.DUE (bit 21) to 0 will produce a #GP fault:

This instructions operation is the same in non-64 bit modes and 64 bit mode.

Operation

MXCSR = MemLoad(m32)

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _mm_setcsr (unsigned int)

Exceptions

intel.

B.4 FXRSTOR - Restore x87 FPU and MXCSR State

Description

See Intel64® Intel® Architecture Software Developer's Manual for the description of the original x86 instruction.

Reloads the x87 FPU and the MXCSR state from the 512-byte memory image specified in the source operand. This data should have been written to memory previously using the FXSAVE instruction, and in the same format as required by the operating modes. The first byte of the data should be located on a 16-byte boundary. There are three distinct layout of the FXSAVE state map: one for legacy and compatibility mode, a second format for 64 bit mode with promoted operandsize, and the third format is for 64 bit mode with default operand size.

Knights Corner follows the same layouts as described in Intel64[®] Intel[®] Architecture Software Developer's Manual.

The state image referenced with an FXRSTOR instruction must have been saved using an FXSAVE instruction or be in the same format as required by Intel64 Intel[®] Architecture Software Developer's Manual. Referencing a state image saved with an FSAVE, FNSAVE instruction or incompatible field layout will result in an incorrect state restoration.

The FXRSTOR instruction does not flush pending x87 FPU exceptions. To check and raise exceptions when loading x87 FPU state information with the FXRSTOR instruction, use an FWAIT instruction after the FXRSTOR instruction.

Note that XMM15-0 registers are logically aliased to the the low 128-bit portions of Knights Corner registers V15 through V0 (ZMM15-0). Therefore, FXRSTOR must restore the contents of the low 128-bit portions of registers V15 through V0.

Any attempt to set reserved bits in control register MXCSR to 1 will produce a #GP fault:

Additionally, any attempt to set MXCSR.DUE (bit 21) to 0 will produce a #GP fault:

Operation

(x87 FPU, MXCSR, XMM) = MemLoad(SRC);

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _fxrstor64 (void*);

Exceptions

intel.

B.5 FXSAVE - Save x87 FPU and MXCSR State

Description

See Intel64® Intel® Architecture Software Developer's Manual for the description of the original x86 instruction.

Saves the current state of the x87 FPU, XMM, and MXCSR registers to a 512-byte memory location specified in the destination operand. The content layout of the 512 byte region depends on whether the processor is operating in non- 64 bit operating modes or 64 bit sub-mode of IA-32e mode.

Bytes 464:511 are available to software use. The processor does not write to bytes 464:511 of an FXSAVE area.

Knights Corner follows a similar layout as described in Intel64[®] Intel[®] Architecture Software Developer's Manual.

All bits set to 0 in the MXCSR_MASK value indicate reserved bits in the MXCSR register. Thus, if the MXCSR_MASK value is ANDd with a value to be written into the MXCSR register, the resulting value will be assured of having all its reserved bits set to 0, preventing the possibility of a general-protection exception being generated when the value is written to the MXCSR register.

Note that XMM15-0 registers are logically aliased to the the low 128-bit portions of Knights Corner registers V15 through V0 (ZMM15-0). Therefore, FXSAVE must save the contents of the low 128-bit portions of registers V15 through V0.

Operation

```
if(64 bit Mode)
{
  if(REX.W == 1){
     MemStore(m512byte) = Save64BitPromotedFxsave(x87 FPU, XMM15-XMM0, MXCSR);
   }
   else {
     MemStore(m512byte) = Save64BitDefaultFxsave(x87 FPU, XMM15-XMM0, MXCSR);
   }
}
else {
  MemStore(m512byte) = SaveLegacyFxsave(x87 FPU, XMM7-XMM0, MXCSR);
```


}

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

void _fxsave64 (void*);

Exceptions

B.6 RDPMC - Read Performance-Monitoring Counters

Description

Loads the 40-bit performance-monitoring counter specified in the ECX register into registers EDX:EAX. The EDX register is loaded with the high-order 8 bits of the counter and the EAX register is loaded with the low-order 32 bits. The counter to be read is specified with an unsigned integer placed in the ECX register.

The Knights Corner co-processor has 2 performance monitoring counters per thread, specified with 0000H through 0001H, respectively, in the ECX register.

When in protected or virtual 8086 mode, the performance-monitoring counters enabled (PCE) flag in register CR4 restricts the use of the RDPMC instruction as follows. When the PCE flag is set, the RDPMC instruction can be executed at any privilege level; when the flag is clear, the instruction can only be executed at privilege level 0. (When in real-address mode, the RDPMC instruction is always enabled.)

The performance-monitoring counters can also be read with the RDMSR instruction, when executing at privilege level 0.

The performance-monitoring counters are event counters that can be programmed to count events such as the number of instructions decoded, number of interrupts received, or number of cache loads. Appendix A, Performance-Monitoring Events, in the IA-32 Intel® Architecture Software Developers Manual, Volume 3, lists the events that can be counted for the Intel® Pentium® 4, Intel Xeon®, and earlier IA-32 processors.

The RDPMC instruction is not a serializing instruction; that is, it does not imply that all the events caused by the preceding instructions have been completed or that events caused by subsequent instructions have not begun. If an exact event count is desired, software must insert a serializing instruction (such as the CPUID instruction) before and/or after the RDPMC instruction.

The RDPMC instruction can execute in 16 bit addressing mode or virtual-8086 mode; however, the full contents of the ECX register are used to select the counter, and the event count is stored in the full EAX and EDX registers.

The RDPMC instruction was introduced into the IA-32 Architecture in the Intel® Pentium® Pro processor and the Intel® Pentium® processor with Intel® MMX™ technology. The earlier Intel[®] Pentium[®] processors have performance-monitoring counters, but they must be read with the RDMSR instruction.

In 64 bit mode, RDPMC behavior is unchanged from 32 bit mode. The upper 32 bits of RAX and RDX are cleared.

Operation

```
if ( ( (ECX[31:0] >= 0) && (ECX[31:0] < 2)
  & (CR4.PCE = 1) || (CPL = 0) || (CRO.PE = 0)))
{
   if(64 bit Mode)
   {
     RAX[31:0] = PMC(ECX[31:0])[31:0]; (* 40-bit read *)
     RAX[63:32] = 0;RDX[31:0] = PMC(ECX[31:0])[39:32];
     RDX[63:32] = 0;}
  else
  {
     EAX = PMC(ECX[31:0])[31:0]; (* 40-bit read *)
     EDX = PMC(ECX[31:0])[39:32];
  }
}
else
{
   #GP(0)
}
```
Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

__int64 _rdpmc (int);

Exceptions

TBD

B.7 STMXCSR - Store MXCSR Register

Description

Stores the contents of the MXCSR control and status register to the destination operand. The destination operand is a 32 bit memory location.

This instructions operation is the same in non-64 bit modes and 64 bit mode.

Operation

 $MemStore(m32) = MXCSR$

Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

unsigned int _mm_getcsr (void)

Exceptions

B.8 CPUID - CPUID Identification

Description

The ID flag (bit 21) in the EFLAGS register indicates support for the CPUID instruction. If a software procedure can set and clear this flag, the processor executing the procedure supports the CPUID instruction. This instruction operates the same in non-64 bit modes and 64 bit mode.

CPUID returns processor identification and feature information in the EAX, EBX, ECX, and EDX registers. The instructions output is dependent on the contents of the EAX register upon execution. For example, the following pseudo-code loads EAX with 00H and causes CPUID to return a Maximum Return Value and the Vendor Identification String in the appropriate registers:

```
MOV EAX, 00H
CPUID
```
Table [B.4](#page-670-0) through [B.7](#page-673-0) shows information returned, depending on the initial value loaded into the EAX register. Table [B.3](#page-669-0) shows the maximum CPUID input value recognized for each family of IA-32 processors on which CPUID is implemented. Since Intel® Pentium® 4 family of processors, two types of information are returned: *basic* and *extended* function information. Prior to that, only the basic function information was returned. The first is accessed with EAX=0000000xh while the second is accessed with EAX=8000000xh. If a value is entered for CPUID.EAX that is invalid for a particular processor, the data for the highest basic information leaf is returned.

CPUID can be executed at any privilege level to serialize instruction execution. Serializing instruction execution guarantees that any modifications to flags, registers, and memory for previous instructions are completed before the next instruction is fetched and executed.

INPUT EAX = 0: Returns CPUID's Highest Value for Basic Processor Information and the Vendor Identification String

When CPUID executes with EAX set to 0, the processor returns the highest value the CPUID recognizes for returning basic processor information. The value is returned in the EAX register (see Table [B.4](#page-670-0) and is processor specific. A vendor identification string is also returned in EBX, EDX, and ECX. For Intel® processors, the string is "GenuineIntel" and is expressed:

EBX = 756e6547h $(*$ "Genu", with G in the low nibble of BL $*)$ EDX = 49656e69h $(* "inel", with i in the low nibble of DL *)$

Table B.3: Highest CPUID Source Operand for IA-32 Processors

ECX = $6c65746eh$ (* "ntel", with n in the low nibble of CL *)

INPUT EAX = 1: Returns Model, Family, Stepping Information

When CPUID executes with EAX set to 1, version information is returned in EAX. Extended family, extended model, model, family, and processor type for the processor code-named Knights Corner is as follows:

- Extended Model: 0000B
- Extended Family: 0000_0000B
- Model: *see table*
- Family: 1011B
- Processor Type: 00B

INPUT EAX = 1: Returns Additional Information in EBX

When CPUID executes with EAX set to 1, additional information is returned to the EBX register:

• Brand index (low byte of EBX) - this number provides an entry into a brand string table that contains brand strings for IA-32 processors. More information about this field is provided later in this section.

Table B.4: Information Returned by CPUID Instruction

- CLFLUSH/CLEVICTn instruction cache line size (second byte of EBX) this number indicates the size of the cache line flushed with CLEVICT1 instruction in 8-byte increments. This field was introduced in the Intel® Pentium® 4 processor.
- Local APIC ID (high byte of EBX) this number is the 8-bit ID that is assigned to the local APIC on the processor during power up. This field was introduced in the Intel® Pentium® 4 processor.

INPUT EAX = 1: Returns Feature Information in ECX and EDX

When CPUID executes with EAX set to 1, feature information is returned in ECX and EDX.

Table B.5: Information Returned by CPUID Instruction (Contd.)

- Table [B.8](#page-677-0) through Table [B.9](#page-678-0) show encodings for EDX.
- Table [B.10](#page-679-0) show encodings for ECX.

For all feature flags, a 1 indicates that the feature is supported. Use Intel® to properly interpret feature flags.

INPUT EAX = 2: Cache and TLB Information Returned in EAX, EBX, ECX, EDX

Knights Corner considers leaf 2 to be reserved, so no cache and TLB information is returned when CPUID executes with EAX set to 2.

INPUT EAX = 3: Serial Number Information

Knights Corner does not implement *Processor Serial Number* support, as signalled by feature bit CPUID.EAX[01h].EDX.PSN. Therefore, all the returned fields are considered reserved.

INPUT EAX = 4: Returns Deterministic Cache Parameters for Each Level

When CPUID executes with EAX set to 4 and ECX contains an index value, the processor returns encoded data that describe a set of deterministic cache parameters (for the cache level associated with the input in ECX).

Table B.6: Information Returned by CPUID Instruction. 8000000xH leafs.

Software can enumerate the deterministic cache parameters for each level of the cache hierarchy starting with an index value of 0, until the parameters report the value associated with the cache type field is 0. The architecturally defined fields reported by deterministic cache parameters are documented in Table[B.5](#page-671-0). The associated cache structures described by the different ECX descriptors are:

- ECX=0: Instruction Cache (I1)
- ECX=1: L1 Data Cache (L1)
- ECX=2: L2 Data Cache (L2)

Table B.7: Information Returned by CPUID Instruction. 8000000xH leafs. (Contd.)

Operation

```
IA32_BIOS_SIGN_ID MSR = Update with installed microcode revision number;
case (EAX)
{
 EAX == 0:
   EAX = 01H; \frac{1}{2} // Highest basic function CPUID input value
   EBX = "Genu";ECX = "ineI";EDX = "ntel";break;
 EAX = 2H:
   // Cache and TLB information
   EAX = 0;
   EBX = 0;ECX = 0;EDX = 0;break;
 EAX = 3H:
   // PSN features
   EAX = 0;EBX = 0;ECX = 0;EDX = 0;break;
 EAX = 4H:
```


```
// Deterministic Cache Parameters Leaf;
 EAX = *see table*EBX = *see table*ECX = *see table*EDX = *see table*break;
EAX = 20000000H;EAX = 01H; // Reserved
 EBX = 0; // Reserved
 ECX = 0; // Reserved
 EDX = 0; // Reserved
 break;
EAX = 20000001H;EAX = 0; // Reserved
 EBX = 0; // Reserved
 ECX = 0; // Reserved
 EDX = 00000010H; // Reserved
 break;
EAX = 80000000H;// Extended leaf
 EAX = 08H; // Highest extended function CPUID input value
 EBX = 0; // Reserved
 ECX = 0; // Reserved
 EDX = 0; // Reserved
 break;
EAX = 80000001H;EAX = 0; // Reserved
 EBX = 0; // Reserved
 ECX[0] = 1; // LAHF/SAHF support in 64 bit mode
 ECX[31:1] = 0; // Reserved
 EDX[10:0] = 0; // Reserved
 EDX[11] = 1; \frac{1}{2} // SYSCALL/SYSRET available in 64 bit mode<br>EDX[19:12] = 0; \frac{1}{2} Reserved
 EDX[19:12] = 0;EDX[20] = 0; // Execute Disable Bit available
 EDX [28:21] = 0; // Reserved
 EDX[29] = 1; // Intel(R) 64 Technology availableEDX[31:30] = 0; // Reserved
 break;
EAX = 80000002H;EAX = 0; // Processor Brand String
 EBX = 0; \frac{1}{2} Processor Brand String Continued
 ECX = 0; \frac{1}{2} // Processor Brand String Continued
 EDX = 0; // Processor Brand String Continued
 break;
EAX = 80000003H;
 EAX = 0; \frac{1}{2} Processor Brand String Continued
 EBX = 0; // Processor Brand String Continued
 ECX = 0; \frac{1}{2} // Processor Brand String Continued
 EDX = 0; \frac{1}{2} // Processor Brand String Continued
 break;
EAX = 80000004H;
 EAX = 0; \frac{1}{2} Processor Brand String Continued
```


```
EBX = 0; \frac{1}{2} Processor Brand String Continued
  ECX = 0; // Processor Brand String Continued
  EDX = 0; \frac{1}{2} Processor Brand String Continued
  break;
EAX = 80000005H;EAX = 0; // Reserved
  EBX = 0; // Reserved
  ECX = 0; // Reserved
  EDX = 0; // Reserved
  break;
EAX = 80000006H;
  EAX = 0; // Reserved
  EBX = 0; // Reserved
  ECX[7:0] = 64; // L2 cache Line size in bytes
  ECX[15:12] = 6; // L2 associativity field (8-way)ECX[31:16] = 256; // L2 cache size in 1K unitsEDX = 0; // Reserved
  break;
EAX = 80000007H;
  EAX = 0; // Reserved
  EBX = 0; // Reserved
  ECX = 0; // Reserved
  EDX = 0; // Reserved
  break;
EAX = 80000008H;
  \text{EAX}[7:0] = 40; // Physical Address bits
  \text{EAX}[15:8] = 48; // Virtual Address bits
  EAX[31:16] = 0; // Reserved
  EBX = 0; // Reserved
  ECX = 0; // Reserved
  EDX = 0; // Reserved
  break;
default, EAX == 1H:
  EAX[3:0] = Stepping ID;
  EAX[7:4] = *see table* // Model
  EAX[11:8] = 1011B; // Family
  \text{EAX}[13:12] = 00B; // Processor type
  EXAMPLE A X [15:14] = OOB; // Reserved
  EAX[19:16] = 0000B; // Extended Model
  EAX[23:20] = 00000000B; // Extended Family
  EAX[31:24] = 00H; // Reserved;
  EBX[7:0] = 00H; // Brand Index (* Reserved if the value is zero *)
  EBX[15:8] = 8; \frac{1}{5} // CLEVICT1/CLFLISH Line Size (x8)<br>
EBX[23:16] = 248; \frac{1}{5} // Maximum number of logical proce
                       // Maximum number of logical processors
  EBX[31:24] = Initial Apic ID;
  ECX = 00000000H; // Feature flags
  EDX = 110193FFH; // Feature flags
  break;
}
```


Flags Affected

None.

Intel® C/C++ Compiler Intrinsic Equivalent

None

Exceptions

None.

Table B.8: Feature Information Returned in the EDX Register (CPUID.EAX[01h].EDX)

Table B.9: Feature Information Returned in the EDX Register (CPUID.EAX[01h].EDX) (Contd.)

Table B.10: Feature Information Returned in the ECX Register (CPUID.EAX[01h].ECX)

*^a*CPUID bit 23 erroneously indicates that POPCNT is not supported. Knights Corner does support the POPCNT instruction. See Appendix[A](#page-621-0) for more information.

Appendix C

Floating-Point Exception Summary

C.1 Instruction floating-point exception summary

Table [C.3](#page-683-0) shows all those instruction that can generate a floating-point exception. Each type of exception is shown per instruction. For each table entry you will find one of the following symbols:

- Nothing : Exception of that type cannot be produced by that instruction.
- Y*both*: The instruction can produce that exception. The exception may be produced by either the operation or the data-type conversion applied to memory operand.
- Y*conv*: The instruction can produce that exception. That exception can only be produced by the data-type conversion applied to memory operand.
- Y*oper*: The instruction can produce that exception. The exception can only be produced by the operation. The data-type conversion applied to the memory operand cannot produce any exception.

APPENDIX C. FLOATING-POINT EXCEPTION SUMMARY

APPENDIX C. FLOATING-POINT EXCEPTION SUMMARY

C.2 Conversion floating-point exception summary

Out-of-range values are dependent on operation definition and rounding mode. Table [C.3](#page-686-0) and Table [C.4](#page-687-0) describe maximum and minimum allowed values for float to integer and float to float conversion respectively. Please note that presented ranges are considered after ``Denormals Are Zero (DAZ)'' are applied.

Those entries in Table [C.4](#page-687-0) labelled with an asterisk(*[∗]*), are not required for Knights Corner.

C.3 Denormal behavior

APPENDIX C. FLOATING-POINT EXCEPTION SUMMARY

APPENDIX C. FLOATING-POINT EXCEPTION SUMMARY

Table C.3: Float-to-integer Max/Min Valid Range Table C.3: Float-to-integer Max/Min Valid Range

Table C.4: Float-to-float Max/Min Valid Range Table C.4: Float-to-ϐloat Max/Min Valid Range

APPENDIX C. FLOATING-POINT EXCEPTION SUMMARY

Appendix D

Instruction Attributes and Categories

In this Appendix we enumerate instruction attributes and categories

D.1 Conversion Instruction Families

D.1.1 *Df*³² Family of Instructions

VMOVAPS VMOVNRAPS VMOVNRNGOAPS VPACKSTOREHPS VPACKSTORELPS VSCATTERDPS VSCATTERPF1DPS

D.1.2 *Df*⁶⁴ Family of Instructions

D.1.3 *D*_{i32} Family of Instructions

VMOVDQA32 VPACKSTOREHD VPACKSTORELD VPSCATTERDD

D.1.4 *Di*⁶⁴ Family of Instructions

VMOVDQA64 VPACKSTOREHQ VPACKSTORELQ VPSCATTERDQ

D.1.5 *S*_{*f*32} Family of Instructions

D.1.6 *S*_{*f*64} Family of Instructions</sub>

APPENDIX D. INSTRUCTION ATTRIBUTES AND CATEGORIES

D.1.7 *Si*³² Family of Instructions

D.1.8 *S*_{*i*64} Family of Instructions</sub>

VFIXUPNANPD VPANDNQ VPANDQ VPBLENDMQ VPORQ VPXORQ

D.1.9 *U*_{f32} Family of Instructions

D.1.10 *U*^{*f*64} Family of Instructions</sup>

D.1.11 *U*_{i32} Family of Instructions

D.1.12 *U*_{*i*64} Family of Instructions</sub>

Appendix E

Non-faulting Undefined Opcodes

The following opcodes are non-faulting and have undefined behavior:

- MVEX.512.0F38.W0 D2 /r
- MVEX.512.0F38.W0 D3 /r
- MVEX.512.0F38.W0 D6 /r
- MVEX.512.0F38.W0 D7 /r
- MVEX.512.66.0F38.W0 48 /r
- MVEX.512.66.0F38.W0 49 /r
- MVEX.512.66.0F38.W0 4A /r
- MVEX.512.66.0F38.W0 4B /r
- MVEX.512.66.0F38.W0 68 /r
- MVEX.512.66.0F38.W0 69 /r
- MVEX.512.66.0F38.W0 6A /r
- MVEX.512.66.0F38.W0 6B /r
- MVEX.512.66.0F38.W0 B0 /r /vsib
- MVEX.512.66.0F38.W0 B2 /r /vsib
- MVEX.512.66.0F38.W0 C0 /r /vsib
- MVEX.512.66.0F38.W0 D2 /r
- MVEX.512.66.0F38.W0 D6 /r
- MVEX.512.66.0F3A.W0 D0 /r ib
- MVEX.512.66.0F3A.W0 D1 /r ib
- MVEX.NDS.512.66.0F38.W0 54 /r

- MVEX.NDS.512.66.0F38.W0 56 /r
- MVEX.NDS.512.66.0F38.W0 57 /r
- MVEX.NDS.512.66.0F38.W0 67 /r
- MVEX.NDS.512.66.0F38.W0 70 /r
- MVEX.NDS.512.66.0F38.W0 71 /r
- MVEX.NDS.512.66.0F38.W0 72 /r
- MVEX.NDS.512.66.0F38.W0 73 /r
- MVEX.NDS.512.66.0F38.W0 94 /r
- MVEX.NDS.512.66.0F38.W0 CE /r
- MVEX.NDS.512.66.0F38.W0 CF /r
- MVEX.NDS.512.66.0F38.W1 94 /r
- MVEX.NDS.512.66.0F38.W1 CE /r
- VEX.128.F2.0F38.W0 F0 /r
- VEX.128.F2.0F38.W0 F1 /r
- VEX.128.F2.0F38.W1 F0 /r
- VEX.128.F2.0F38.W1 F1 /r
- VEX.128.F3.0F38.W0 F0 /r
- VEX.128.F3.0F38.W1 F0 /r

Appendix F General Templates

In this Chapter all the general templates are described. Each instruction has one (at least) valid format, and each format matches with one of these templates.

F.1 Mask Operation Templates

Mask m0 - Template

VMASKMask m0

(intel) Mask m1 - Template

Mask m2 - Template

VMASKMask m2

(intel) Mask m3 - Template

Mask m4 - Template

VMASKMask m4

(intel) Mask m5 - Template

APPENDIX F. GENERAL TEMPLATES

F.2 Vector Operation Templates

(intel) Vector v0 - Template

(intel)

(intel) Vector v1 - Template

Vector v10 - Template

APPENDIX F. GENERAL TEMPLATES

Vector v11 - Template

Vector v2 - Template

(intel)

APPENDIX F. GENERAL TEMPLATES

6 5 4 3 2 1 0

(intel) Vector v3 - Template

Vector v4 - Template

APPENDIX F. GENERAL TEMPLATES

Vector v5 - Template

APPENDIX F. GENERAL TEMPLATES

Vector v6 - Template

Vector v7 - Template

(intel)

Vector v8 - Template

APPENDIX F. GENERAL TEMPLATES

Vector v9 - Template

APPENDIX F. GENERAL TEMPLATES

F.3 Scalar Operation Templates

Scalar s0 - Template

Description

(intel) Scalar s1 - Template

Description

